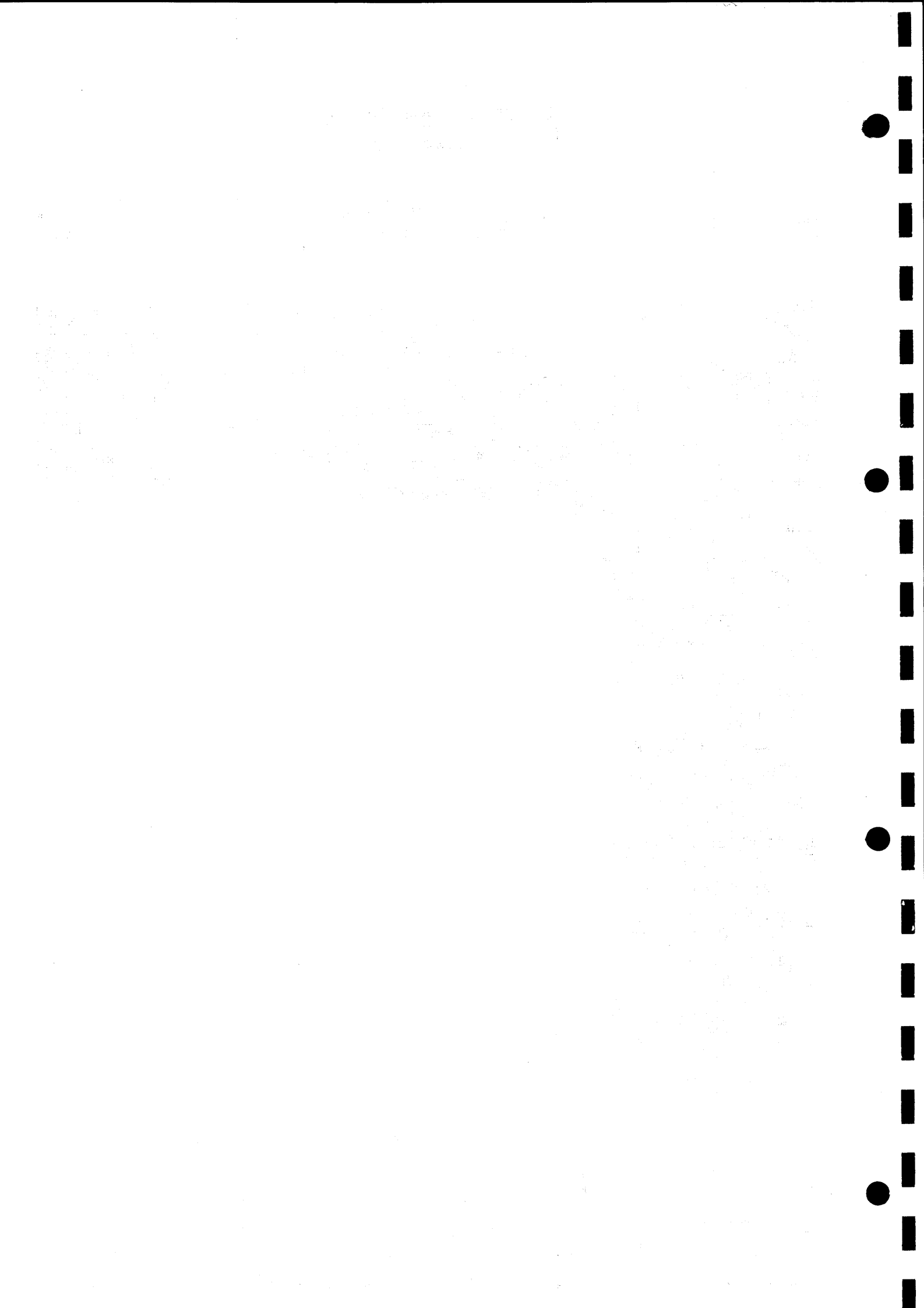


**MUSIL**

**Programming Guide**

**RF**

**36000**



# MUSIL Programming Guide

A/S REGNECENTRALEN  
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Author and  
Text Editor:           Joan Rosenstein

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ABSTRACT:            This manual shows how to program an RC 3600 in the MUSIL  
high-level language.

SUPPORTING DOCUMENT:  
                          RC 3600 MUSIL Programmer's Reference Card (RCSL 42 - i 0355)

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**Part One**  
**Basic System Architecture**



## 1.1 Introduction

In this Part One, we shall see how the RC 3600 system is built up and operated. We shall also take a brief overall look at the applications software available on the system.

## 1.2 The RC 3600 System

The RC 3600 is a system that is composed of hardware and systems software. The hardware is a 16-bit per word mini-computer which is particularly reliable and sturdy, but not optimally flexible nor specifically designed for the support and terminal functions that an RC 3600 system is designed to perform. This flexibility and suitability is provided by the systems software.

The RC 3600 system is a satellite system. That means that it is ultimately associated with a larger host computer. The job of the host computer in an RC 3600/mainframe configuration is to perform computation and data management, often in the form of data base or file management. The job of the RC 3600 in this configuration is to take over those functions that can be separated from the mainframe, such as I/O, peripheral management, data entry and collection, data conversion, and various tasks associated with communications.

The RC 3600 system can be used with or without standard or custom-made applications programs supplied by Regnecentralen. For those users who wish to program their own applications, a special high-level programming language is provided. This language, called MUSIL, is designed for programming support functions. Therefore, it is strong on I/O handling and weak on computational facilities.

The central idea in RC 3600 structure is to complement each element of hardware with systems software. We shall first survey these hardware elements.

### 1.2.1 RC 3600 Hardware

RC 3600 hardware consists of a central unit with its associated core memory, peripheral units with their associated channels and controllers, a real-time clock, an inter-

rupt system, and a bootstrap loader. With this equipment alone, one can process only one job at a time, each sort of I/O device must have its own data transfer protocol, and high-level programming languages cannot be used. There is also a direct memory access channel for use by the faster peripherals to access core without going through the central processor, but the RC 3600 hardware alone cannot utilize this speed maximally, because only one job can proceed at a time with this hardware, which means that the whole machine is limited in speed by the slower peripherals in I/O-oriented tasks, which are the main tasks of the RC 3600. To solve these problems, various software elements are used to complement each hardware element.

### 1.2.2 RC 3600 Systems Software

RC 3600 systems software is composed of a monitor, the systems process S, I/O utility routines, and peripheral device drivers.

- 1.2.2.1 The Monitor has the job of implementing multiprogramming. It does this by complementing the hardware interrupt system with a software interrupt system utilizing a software real-time clock. It provides the means for each process to communicate with the others. It can, therefore, provide a means for several different jobs to be executed at the same time by providing each of them in turn with time slices regulated by the clock. In this way more than one job can run at the same time and slow I/O processes do not slow down the overall performance of the system.
- 1.2.2.2 The Systems Process S implements a software core allocation system. It creates process descriptors for each process it loads, and these descriptors can be used by the monitor in its task of mediating information among the various processes. S also replaces the buttons and switches that give the operator direct access to the hardware with an operating system that gives the operator access to the system as such.
- 1.2.2.3 The Drivers replace the individual peculiarities of the various peripheral device types with a single I/O protocol, enabling the monitor to treat I/O processes on the same level with any other processes. Those routines that all the drivers use in common are gathered together in the I/O utility procedures, so as to make each driver as small as possible.

## 1.3 RC 3600 Applications Software

Most RC 3600 users require one or more standard or custom-made applications programs from Regnecentralen. For such customers full job runs are usually provided on one medium, e.g., on one magnetic tape or in one card deck. These runs contain both the systems and the applications software. Other customers require their own programming capability. For them it is necessary to have a run medium with systems software and a program production package.

### 1.3.1 Applications Run Media

Those using ready-made programs receive a medium with a monitor, a "basic system", I/O utility routines, drivers, and one or more applications programs in object code, along with an interpreter to execute the object code.

1.3.1.1 Basic Systems consist of a systems process S, a console device driver, and an autoloader device driver. The last is necessary because S must have a device driver from which to load application modules. The console device driver can be of two sorts, a driver for the Operator Control Panel (OCP) or a driver for a keyboard device. Basic systems have been developed to keep the space needed for systems software as small as possible.

### 1.3.2 Program Production Packages

These packages can be supplied separately or on a medium with the necessary systems software, which in this case also must include an interpreter. The package itself consists of a MUSIL compiler, a MUSIL Text Editor, any necessary drivers, and a program generator.

1.3.2.1 The MUSIL Compiler converts MUSIL source code into MUSIL object code, which is executed by a separate piece of software called the MUSIL Interpreter.

The MUSIL compiler can also integrate assembly-coded subroutines, called "code procedures" into the object code output from the compiler. In addition, it can copy parts of MUSIL programs into locations in other MUSIL programs.

1.3.2.2 MUSIL Text Editor allows editing on MUSIL source text by character, string, line, or page. It operates by taking in the source text as data to itself.

- 1.3.2.3 Program Generators combine one or more compiled MUSIL programs onto a single run medium along with the necessary systems software. They can also integrate into the run "command files", which are files containing code that substitutes for direct operator action, which makes a run more automatic than it would ordinarily be.

## 1.4 RC 3600 Operation

RC 3600 operation is both simple and flexible. Standard programs can be run automatically by a few simple commands or can be closely controlled by the operator through decision points in the programs that give rise to requests for the input of "runtime parameters" by the operator. Error messages are also complete and easily understood. The programmer who wishes to operate an RC 3600 system should obtain an RC 3600 Operator's Guide and/or an RC 3600 Data Conversion Operator's Reference Card.

Program writing can also be done at the RC 3600 console. Programmers who wish to create or edit MUSIL programs at the machine should also provide themselves with an RC 3600 MUSIL Programmer's Reference Card, which contains information on program writing, compilation, and editing, as well as program production through run generation.

### 1.4.1 Operation Principles

The operator may communicate with any loaded process, including the systems process S. Messages to and from S are indicated on a keyboard device by CTRL G (BELL) and S and the symbol >, respectively, and on the OCP by the LOAD button and LOAD lamp.

To load a medium, the command LOAD is used, and to select a file from a loaded medium, the command INT, meaning "interpret", is used.

The word "file" is also used to indicate a device whose medium does not contain a catalog. Thus, the paper tape reader is operated as a file. On the other hand, a disc contains, normally, more than one file and a catalog, so that in operating with a disc one must specify which file on the disc one wishes to access. This use of "file" simplifies RC 3600 operation for non-cataloged media.

Most standard programs can be operated from an OCP, but more complicated programs and program writing require a keyboard device. The user of an RC 3600 system with

an OCP can request from Regnecentralen one or more command files to make the use of his machine more flexible. For example, with an OCP all programs must be loaded from the autoload device, but with the proper command files another device can be used for program load, imitating this capability of the keyboard devices.







**Part Two**  
**The MUSIL Program**



## 2.1 MUSIL Program Structure

MUSIL programs must be written in modular form. Comments may be placed anywhere in the program, as long as they are placed outside of words, numbers, and text strings. (It is, of course, most common to place them between statements or between rows of a table.) Comments are signaled by exclamation points, thus:

! THIS IS A COMMENT !

MUSIL programs are presented in five sections, which must be given in the following order:

The Constant Section is the first section. It is normally present. In it are defined all the constants that will be used in the program, both numerical constants and text string constants. Constants may not be defined anywhere else in a MUSIL program.

The Type Section may be absent. If it is not, then it must be the second section. In it are defined types, or categories, of variables. It is used mostly for file type definitions, for file types are usually long. This section provides only a convenient short-hand type of definition for variable structures. It cannot itself define variables.

The Variable Section is used to define variables of both numerical and text string type, as well as records and files. Variables cannot be defined anywhere else in the program.

The Procedure Section is, properly-speaking, a part of the Main Program Section that follows, but it is convenient to speak about it as a separate section. In it the programmer can define his own procedures, usually I/O exception-handling procedures, that he can call later on from the Main Program Section by name.

The Main Program Section contains the program's instructions.

MUSIL commands can be viewed as of two types: I/O commands and other commands. In this Part Two we shall discuss only the latter. I/O commands are discussed in Part Three.

## 2.2 The Constant Section

In the Constant Section we define numerical or text string constants as well as tables of numerical or text string constants. Conversion tables are often found in this section, but conversion tables can also be made into separate programs. (Certain conversion table programs can be obtained from Regnecentralen.)

### 2.2.1 Identifiers

The definition of a constant has the following form:

name = value,

The "name" may consist of any sequence of letters of the alphabet or numerals, but it must begin with a letter and may not include characters other than letters and numbers. The name may be of any length, but only the first 7 characters and the total character count will be used by the program. That is, to the program the following names are the same

MYNAME11      and      MYNAME13

but

MYNAME11      and      MYNAME111

are different.

A "value" may be a number, a text string, or a table of numbers or text strings.

### 2.2.2 Numbers

A number may be assigned to an identifier (a "name"). The number may be decimal, octal, or binary; but whatever the number, its binary equivalent must not exceed 16 bits. For decimal numbers this means that the number must fall between

-32768      and      +32767

For octal numbers the range is

0            to            8'177777

Octal numbers are generally used to express bit patterns conveniently.

Numbers may have spaces between a sign and the numerals, but there must not be spaces within the number itself. Identification of numbers is done as follows:

NUM1 = 2'011001,	! A BINARY NUMBER !
NUM2 = 8'775,	! AN OCTAL NUMBER !
NUM3 = -23005,	! A DECIMAL NUMBER !

In the absence of a sign, a positive sign will be assumed.

Decimal points may not be used within numbers.

### 2.2.3 Text Strings

Strings representing ASCII texts may be assigned a name. This is done according to the following model:

```
name = 'text string',      or
name = "text string",     or
name = "text string",     or
name = 'text string',
```

Either single or double quotation marks may be used, and mixing single and double quotation marks within the same identification is also allowed. An example of the identification of a text string might be

```
ALPHA22 = "ERROR HAS OCCURRED",
```

After ALPHA22 has been thus defined, it can be used as the name of the text string ERROR HAS OCCURRED.

Text strings cannot be operated on arithmetically. Once they have been named, however, they can be

```
assigned to a variable as its current value,
used in text comparisons,
output on the operator console,
output to a peripheral device.
```

Text strings may include byte values given by their numeric representation. One common use of this facility is to output control characters, for example,

```
ALPHA = '<45>',
```

places the binary value of decimal 45 into location ALPHA. The angular parentheses indicate a byte value. Since the ASCII code for decimal 45 is a minus sign, ALPHA can now be regarded as containing a textual minus sign.

If we write, for example,

```
V = "<8'126>",
```

then the ASCII code for the letter V can be retrieved from location V for later use.

As another example, if we write

```
CR = "<13>",
```

then the ASCII code for a Carriage Return can be retrieved from location CR for use in controlling a line printer.

Strings of such ASCII characters can also be placed together under one "name", so that they can be called together as a sequence of ASCII characters. If we write, for example,

```
ALPHA31 = "<45><0><10>",
```

then ALPHA31 will contain the ASCII codes for

```
Minus sign  NUL  Line feed
```

No punctuation is used between the factors of such a string of ASCII characters.

#### 2.2.4 Tables of Constants

Tables of numerical or text string constants can be defined in the Constant Section. Tables are in fact text strings themselves, so that their elements cannot be operated on arithmetically in a direct way.

Tables are identified as follows:

```
name = # element1 blank element2 blank ....#,    or
name = "<element1><element2> ....." ,
```

where the quotation marks may, again, be single or double, or any mixture of single or double quotation marks. The blanks denote spaces and/or new lines.

As examples we can take the following:

```
LPTTABLE = #14 0 64 89 56 8'377 0 65#,
LPTTABLE = "<14><0><64><89><56><8'377><0><65>",
```

which are equivalent definitions for the same table identification.

Note that the following definitions are absolutely equivalent:

```
ALPHA = #45#,    and    ALPHA = "<45>",
```

Table definitions are, therefore, very useful for writing device conversion tables and can include all characters, as well as control characters.

### 2.2.5 Section Structure

The Constant Section begins with the key word

```
CONST
```

which is not followed by any punctuation. Directly after CONST come the desired identifications in any order. Each identification is followed by a comma, except for the last identification, which is followed by a semicolon, thus:

```
CONST
      ALPHA  = 45,
      BETA1  = -8'377,
      GAMMA  = "MOUNT TAPE";
```

Because the whole MUSIL program can be regarded as one single compound statement, identifications may be entered with any spacing. Blanks will be ignored, as long as they do not occur within names or values.

### 2.2.6 Some Cautions On the Use of Constants

Values are stored in their named locations left-justified. Numerical constants are stored as given in the program. String constants are marked by a binary zero at the end of the text. When text string are read out to another location or to a peripheral device, this binary zero is stripped off. If, therefore, a text string is assigned to a variable during program execution and then output on a console device, the console will not stop printing at the end of the text, unless a binary zero has been appended to the text by the programmer.

To avoid such annoying occurrences, a binary zero should be placed at the end of each text string that will be operated on in some way before being output to the console. This is done thus:

```
ALPH = "OUTPUT THIS TEXT<0>",
```

As an example of what else can happen if this is not done, suppose I have assigned the text THIS MESSAGE IS WRONG to ALPHA20. Suppose that later on in the program I wish to change the contents of ALPHA20 to THIS IS ALPHA and then output the contents of ALPHA20. What I want to output is

```
THIS IS ALPHA
```

but what I will get is

```
THIS IS ALPHAIS WRONG
```

Using the binary zero after the text strings would solve this problem.

String constants that will not be moved around in the program do not, of course, need to conclude with this explicit binary zero.

Text strings may mix ASCII characters and byte values thus:

```
MTTEXT = "<7<10>MT ERROR",
```

The free use of blanks allows long tables to be presented in a way that makes them easy to check. One neat way to insert tables is the following.

```
LPTTABLE = # 0  0  0  0  0  0  0  0  0  0  !0th !
              0  0  0  0  0  0  0  0  0  0  !9th !
              33 34 35 36 37 38 39 40 41  !18th !
              45 46 47 48 49 50 51 52 53  !27th !
              60 61 62 63 64 65#          !36th !
```

where the numbers in exclamation points are comments that give the ordinal position of the first number in each line.

### 2.2.7 Common Errors

The most common errors in the Constant Section are

- |                                    |  |
|------------------------------------|--|
| 1) CONST,                          | Punctuation after CONST.   |
| 2) ALP = "2"                       | No comma after an identification.  |
| 3) LASTONE = 5,                    | Comma after the last definition in the section.                                      |
| 4) ANY = 7;                        | Semicolon after an identification that is not the last.                              |
| 5) GA.1 = 2,                       | A name that does not consist of only numbers and letters.                            |
| 6) 1FIVE = "A",                    | A name that begins with a number.  |
| 7) TW 00 = 2,                      | A blank within a name.   |
| 8) J = 3.4,                        | Presence of a decimal point.   |
| 9) K = 55 6,                       | Presence of a blank within a value.  |
| 10) L = 2,330,                     | Presence of a comma within a value.  |
| 11) B = "<8'128>",                 | Wrong ASCII value.   |
| 12) G = "<7>,<10>",                | Comma between ASCII values of a text string.   |
| 13) TABLE = # 0,33,37#             | Commas within a table.   |
| 14) ALPHA134 = 5,<br>ALPHA137 = 6, | Doubly-defined name (program disregards all characters of a name after the first 7). |
| 15) HIGH = 50000                   | Value too large.   |





```

0165
0166 CRTABLE=          ! CR CONTROLLER FORMAT TO ERCDIC
0167          0      1      2      3      4      5      6      7 !
0168 #
0169 !   0 !   64 241 242 243 244 245 246 247
0170 !   8 !  249 49  50  51  52  53  54  55
0171 !  16 !  248 121 122 123 124 125 126 127
0172 !  24 !   56 57 58 59 60 61 62 63
0173 !  32 !  240 97 226 227 228 229 230 231
0174 !  40 !  233 33 34 35 36 37 38 39
0175 !  48 !  232 105 224 107 108 109 110 111
0176 !  56 !   40 41 42 43 44 45 46 47
0177 !  64 !   96 209 210 211 212 213 214 215
0178 !  72 !  217 17 18 19 20 21 22 23
0179 !  80 !  216 89 90 91 92 93 94 95
0180 !  88 !   24 25 26 27 28 29 30 31
0181 !  96 !  208 161 162 163 164 165 166 167
0182 ! 104 !  169 225 98 99 100 101 102 103
0183 ! 112 !  168 160 170 171 172 173 174 175
0184 ! 120 !  104 32 234 235 236 237 238 239
0185 ! 128 !   80 193 194 195 196 197 198 199
0186 ! 136 !  201 1 2 3 4 5 6 7
0187 ! 144 !  200 73 74 75 76 77 78 79
0188 ! 152 !   8 9 10 11 12 13 14 15
0189 ! 160 !  192 129 130 131 132 133 134 135
0190 ! 168 !  137 65 66 67 68 69 70 71
0191 ! 176 !  136 128 138 139 140 141 142 143
0192 ! 184 !   72 0 202 203 204 205 206 207
0193 ! 192 !  106 145 146 147 148 149 150 151
0194 ! 200 !  153 81 82 83 84 85 86 87
0195 ! 208 !  152 144 154 155 156 157 158 159
0196 ! 216 !   88 16 218 219 220 221 222 223
0197 ! 224 !  112 177 178 179 180 181 182 183
0198 ! 232 !  185 113 114 115 116 117 118 119
0199 ! 240 !  184 176 186 187 188 189 190 191
0200 ! 248 !  120 48 250 251 252 253 254 255
0201 #;
0202

```

## 2.3 The Variable Section

While the Constant Section assigns names to values, the Variable Section names locations in core for the values that the program will use. In order to do this, the program must be told what sort of values will fill these locations.

### 2.3.1 Identifiers

Variables may be integers, text strings, files, or records. The format for a variable definition is

```
name:type;   or   name, name, . . . ., name:type;
```

The restrictions on "name" are the same as for the Constant Section. "Type" is declared as in the following paragraphs.

File and record variable definition will be discussed in Part Three.

### 2.3.2 Integer Variables

If we wish to declare a location in such a way that integers can be stored in it, then we write

```
name:INTEGER;
```

For example,

```
D:INTEGER;           or
I:INTEGER;           or
NUMBER5:INTEGER;
```

Blanks may be freely used after or before the colon, e.g.,

```
FIRST :  INTEGER;
NEXT  :  INTEGER;
```

so as to give the coding a pleasing appearance and make it easy to read.

Allowable numerical values that can go into these locations are between

```
-32768   and   32767
```

### 2.3.3 Text String Variables

Locations can also be declared for text strings. In this case the number of bytes to be reserved must be declared. The format is

```
name:STRING(number of bytes);
```

For example,

```
TEXT1      :   STRING(20);      or  
TXT,TEXT,A :   STRING(6);
```

Text strings assigned later on in the program that are smaller than the number of bytes declared for their location will be left-justified. Implicit binary zeroes will not be assigned.

### 2.3.4 Section Structure

The Variable Section begins with the word

```
VAR
```

which is not followed by punctuation. Each declaration in this section ends with a semicolon, including the last declaration.

### 2.3.5 Some Cautions On the Use of Variables

There are only a few common mistakes in the Variable Section:

```
An incorrect name  
A punctuation error
```

2.3.6 Variable Section Example

```

0203
0204 VAR
0205
0206 OPDUMMY:          STRING(2);          ! RUNTIME PARAMETERS !
0207 PROGNO:          INTEGER;
0208 BLOCKNO:         INTEGER;
0209 FILENO:          INTEGER;
0210 REWIND:           INTEGER;
0211 FIXRECS:         INTEGER;
0212 MAXCOL:           INTEGER;
0213 MINCOL:           INTEGER;
0214 BLOCKED:          INTEGER;
0215
0216 OPTXT:            STRING(20);         ! COMMUNICATION AREA !
0217 OPSTRING:        STRING(20);
0218 OPDEC:            STRING(10);
0219
0220 OPCONT:           STRING(2);          ! INTERNAL VARIABLES !
0221 NEXTCONT:         STRING(1);
0222 GLCONT:           STRING(1);
0223 WBLOCKED:        INTEGER;
0224 ERRORNO:         INTEGER;
0225 MASK:             INTEGER;
0226 TOM:              INTEGER;
0227 SIGN:             INTEGER;
0228 Q:                INTEGER;
0229 PAR:              INTEGER;
0230 LENGTH:           INTEGER;
0231 OPENED:           INTEGER;
0232 P1:               INTEGER;
0233 P2:               INTEGER;
0234 P3:               INTEGER;
0235 S1:               STRING(2);
0236 S2:               STRING(2);
0237 NEXTMT:           INTEGER;
0238 INLENGTH:         INTEGER;
0239 OUTLENGTH:        INTEGER;
0240 CARDSREAD:         INTEGER;
0241 SAVEDSUSPEND:     INTEGER;
0242

```

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## 2.4 The Procedure Section

The Procedure Section has no key word to begin it. It consists only of procedure definitions, one after another.

### 2.4.1 Defining a Procedure

In the Procedure Section the programmer defines his own procedures. This is done according to the following format:

```

PROCEDURE name;
  BEGIN
    ..... ;
    ..... ;
    .....
  END;
```

Within the BEGIN ..... END described above may be found any of the sort of statements that can be used in the Main Program Section. The Procedure Section is useful for defining what shall occur in the case of an I/O exception situation.

The variables and constants used within the Procedure Section must have been previously defined in their appropriate sections.

### 2.4.2 Executing a Procedure

To start the execution of a procedure from within the Main Program Section, one writes simply

```

procedure name;
```

For example, suppose we have within the Procedure Section

```

PROCEDURE ENDGAME;
  BEGIN
    ..... ;
  END;
```

then, to call this procedure from the Main Program Section, we write

```

ENDGAME;
```

simply. Procedures cannot be called by other procedures. *before they are defined.*

### 2.4.3 Code Procedures

To incorporate code procedures within a MUSIL program during compilation, the programmer must indicate in his program where the MUSIL compiler is to put these procedures. One does this by writing in the Procedure Section

```
PROCEDURE name (parameter specification1, ...  
                ..., parameter specification5)  
CODEBODY external identification;
```

The parameter specification and the external identification can be obtained from Regnecentralen.

To call the code procedure from the Main Program Section, one then writes

```
name (parameter1, ..., parameter5)
```

Code procedures and instruction in their use are supplied by Regnecentralen.

### 2.4.4 Some Cautions On the Use of Procedures

The most common errors are

Forgetting to define the procedure's constants and variables in the Constant and Variable Sections, respectively.

Trying to jump from a point within one procedure to a point within another procedure.

Trying to call a procedure from another procedure before the first procedure has been defined.

Trying to call procedures recursively.

### 2.4.5 Examples of Procedures

Such examples will be given in Part Three.

## 2.5 The Main Program Section

The statements that actually control a job are found in the Main Program Section.

### 2.5.1 Section Structure

The Main Program Section begins with the key word

BEGIN

not followed by any punctuation. It ends with the word

END;

Statements in this section are separated from one another by semicolons. Spaces may be freely used between words.

### 2.5.2 Arithmetic Operators

MUSIL includes the following arithmetic operators:

+ addition  
 - subtraction  
 \* multiplication  
 / division

and parentheses may be used freely.

Arithmetic operations are executed from left to right in order of their priority, which is (from high to low)

plus and minus signs  
 multiplication and division  
 addition and subtraction

Thus  $-5+6*7-2/3$   
 is equivalent to  $-5 + (6 * 7) - (2/3)$

### 2.5.3 Relational Operators

The available relational operators are

<	less than	>=	greater than or equal to
>	greater than	=	equal to
<=	less than or equal to	<>	not equal to

These symbols can also be used for text string comparisons, in which case the strings will be compared lexicographically, that is, on the basis of their numerical ASCII values.



#### 2.5.4 Monadic Operators

There are four monadic operators:

+ number	The plus sign
- number	The minus sign
BYTE textstring	The integer value of the first character of the text string
WORD textstring	The integer value of the first two characters of the text string

For example, suppose TXT contains

"<2'11001001>2'11110011>"

which is equivalent to "A", then

BYTE TXT	gives the integer	0000000011001001
WORD TXT	gives the integer	1100100111110011

#### 2.5.5 Logical Operators

There are three logical operators:

operand 1 AND operand 2  
yields the integer value of the logical AND operation as performed on the current value of the two operands

operand 1 SHIFT operand 2  
shifts the value of the first operand to the left if the second operand is positive, and to the right if the second operand is negative, shifting the number of bits equal to the numerical value of the second operand; the shift is not cyclical: bits shifted out of the word are lost and the vacant positions are filled with zeroes

operand 1 EXTRACT operand 2  
extracts bits from the first operand; the number of bits extracted is equal to the current numerical value of the second operand

*from right* ↙

The operands involved in the logical operators must be integral and the result will be an integer.

For example, let VAR1, VAR2, A, and INT be integer variables, and let the current values be

```
VAR1    2'0000000010011011
VAR2    2'0000000011100000
INT     2'1111000010111111
A       2'1111000000001111
```

then

```
VAR1 AND VAR2    gives    0000000010000000
A SHIFT 2        gives    1100000000111100
A SHIFT (-2)     gives    0011110000000011
INT EXTRACT 8    gives           10111111
```

### 2.5.6 Operators In General

The priority from high to low for all operators is

```
Monadic operators
Multiplying and logical operators
Adding operators
Relational operators
```

Division of a number by zero, or division of zero by zero, will not give rise to an error message. The result of such operations will be

-1

No indication of integer overflow is given.

When text strings are compared, the comparison will take place only on that number of characters that is the smaller of the two text strings, for example,

```
Let ALPHA, declared of length 2, contain TR
and let BETA, declared of length 5, contain TRANS
Then, the relation
```

ALPHA < BETA

could give misleading results, for the comparison will take place only on the first two characters, which are TR in both cases.

Note that the following is not allowed:

"text string" operator "text string"

for operators can operate only on named values.

2.5.7 Assignment

The symbol used for assignment is := so that

```
A:=B;
```

assigns the current value of B to A.

Integer values can be assigned directly, thus:

```
INT1:=5;
```

as long as INT1 had been declared as an integer variable:

```
INT1:INTEGER;
```

But text strings cannot be directly assigned. That is, it is not allowed to write

```
TEXT3:="REWIND TAPE";
```

even if TEXT3 had been previously defined in the Variable Section as a text string variable. To put REWIND TAPE into TEXT3, one must have in the Variable Section something like

```
TEXT3:STRING(11);
```

and in the Constant Section

```
T3="REWIND TAPE",
```

and then one can write in the Main Program Section

```
TEXT3:=T3;
```

Then REWIND TAPE will be in TEXT3.

Variables defined as integers may have only integers assigned to them. And variables defined as text strings may have only text strings assigned to them. If, for example, INT1 is an integer variable and TEXT2 is a string variable, then one may not write either

```
INT1:=TEXT2;           or           TEXT2:=INT1;
```

Multiple assignments are not allowed. That is, one may not write

```
INT1, INT2:=0;         or           INT1:=INT2:=0;
```

When text strings are assigned, the number of characters that are moved is equal to the lesser of the number of characters defined for the variables involved. That is, if in the Variable Section we have

```
TEXT1:STRING(10);
TEXT2:STRING(20);
```

and in the Main Program Section we have

```
TEXT1:=TEXT2;           or           TEXT2:=TEXT1;
```

then, only ten characters will be moved in either case. In the first case only the first 10 characters of TEXT2 will be moved into TEXT1. In the second case the 10 charac-

ters in TEXT1 will be moved into the ten left-most positions of TEXT2, leaving the remainder of TEXT2 unchanged.

### 2.5.8 Labels

MUSIL statements may be labeled, so as to identify them. The label must be a unique numeric value between 0 and 65535. The format for labeling is

```
label: statement;
```

For example, 35: GAMMA:=5;

labels the statement GAMMA:=5; with the identifying label 35.

The use of spaces before or after the label's colon is optional.

Note that the statement

```
label: END;
```

must be preceded by a semicolon.

### 2.5.9 Compound Statements

The sub-statements of compound statements are set off by a BEGIN ... END combination. Thus:

```
BEGIN
    statement 1;
    statement 2;
    .....
    statement n
END;
```

There is no punctuation after the BEGIN or before the END (unless the END is labeled). Spaces may be freely used to improve readability.

BEGIN ... END phrases may be nested up to the number of 30.

### 2.5.10 Unconditional Branching

This is represented by the statement

```
GOTO label;
```

with no space within the GOTO.

If the GOTO statement is used to go to a labeled END statement, then the statement before the END statement must conclude with a semicolon, thus:

```

.....
.....
GOTO 60;
.....
.....
ALPHA:=40;
60:  END;
```

No error is committed, however, if the semicolon is used also in other cases.

### 2.5.11 Conditional Branching

MUSIL has the usual

```
IF relation THEN statement,
```

as well as an

```
IF relation THEN statement1 ELSE statement2
```

If the relation is true, then statement<sub>1</sub> will be executed. If not, then control will pass to the next statement. The IF-relation may be any allowable relation, and the THEN-statement may be any allowable statement, including compound statements. For example,

```

IF ALPHA=5 THEN
  BEGIN
    IF BETA<10 THEN U:=0;
    IF GA>= 9 THEN V:=1;
    .....
  END;
```

The relational expression may contain only variables or constants. The following are not allowed.

```

IF "TEXT1"="TEXT2" THEN ...   or
IF TEXT="REWIND" THEN ...
```

With respect to compound statements, note that in our example above, that if ALPHA was not 5, then program flow would have passed to the first statement after the END statement, bypassing the intermediate IF ... THEN statements.

2.5.12 Repetitive Statements

There are two repetitive statements in MUSIL:

```
WHILE ... DO
REPEAT ... UNTIL
```

The format of the first is

```
WHILE relation DO statement;
```

This instruction allows the repetition of an operation as long as the relational statement remains true, e.g.,

```
WHILE X>Y DO
  BEGIN
    ..... ;
    ..... ;
    .....
  END;
```

There is no punctuation after the DO. In the above example if X is never greater than Y, then the DO statement will never be executed.

The format of the second is

```
REPEAT statement UNTIL relation;
```

An example for this command might be

```
REPEAT
  BEGIN
    ..... ;
    ..... ;
    .....
  END
  UNTIL X=Y;
```

There is no punctuation after REPEAT or before UNTIL.

In this example, if X is equal to Y when END is reached, then the statement will be executed once.

### 2.5.13 Commands For the Operator Console

MUSIL contains 10 commands that are useful for communications between the program and the operator console. They are called Standard Procedures. Six of them will be described here, and the rest in Part Three.

OPMESS(string variable name);

outputs the text string contained in the string variable or constant named to the operator console. It continues to output the contents of the variable or constant until a binary zero <0> is reached. At most 80 bytes will be output in this way. If the text is less than 80 bytes and does not end in a binary zero, then the output will continue for the full 80 bytes anyway, outputting whatever is in core following the text desired. Text to be output by OPMESS should be in ASCII code.

Example:                   OPMESS(ALPHA);           where ALPHA contains a text string, will output the value of ALPHA on the operator console. Thus,

```
CONST
  ALPH="REWIND";
BEGIN
  OPMESS(ALPHA)
END
```

will output REWIND on the operator console.

OPIN(string variable name);

is the reverse operation. It allows the operator to insert a text string of up to 80 bytes into a variable previously defined as STRING in the Variable Section. In order to give the operator time to input this text, OPIN must be followed by

OPWAIT(integer variable name);

which makes the system wait for the operator's input. The number of characters that the operator actually inputs will be stored in the integer variable by the system.

```

Example:      VAR
              LENGTH:INTEGER;
              MAGTEXT:STRING(80);
              BEGIN
              OPIN(MAGTEXT);
              OPWAIT(LENGTH)
              END;

```

will allow the operator to put up to 80 bytes of text into the variable represented by `MAGTEXT`, while the `OPWAIT(LENGTH)` will give the operator time to do this and will store the number of characters input in `LENGTH`.

Operator input will normally be terminated, when a control key is used (CR, LF, ESCAPE, etc.).

If `OPIN` has been used, but it is desired to see if a text has actually been accepted for input from the console, then the command

#### OPTEST

can be written. If a text has been accepted for input, then `OPTEST` will give a non-zero value. If the `OPIN` operation has been unsuccessful, then the value of `OPTEST` will be zero.

```

Example:      VAR
              LENGTH:INTEGER;
              MAGTEXT:STRING(80);
              BEGIN
              OPIN(MAGTEXT);
              WHILE OPTEST = 0 DO
              BEGIN
              ..... ;
              .....
              END;
              OPWAIT(LENGTH);
              END;

```

will allow the operator to insert a text into `MAGTEXT`. If the input is successful, then `OPTEST` will be non-zero at some point and the `WHILE ... DO` statement will cease execution.



OPTEST is a standard function. Its current value can be used profitably to control program branching with respect to whether or not operator action has occurred, for example

```
IF OPTEST=0 THEN PSTOP;
```

where PSTOP is a user-defined procedure.

The RC 3600 system operates in binary. Thus all decimal numbers that are input by the operator with OPIN must be converted to binary before they can be used by the machine. This is done with

```
DECBIN(decimal value name, binary value name);
```

There must have been defined in the Variable Section two variables. One will be used to store the number inserted in decimal by the operator. The other will be used to store its binary equivalent. The variable with the binary equivalent is the one that must be used for all subsequent MUSIL statements that work with this inserted value.

```
Example:      VAR
                DEC:STRING(10);      ! NB: DECIMAL INPUT IS!
                LENGTH,BIN:INTEGER;  ! DEFINED AS A TEXT STRING!
            BEGIN
                OPIN(DEC);
                OPWAIT(LENGTH);
                DECBIN(DEC,BIN);
                IF BIN=0 THEN PSTOP;
            END;
```

The decimal value being converted must have no sign. It will be converted into a 16-bit binary value. There will be no check for overflow, so that the number must be less than or equal to 32767. If a non-numeric character appears within the input, then conversion will proceed up to that point and then stop. A plus or a minus sign or a decimal point is considered to be a non-numeric character.

```
BINDEC(binary value name, decimal value name);
```

is used with OPMESS. It takes the binary value and stores its ASCII equivalent concluding with a binary zero byte. The decimal value can then be output to the operator console by OPMESS.

```

Example:      VAR
              DEC:STRING(6);
              BIN:INTEGER;
              BEGIN
                BIN:=2'1001;
                BINDEC(BIN,DEC);
                OPMESS(DEC)
              END;

```

will output decimal 00009 to the operator console for inspection.

The decimal value variable must be defined in the Variable Section as a STRING with a minimum of 6 bytes. If it must be output with a sign, then the sign must be defined separately:

```

CONST
  PLUS="+";
VAR
  DECSIGN:STRING(1);
  DEC:    STRING(6);
  BIN:    INTEGER;
BEGIN
  BIN:=2'1001;
  BINDEC(BIN,DEC);
  DECSIGN:=PLUS;
  OPMESS(DECSIGN);
  OPMESS(DEC)
END;

```

which will output first + and then 00009 on the operator console.

(The binary value will be converted to exactly five decimal digits.)

To output +9 directly, one can make use of the instructions

MOVE and INSERT

The MOVE and INSERT commands are, thus, very useful in connection with commands to the operator console, but they are most often used for I/O. Therefore, they will be discussed in Part Three.

## 2.5.14 Some Cautions For the Main Program Section

The most common errors relating to the above material are

- 1) Forgetting a BEGIN or putting punctuation after it.
- 2) Forgetting the semicolon after an END.
- 3) Forgetting a semicolon after a statement that precedes a labeled END.
- 4) Forgetting the parentheses around a negative (i.e., right-hand) SHIFT.
- 5) Forgetting that bits are lost when SHIFT is used.
- 6) Mis-sequencing operators.
- 7) Dividing by zero inadvertently.
- 8) Comparing text strings of unequal length and forgetting that the comparison will be only on the lesser number of characters.
- 9) Trying to operate on text strings instead of on their names.
- 10) Forgetting the colon in the assign symbol.
- 11) Forgetting to define variables and/or constants before using them.
- 12) Trying to assign text strings to integer variables or integers to text string variables.
- 13) In assigning texts to string variables longer than the text, forgetting to take care of the excess text remaining from a previous text assignment.
- 14) Forgetting the colon after a label.
- 15) Spelling GOTO as two words.
- 16) Using constants within relational statements instead of variable or constant names.
- 17) Illogical entries into compound statements.
- 18) Forgetting that the REPEAT ... UNTIL statement will always be executed at least once.
- 19) Forgetting to convert binary numbers to decimal when the decimal value is to be output by OPMESS.
- 20) Forgetting the OPWAIT instruction.
- 21) Forgetting to convert decimal numbers input by the operator to binary before trying to use them.
- 22) Trying to use plus or minus signs with DECBIN.
- 23) Trying to input a number greater than 32767 from the operator console.
- 24) Forgetting the semicolon after a procedure name.
- 25) Using illogical jumps.
- 26) Trying to use values assigned within a procedure before the procedure has been activated.
- 27) Forgetting to define a procedure's constants and variables in the Constant and Variable sections.



**Part Three**  
**I/O Commands**



## 3.1 Overview

In this section we shall examine the most important MUSIL commands, the I/O commands. I/O commands deal with the physical transfer of data from core to a peripheral device, or from a peripheral device to core. Two kinds of operations are involved in this sort of data transfer.

Control operations do not result in any direct transfer of data, but they are necessary for data transfer. Typical control operations are, for example, the opening or closing of a file, the positioning of a magnetic tape, etc.

Transput operations call for the actual data output or input. Such operations are, thus, of two types: input or output mode. It is the output data that is the purpose of the entire data processing operation.

Both control and transput operations are performed in conformity with messages sent to the appropriate driver process. As noted in Part One, the driver always reports on the success or failure of an I/O operation, that is, it reports on its status.

The status of an I/O operation is put into a status word which is accessible by the programmer as well as by the system itself. The status word tells which aspects of the I/O operation have been successful and which have failed. If a failure is such that data processing can not proceed without some special exception handling procedure, then one of two things can happen.

If the programmer has not provided his MUSIL program with an applicable exception handling procedure, then the system will stop the processing of the job and display an error message on the operator device to inform the operator of what has occurred. Such error messages have the form

device name ERROR error number

For example

LPT ERROR 21

which means that processing can not continue because the line printer is off-line.

After the operator has corrected the situation, the job must be restarted from the beginning.

If the programmer has provided his program with an applicable exception handling procedure, then control will pass to it. Such procedures are called GIVEUP procedures and they are called when the status word is compared with a programmer-generated GIVEUP mask that tells the system which exception situations will be handled by the program instead of by the system.

The status of an I/O operation is only one of the components of the messages that pass between an I/O device driver and the program and system via the monitor. Other information that is passed along concerns the identification of the I/O device and data and constant reporting on where the data is at any time. This information is accessible to the programmer at any time via the file descriptors, which provide names for the locations from which the information can be accessed by the programmer.

I/O operations in MUSIL are arranged in such a way that the programmer is able to avoid all housekeeping tasks associated with I/O, while still retaining the option to assume control of some of these tasks if he so desires. This in fact is the reason why MUSIL was created, it having been felt that no existing programming language fully satisfied this objective in a convenient and logical way. For example, MUSIL provides for an automatic transfer of data between the peripheral device and the buffers, and it also provides the programmer with instructions by which he may assume control of this function. The instructions that give the programmer such direct control are called primitive operations.

Primitive operations are those operations that high-level instructions use to perform their functions. In MUSIL the primitive operations dealing with buffer control are available also to the programmer, though in normal situations he would have no need for them.

## **3.2 The Organization of Data**

I/O data is organized into groups called bytes, records, blocks, and files. This organization represents a hierarchy of data organization that makes it possible to deal with as much data at a time as any given I/O situation allows.



### 3.2.1 Bytes

For the RC 3600 bytes are groups of 8 bits. They correspond to "characters" which are defined as letters of the alphabet, numerals, or special symbols, for example, punctuation. Most work with MUSIL is done using bytes composed of bit patterns from the ASCII code, but for data any code may be used in working with an RC 3600.

The RC 3600 is a 16-bit per word machine; thus, each word in the machine has room for two bytes. MUSIL contains commands that allow the transport of data byte by byte.

### 3.2.2 Records

Records are groups of bytes. The concept of a record is a logical idea. There is no way to define a "record" in a way that covers all possible record types, but a record is usually considered to be the smallest piece of information that is of interest to the end-user.

Records are components of larger sets of information, called "files". Records within a file can be of several types.

Fixed length records are the components of a file that consists of records which all have the same number of bytes in them. Variable length records are components of a file in which the length of the records varies from one record to another.

### 3.2.3 Blocks

Some data media allow records to be blocked, that is, grouped together. Blocking is a physical concept. A block of data is a quantity that is read into memory or written out of memory in one physical operation. Blocking involves some physical delineation on some medium, such as an interblock gap of a magnetic tape or a control character input to a line printer within the data that it receives from memory. On the slower peripheral devices one can observe the occurrences of blocks, for when the end of a block is reached one can often see the device pause for a moment. On some media blocks can also be directly observed. A punched card, for example, is usually treated as a one-record block.

Both fixed and variable length records may be grouped together in blocks. The size of blocks is usually determined by the programmer, whose decision depends on the parameters of the application he is programming for.

Record length format may also be classified as undefined. This means that each record will be treated as a separate block. This is done, for example, when the block size of a magnetic tape input file is unknown to the programmer who must process the tape.

The programmer's decision on blocking strategy is related to his decision on the number of "buffers" that he will employ. Buffers are sections of core that are reserved to hold input or output data for transport from or to peripheral devices. They will be more fully discussed below, but here it should be noted that programmer decisions with respect to blocking, buffer size and number of buffers are very crucial for the speed of an I/O operation. For example, the common situation in which a magnetic tape containing print line images is to be printed out on the RC 3632 line printer operating at 1800 lpm requires 7 buffers for the print lines and one buffer for the magnetic tape input for optimal throughput. Such information can be calculated by the programmer, or it can be provided to him by Regnecentralen on the basis of experience.

3.2.3.1 Blocking methods. There may be any number of records in a block, including fractions of records, but the most normal cases are

- 1) To have one record per block, and
- 2) To have an integral number of records per block.

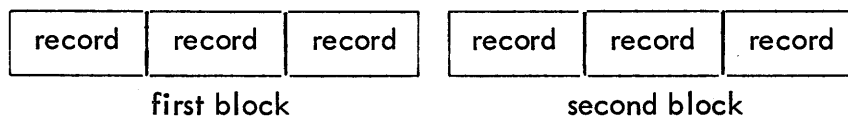
In the first case we say that the records are "unblocked". In the second case we say that the records are "blocked". In many cases the last block of a blocked file of records may not be completely filled with information.

Six types of records are common:

- 1) Fixed length unblocked, where all the records have the same length and there is one record per block.



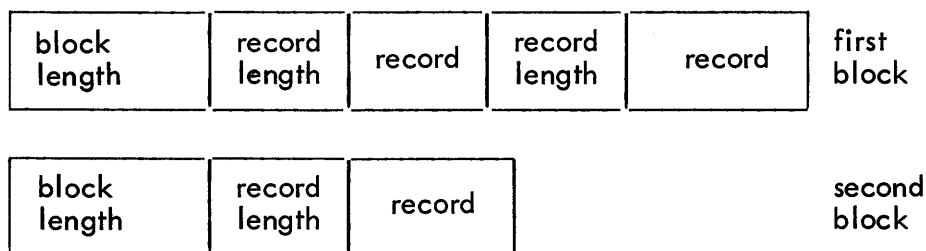
- 2) Fixed length blocked, where all the records have the same length and are grouped in blocks.



- 3) Variable length unblocked, where record length varies and there is one record per block.

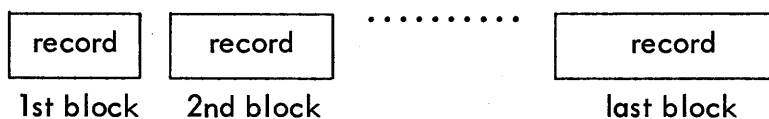


- 4) Variable length blocked, where record length varies and records are grouped in blocks.

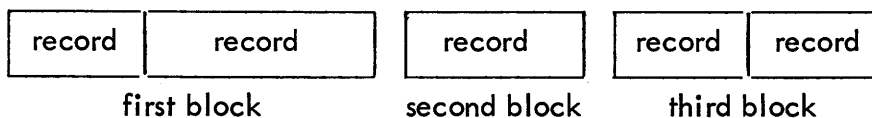


There may be any number of blocks and of records in a block.

- 5) Undefined length unblocked, where there is one record per block but there is no information about their length.



- 6) Undefined length blocked, where there is no information about the record length but records are grouped in blocks.



### 3.2.4 Files

In MUSIL we define a "file" as a set of data stored on some device. For this reason we can use the device name as the file name, if the medium has no catalog. If there is a catalog on the medium, as is usually the case for a disc, for example, then we cannot use the device name directly as the file name.

This usage of the concept of a file allows for simpler programming for non-cataloged media, the common situation on systems without disc support.

## 3.3 File Descriptors

We have stated that each file in a MUSIL program is defined by a file descriptor. The file descriptor gives the structure of the file and the nature of the data in it. The file descriptor also contains current information on the condition of the data in the file that are undergoing processing. The file descriptor allows both the system and the programmer to know and control what is going on during I/O operations.

Three kinds of information are found in a file descriptor:

- 1) file identification
- 2) control information
- 3) buffer information

File identification includes

- 1) The name of the file.

As explained above, uncataloged media are considered to contain but one file. Thus, in this case the name of the file is the name of the device on which the file appears. For example, the file on the paper tape reader is called PTR.

- 2) The kind gives information about the sort of device the file is on. The kinds include

Is the medium blocked? Is the device positionable? Is automatic error recovery wanted? For example, for magnetic tape it is usual to repeat a read or write operation that was not successful the first time, but one might specify a non-repeatable kind for a magnetic

tape file in certain situations to save read time, or to validity-check tapes.

- 3) The mode gives information about the data in the file. The mode contains an "operation code" that says if the data is input or output data.
- 4) Whether conversion is to be done, and if so, a reference to the conversion table.

For character-oriented devices (such as paper tape and card equipment and line printers), conversion is most commonly done during I/O, so that the conversion will most frequently be handled by the driver process. For block-oriented devices, conversion must be done in the MUSIL program.

#### Control information

- 1) The position of the file. For magnetic tapes this would be the "file number" in IBM usage, that is, the number of tape marks that have been passed. The position also includes the block number when relevant.
- 2) The status tells whether or not an I/O operation has been successful, and if not, why not.

#### Buffer information includes

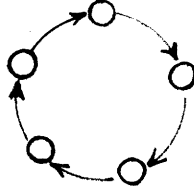
- 1) Buffer size. How big they are in bytes.
- 2) Used buffer. Which buffer is currently in use and the size of the block in it (in bytes).

From the above it can be seen that the file descriptor contains two types of information: permanent information identifying the file and its nature, and information on the current state of the file.

The information in the file descriptor can be accessed by the programmer through a system-defined record called FILEDESC. The use of this record will be described below.

## 3.4 Buffer Strategy

The MUSIL programmer is responsible for defining the number and size of the buffers to be used. The buffers so defined will then constitute a cyclical buffer pool.



At any time during I/O processing one of the buffers will be actually trans-putting data. This is the used buffer. If something goes wrong during data transfer, in the absence of a programmer-defined GIVEUP procedure within the MUSIL program the system will stop the job and display a device error message to the operator, as explained above. We shall call this the standard procedure in the description below.

Buffer strategy proceeds as follows:

### 3.4.1 Input

When a file is OPENed for data input to memory the system responds to the OPEN command by making one of the buffers the used buffer, by establishing a pointer from the file descriptor to this buffer. Block length in the file descriptor is set to zero also. The following operations then occur:

- 1) Input of the first block of the file is started into the used buffer.
- 2) Input of the second block is started into the second buffer and it is made the used buffer.
- 3) The process is continued until there are no more free buffers.
- 4) The first buffer is checked to see if data transfer to it was successfully completed.
- 5) If not, the standard procedure or a GIVEUP procedure is followed, if so, then the first buffer becomes the used buffer again, and data can be processed in it, after which it is used for more input.
- 6) This process is followed then for the second buffer, for the third, and so on, until all the data in the file is input.

The buffer pool can be viewed as a continually cycling wheel of buffers. This allows processing to proceed while data is being input. The optimal input speed is achieved when the number and size of the buffers are such as to allow continuous input at the maximum speed that the device allows.

It should be noted that if an input file is CLOSEd during the input process, then the buffer wheel will keep turning to empty the buffers, but the data remaining at the time of CLOSing will be lost.

### 3.4.2 Output

When a file has been OPENed to receive output, the buffer pool is similarly activated.

- 1) The first block of output is put into the first buffer, which is the used buffer.
- 2) The rest of the buffers are filled with output, each becoming the used buffer in its turn.
- 3) The data transfer from the first buffer to the device is checked for success.
- 4) If unsuccessful, then the standard procedure or a GIVEUP procedure is followed, and if successful, then the first buffer is ready to receive more data from core and becomes the used buffer.
- 5) The process is repeated for all the buffers and the "wheel" turns until the external file is complete.
- 6) The file can now be CLOSEd. E.g., on magnetic tape a double tape mark is written.

## 3.5 Exception Handling

There are two kinds of "exceptions": errors, such as parity errors, and normal stopping points for operations, such as reaching the end of a tape. Exceptions are reported to the system via the status word in the file descriptor.

After the attempted execution of a driver process operation the success or failure of the operation is reported in the status word. If all the bits of the status word are zero, then the operation has encountered no exceptions and processing can proceed. Though the specific events represented by the bits

of the status word are different for different devices, the general over-all representation is as follows:

<u>Bit</u>	<u>Interpretation</u>	<u>Action</u>
0	disconnected	hard error
1	off-line	hard error
2	device busy	if device kind is repeatable, the operation will be repeated; if not, then there is a hard error
3	device mode 1	ignored
4	device mode 2	ignored
5	device mode 3	ignored
6	illegal	hard error
7	End of File	hard error
8	block error	as for bit 2
9	data late	as for bit 2
10	parity error	as for bit 2
11	end of medium	hard error
12	position error	hard error
13	rejected	hard error
14	timer error	hard error
15	(not processed)	ignored

The standard system response to an exception is to try to repeat the operation. It does this up to five times, after which the error is a hard error. Such errors cause the program to stop running and an error message to be displayed on the operator console. The error message consists of the device name and an error number. This number is equal to 20 plus the number of the bit involved. For example, a line printer that is disconnected at the time the system wants to print with it will give rise to the message

LPT ERROR 20

for the name of a line printer is LPT and the status bit involved is number 0, and 20 plus 0 equals 20.

Consult the RC 3600 Data Conversion Operator's Reference Card for the various device error numbers and their interpretations.



The programmer may, however, elect to write a procedure that bypasses some or all of these standard exception-handling facilities of the driver involved. He does this by "turning off" those bits whose actions he wishes to bypass by specifying a

### 3.5.1 GIVEUP Mask

The GIVEUP mask is formed by creating a constant word that has 1's in those bit positions that correspond to the bits of the status word that the programmer wishes to prevent from initiating the standard system response. For example, if the programmer wishes to write his own exception handling routine for the occurrence of a position error, then his GIVEUP mask must contain a 1 in its bit 12.

GIVEUP masks are usually written in binary for easy readability, for example,

```
2'1100001111111111
```

GIVEUP masks are part of the corresponding file descriptor and are associated with

### 3.5.2 GIVEUP Procedures

Such procedures provide the routine that the system is to follow for each exception condition that the programmer has signified his desire to control by the 1's in the GIVEUP mask.

The programmer need not, however, provide for the standard repetition of the operation, as the GIVEUP procedure will be consulted by the system only after it has tried up to five repetitions of the operation that gave rise to the error.

## **3.6 Record and File Variables**

Record and file variables are defined within the Variable Section. We shall now describe how to do this.

3.6.1 Record Variables

Locations may be reserved for assignment for records. The general format is

```
name: RECORD
      record structure
      END;
```

For example, if we write

```
VAR
      TWOPARTS: RECORD
                HEAD:STRING(4);
                TAIL:STRING(4)
                END;
```

then we have reserved space in location TWOPARTS for the assignment of records of 8 bytes, with the first four bytes being given the name HEAD and the last four bytes being given the name TAIL. Subsequent to this definition HEAD and TAIL can be considered to be defined with respect to TWOPARTS, but not defined in themselves. Thus, later on in the Main Program section, HEAD and TAIL are defined only when named together with TWOPARTS, thus

```
TWOPARTS.HEAD      and
TWOPARTS.TAIL
```

The program will not become confused if the same sub-names are used for parts of other records. That is, given

```
TWOPARTS.HEAD      and
INPUTLINE.HEAD
```

the program will know that parts of different records are being referred to.

Further refinements in record definition are also possible. Consider the following:

```
VAR
      S:  STRING(1);
      INREC:  RECORD
                CCW:S;
                TEST:S;
                LINE:STRING(132);
                STOPF:STRING(2) FROM 1
                END;
```

This coding defines S as a one-character string location. It defines INREC as a record with the first two bytes as text strings of one byte each. The next 132 characters are called, collectively, LINE and are also text.

Furthermore, STOPF is defined as the name for the first two characters of the record INREC, these two characters being taken together. Thus, later on in the program we can make assignments to either of the first two characters of INREC individually, using their names CCW and TEXT, or we can assign values to them together, using their common name STOPF, viz., via

```
INREC.STOPF
```

Observe that in general each definition within a record definition ends with a semicolon, except that no semicolon is used after RECORD or before END.

As a final example, suppose that instead of giving the first two characters of INREC a name STOPF, we wanted to give a name to the first ten characters of LINE. Then, instead of writing

```
STOPF:STRING(2) FROM 1
```

we would write

```
NEWNAME:STRING(10) FROM 3
```

Note, then, that the characters of a record definition are numbered up from 1. Note particularly that a record may not be longer than 256 bytes.

### 3.6.2 File Variables

File variables are also defined in the Variable section. The most general format for a file variable definition is

```
name:   FILE
        'device',
        kind,
        number of buffers,
        buffer size,
        type;
        CONV conversion table name
        OF RECORD
            record structure
        END;
```

The meaning of the parts of the file variable definition are

- 1) Device is the name for the device the file is, or will be, on. Allowable code names are

MT0, MT1, MT2, MT3	Magnetic tape units on the first magnetic tape channel
MT4, MT5, MT6, MT7	Magnetic tape units on the second magnetic tape channel
LP0, LP1	Line printers
SP1, SP2	Serial printers
CDR, RDP	Card reader, Card reader punch
PTR, PTP	Paper tape reader, Paper tape punch
FD0, FD1	Flexible disc drive
DKP0, DKP1	Disc cartridge drive
PLT	Incremental plotter
CT0, CT1	Cassette tape unit
CP0	Charaband printer

The device name must be enclosed in quotation marks.

- 2) Kind gives information about the device. Allowable kinds are derived from the binary representation of the kind. The binary representations are

bit 15	character-oriented
14	block-oriented
13	positionable
12	repeatable (automatic error recovery)
11	cataloged medium (i.e., disc)

Thus, for a block-oriented, positionable, and repeatable device (such as magnetic tape), the kind word would be

2'0000000000001110

which in decimal is 14, so that kind here would be 14.

- 3) Number of buffers is selected by the programmer.
- 4) Buffer size is the size of a single buffer. This is also determined by the programmer.

5) Type refers to the record format. Allowable types are

UB	undefined, blocked	
U	undefined	
F	fixed	
FB	fixed, blocked	
V	variable	(IBM V format)
VB	variable, blocked	(IBM VB format)

6) GIVEUP procedure name need not occur. It is present if the programmer wishes to define his own device error routines. The procedure name must, of course, refer to a user-defined procedure in the Procedure section.

7) GIVEUP masks occur together with GIVEUP procedures. Inserting a 1 bit in any position in the word will cancel the corresponding action in the system's standard error routine.

8) Record structure can be given in terms of the name of a previously-defined record type (cf. 4.1), or can be specified directly, in the same way that records were defined, thus:

```

TYPE
    PLINE    = RECORD P:STRING(50) END;
VAR
    IN       : FILE
              'MT0', 14, 1, 60, FB
              OF PLINE;

```

or

```

VAR
    IN       : FILE
              'MT0', 14, 1, 600, FB
              OF STRING(50);

```

3.6.3 Example of File Definitions

```

0378
0379
0380
0381 IN:      FILE          ! INPUT FILE DESCRIPTION
0382          'MT0',        ! NAME OF INPUT DRIVER
0383          14,           ! KIND= REPEATABLE,
0384                          ! POSITIONABLE,
0385                          ! BLOCKED.
0386          2,           ! BUFFERS
0387          1995,        ! BUFFERSIZE
0388          FB;          ! FIXED BLOCKED
0389
0390          GIVEUP
0391          MTINERROR,    ! MT ERROR PROCEDURE
0392          2'011000111111110 ! GIVE UP MASK
0393          OF RECORD    ! RECORD STRUCTURE
0394          CCW:         STRING(1);
0395          DATA:      STRING(132)
0396          END;
0397
0398
0399
0400 OUT:     FILE          ! OUTPUT FILE DESCRIPTION
0401          'LPT',        ! NAME OF OUTPUT DRIVER
0402          2,           ! KIND= BLOCKED
0403          5,           ! BUFFERS
0404          133,        ! BUFFERSIZE
0405          U;          ! UNDEFINED
0406
0407          GIVEUP
0408          LPERROR,     ! LP ERROR PROCEDURE
0409          2'1100001011110110; ! GIVE UP MASK
0410
0411          CONV
0412          LPTTAB       ! CONVERSION TABLE
0413
0414          OF RECORD    ! RECORD STRUCTURE
0415          CCW:         STRING(1);
0416          DATA:      STRING(132)
0417          END;

```

## 3.7 Using the File Descriptor

The file descriptor contains a great deal of information about its corresponding file. The programmer may from time to time wish to access this information.

### 3.7.1 Accessing the File Descriptor

The most common use of the file descriptor is in accessing the status word for use in a GIVEUP procedure, but it is also common to use the file descriptor for displaying information about the file. To enable the programmer to have easy access to the file descriptor, there is a system-defined record for each file that is defined as follows:

```

FILEDESC = RECORD
  ZNAME   : STRING(6);      file name
  ZMODE   : INTEGER;       operation mode
  ZKIND   : INTEGER;       file kind
  ZMASK   : INTEGER;       GIVEUP mask
  ZFILE   : INTEGER;       file position
  ZBLOCK  : INTEGER;       block position
  ZCONV   : INTEGER;       conversion table address
  ZFORM   : INTEGER;       record format
  ZREM    : INTEGER;       number of bytes remaining in the
                           current block
  ZLENGTH : INTEGER;       record length
  ZFIRST  : INTEGER;       address of the first byte of the
                           current record
  ZTOP    : INTEGER;       address of the top byte of the
                           current record (that is, the first
                           byte of the next record)
  Z0      : INTEGER;       the status word
  ZUSED   : INTEGER;       address of the used buffer
  ZSHAREL : INTEGER;       block length for the buffer
END;
```

To access the current value of any field of the file descriptor, one writes

```
filename.field
```

For example,

```
IN.Z0
```

signifies the status word of the file called IN. To access this status word, then, we must have a command something like

```
STATIN:=IN.Z0;
```

which will put the status word into the previously-defined integer variables STATIN, or we could write directly something like

```
IF IN.Z0 = 2'1100000000000000 THEN
```

```
.....
```

To display the current block number, for example, of the current file, one could write

```
BINDEC(IN.ZBLOCK,OUTPUT);
OPMESS(OUTPUT);
```

where OUTPUT had been previously defined as a string variable.

### 3.7.2 Examining the Status Word

The status word can be examined by writing

```
BINDEC(filename.Z0,string variable name);
OPMESS(string variable name);
```

but since this event occurs frequently, a special command is available for displaying the contents of the status word on the operator device.

```
OPSTATUS(filename.Z0,string name);
```

In order to use this command, we must in the Constant Section define a string constant that will be capable of outputting a text for each bit of the status word that is non-zero. This is best illustrated by an example.

Let us define a constant called ERRORS, thus:



CONST

```

ERRORS = "DISCONNECTED <10><0>
          OFF-LINE      <10><0>
          BUSY          <10><0>
          BYTE OR NOISE <10><0>
          HARDWARE      <10><0>
          WRITE RING    <10><0>
          UNIT RESERVED <10><0>
          EOF           <10><0>
          BLOCK SIZE    <10><0>
          OVERRUN       <10><0>
          PARITY        <10><0>
          EOT           <10><0>
          POSITION        <10><0>
          DRIVER         <10><0>
          DENSITY        <10> ";

```

which is the interpretation of the error messages for a magnetic tape unit arranged so that each diagnostic will be printed on a separate line, that is <10> is a line feed and <0> is a carriage return.

Then, the OPSTATUS command

```
OPSTATUS(filename.Z0,ERRORS);
```

will display on the operator device the lines of ERRORS that correspond to positions of the status word that are non-zero. That is, if the status word contains at the time of inquiry

```
1000000000000000
```

then

```
DISCONNECTED
```

will be displayed, and a line feed and carriage return will be accomplished.

3.7.3 Example of a GIVEUP Procedure

```

0543                                     ! RC36-00099 PAGE 11 !
0544
0545
0546 PROCEDURE MTINERROR;
0547 BEGIN
0548     IF IN.Z0 AND 256 <> 0 THEN GOTO 9; ! EOF !
0549     IF IN.Z0 AND 8'041000 = 0 THEN BLOCKNO:= IN.ZBLOCK;
0550     IF IN.Z0 SHIFT 1 < 0 THEN OPMESS(MTMOUNTTAPE);
0551     IF IN.Z0 SHIFT 1 >= 0 THEN
0552         BEGIN
0553             OPMESS(MTTXT);
0554             MASK:=IN.Z0;
0555             SHOWERROR;
0556         END;
0557     REPEAT OPSTOP UNTIL STOPPED <> 0;
0558     IF STOPPED = 1 THEN GOTO 1;
0559     OPMESS(RUNTXT);
0560 END;
0561
0562 PROCEDURE LPERROR;
0563 BEGIN
0564     NEXTLP:= OUT.Z0 AND 8'000020;
0565     OUT.Z0:=OUT.Z0 - NEXTLP;
0566     IF OUT.Z0 SHIFT 1 < 0 THEN
0567         OUT.Z0:=OUT.Z0 AND 8'041342;
0568     IF OUT.Z0 = 8'040000 THEN IF NEXTLP <> 0 THEN
0569         OUT.Z0:= NEXTLP;
0570     IF OUT.Z0 AND 8'001342 <> 0 THEN
0571         OUT.Z0:= OUT.Z0 AND 8'001342;
0572     IF OUT.Z0<>0 THEN
0573         BEGIN
0574             OPMESS(LP TXT);
0575             MASK:=OUT.Z0;
0576             SHOWERROR;
0577             NEXTLP:=0;
0578             REPEAT OPSTOP UNTIL STOPPED <> 0;
0579             IF STOPPED = 1 THEN GOTO 1;
0580             OPMESS(RUNTXT);
0581             IF OUT.Z0 AND 8'141362 <> 0 THEN
0582                 REPEATSHARE(OUT);
0583         END;
0584 END;
0585

```

Recommended GIVEUP procedures for many situations are available from Regnecentralen.

### 3.8 Accessing File Contents

Once a file has been defined in the Variable Section and opened in the Main Program Section, its records can be read into memory or output to a medium.

When a record has been read into memory, it may be accessed for assignment or comparison with respect to the data it contains. Such operations, however, can take place only on previously-defined (that is, named) records or parts of records. The format for accessing the data in a file is

filename↑

for accessing the data as a whole, and

filename↑.fieldname

for parts of a record.

For example,

```

VAR
    ALPHA:STRING(2);
    IN: FILE
        "MT0", 14, 1, 1340, FB
        OF RECORD
            CCW      :STRING(1);
            SELECT1  :STRING(1) FROM 1;
            DATA    :STRING(2) FROM 2
        END;
BEGIN
    .....
    .....
    ALPHA:=IN↑.DATA;
    .....
    .....
END;
```

puts the current contents of DATA from file IN into ALPHA. Open and get record statements are defined below.

## 3.9 I/O Commands

Though buffer strategy in MUSIL is based on the transfer of blocks of data, actual processing is almost always performed on records or characters. There are two levels of I/O commands in MUSIL, therefore: the higher-level record and character commands and the primitive procedures that the higher-level commands utilize.

We shall now explain the I/O commands in detail.

### 3.9.1 Opening and Closing Files

OPEN(filename,mode);

This command ensures that file identification is established, reserves the peripheral device involved, prepares for conversion if necessary, and initializes the file. The "mode" is input or output and the conditions of transport, for example, odd or even parity. It is represented by a decimal number that varies from device to device. Operation mode numbers can be found in the Appendix.

CLOSE(filename,release);

closes the file. For input files the closing process concludes all pending data transfers, but it does not check the data transfer, since this data will be lost anyway. For output files the closing process completes and checks all pending data transfers and writes a terminator to the file, for example, a file mark in the case of magnetic tape. "Release" may be any integer. If it is zero, then the device will not be made available to any other program and the device will not be automatically set off-line. If "release" is not zero, then the device will be set off-line. For example, in the case of magnetic tape, release not equal to zero will cause the tape to be rewound and the tape unit to be set off-line.

### 3.9.2 Record-by-Record Data Transfer

This is the most common transport means in MUSIL. It can be done after a file has been OPENed.

GETREC(filename, variable name);

This command has the general effect of making "the next" record available for processing. The first time it is used after an OPEN command on the file, it must start the transfer of data into the buffer wheel and establish control over the turning of the wheel. It does this by calling the primitive commands INBLOCK, TRANSFER, and WAITTRANSFER (explained below). When used subsequently, it simply makes the next record in the block the "current" record, so that this record can be processed. In this way GETREC can be used to step through the records in the buffers until it again becomes necessary to input a block of data. The specific actions of GETREC in various situations are best described by an example.

Let us examine the following situation:

```

VAR
    SIZE:    INTEGER;
    INFILE:  FILE
            ..... !FILE DEFINITION!
            END;

BEGIN
    OPEN(INFILE, mode);
    .....
    GETREC(INFILE, SIZE);
    .....
    .....
END;
```

The GETREC command would then give rise to the following:

For file record format undefined and unblocked, U:

The general effect is that the contents of a buffer become available for processing. If all the buffers are empty, then the instruction starts the transfer of data into the buffer wheel as well. Let us say we are using cards and have only one buffer, then the effect of the instruction is to read a card. For paper tape and one buffer, the effect is to read and process as much of the tape as will fit into one buffer. The size of the block read is put into SIZE.

For file record format undefined and blocked, UB:

The general effect is to make a number of characters available for processing, this number being equal to the current value of SIZE. If the buffers are empty at the time of call, then SIZE characters must be input. If there is less than SIZE characters in the buffer, then the next block is input until SIZE characters can be available for processing. This situation can be taken advantage of for reading the first character of a tape block to see what kind of tape block one is dealing with. For example, if SIZE and MTO have been previously defined, and we write

```

OPEN(MTO, mode);
SIZE:=1;
GETREC(MTO, SIZE);
IF BYTE MTO = some binary code THEN
    BEGIN
        SIZE:=MTO.ZREM;
        GETREC(MTO, SIZE);
        .....
        .....          !PROCESSING OF BLOCK!
    END;
```

then the effect is to read and examine the first byte of the tape and then decide if we want to read in the rest of the first block.

For file record format fixed length and unblocked, F:

In this case the record format has been defined in the Variable Section, as has its length. The command will make a record available, reading in a record, if necessary. The number of bytes made available will be put into SIZE. This will be the record length in this case. In case a record of the correct length cannot be gotten, the exception procedure is called with status bit 8 (block length). If the error is accepted, the record is delivered as a short block.

For file record format fixed length and blocked, FB:

The instruction will make a record available. If it cannot find one in the buffers, then it will first begin to read the next blocks into the buffer wheel. The record length will appear in SIZE. Incorrect record lengths are handled as for F formats.

For file record format variable length and unblocked, V:

The next record will be made available. If the buffers are empty, input will proceed into the buffer wheel. The first four bytes of the record, which contain the record length, are decoded and put into SIZE. The variable length format used is the IBM V format. Incorrect record lengths are handled as for F formats.

For file record format variable length and blocked, VB:

This is the means of handling the IBM VB format. The next record is made available, by reading the next blocks into the buffer wheel, if necessary. The first four bytes of a new block, containing the block length, are decoded. In any case the first four bytes of the record, containing the record length, are decoded and placed in SIZE.

PUTREC(filename, name or number or expression);

is the command that makes space for a record available for output and starts the buffer wheel turning for output by calling the primitive procedures OUTBLOCK, TRANSFER, and WAITTRANSFER. The specific actions are

For file record format undefined and unblocked, U:

PUTREC(FILENAME, SIZE);

A previous buffer is output, then space is reserved in the buffer for the next record to be output the next time PUTREC is called. The name, number, or expression is the length of the record.

For file record format undefined and blocked, UB:

The command makes space for the next record of SIZE bytes in the buffer. If there is no room for it, it outputs a buffer and reserves space in the new buffer. The name, number or expression in this case is the length of the record.

For file record format fixed and unblocked, F:

The number, name, or expression is ignored. The record length is that given in the record definition. The effect of the command is to output a previous record and make space for the current record in a new buffer.

For file record format fixed and blocked, FB:

If the record can fit into a buffer, it is put there; if not, a block will be output to make room for it. The name, number or expression in the command is ignored. The SIZE is that of the record definition.

For file record format variable and unblocked, V:

This command utilizes the IBM V format. A block will be output to make room for a new record. The block size and record size (which are equal) are computed, so that the output medium will be in IBM V format. The record of length equal to the name, number or expression will have room in the space reserved.

For file record format variable and blocked, VB:

If there is no room in the buffer, a block is output to make room for the current record. In any case the record descriptor is computed and the block descriptor is up-dated. The record size must be in the name, number, or expression.

Some general notes on GETREC and PUTREC should be made. First, the second factor of the command has different general functions in the two expressions. For GETREC, it relates to the length of the record coming in from the buffer, a given fact. In PUTREC on the other hand, the record length must be given to the program so that it knows how much output space to look for in the output buffer wheel (except for F formats).

Secondly, GETREC and PUTREC are not equivalent to READ and WRITE, respectively. They do not, in fact, move data.

### 3.9.3 Character-by-Character Data Transfer

Data transfer by character can be done on files that have records which are undefined and unblocked. The general effect of character I/O commands is to input or output data character by character.

INCHAR(filename, integer variable name);

takes the next available byte from the input buffers and places it in "integer



variable name". If the buffers are all empty, then INCHAR will call INBLOCK, TRANSFER, and WAITTRANSFER, in order to read data from the input device.

OUTCHAR(filename, value);

checks to see if there is room for one byte in the output buffers. If not, then it calls OUTBLOCK, TRANSFER, and WAITTRANSFER, in order to make room for the byte. The byte that is put into the buffer is the last byte of the "value", which may be a number, the current contents of a variable, or the value of an expression.

OUTTEXT(filename, string name);

outputs the string contained in the "string name". Output continues until a binary zero is reached. Thus, the string to be output must contain such a binary zero. This command is useful, for example, for putting text onto print-out. OUTTEXT is a shorthand command offered in MUSIL as a convenience. Its effect can be obtained by using combinations of MOVE and OUTCHAR.

#### 3.9.4 Primitive Procedures

The normal MUSIL I/O commands are arranged in a hierarchy in which the commands higher up in the hierarchy operate by calling on the commands lower down:

highest level:	GETREC	PUTREC	INCHAR	OUTCHAR
intermediate level:	INBLOCK		OUTBLOCK	
lowest level:	TRANSFER		WAITTRANSFER	

The intermediate and lowest level commands are "primitive procedures" when they are used directly and explicitly by the programmer. It should be noted that few situations will arise in which it will be necessary for the programmer to use them explicitly, while the higher-level commands always use them.

TRANSFER(filename, length, operation mode);

is the command that actually starts the transfer of physical data to and from the peripheral devices. In the input mode TRANSFER causes the used buffer to be filled from the data medium and moves the used buffer pointer to the next buffer in the buffer wheel. In output mode TRANSFER writes the contents of the used buffer to the medium and moves the used buffer pointer to the next buffer in the output buffer wheel. The number of bytes transput is in "length", which must be an integer expression. It should be noted that TRANSFER does not look to see if the used buffer is ready to receive or output new data. Therefore, TRANSFER is always followed by

WAITTRANSFER(filename);

which looks to see if the used buffer is ready with data input or finished with output. If not, it delays the transput until the used buffer is ready.

INBLOCK(filename);

administrates the data input operation. Its effect is to call TRANSFER until all the buffers are started in the buffer wheel. Then it calls WAITTRANSFER to prepare the first input.

OUTBLOCK(filename);

administrates the data output operation. Its effect is to call TRANSFER to start the output buffers. Then it calls WAITTRANSFER in order to ready the next buffer.

REPEATSHARE(filename);

can be used only within a GIVEUP procedure. This command restarts a rejected operation and then returns to the internal waittransfer, so that the operation can be completed before returning control to the next command in the program.

## 3.10 Data Manipulation

In this section we shall explore some commands that are used for the direct manipulation of data. These commands are understandable both as simple MUSIL commands and as I/O handling commands. They facilitate the solution of data conversion problems.

```
WAITZONE(filename);
```

is mentioned here because it is a lower-order command that is used by the following command, as well as being useful in general to the programmer. Its effect is to halt processing in an orderly way, so that processing can later be resumed without trouble. For input files it empties the buffer wheel, skipping data. For output it assures that all output operations have been completed. For example,

```
IF operator action is called for THEN
    BEGIN
        WAITZONE(filename);
        .....           !OPERATOR ACTION!
        IF ready now THEN
            .....           !RESUME PROCESSING!
    END;
```

```
SETPOSITION(filename, file number, block number);
```

allows one to position a positionable medium, such as a magnetic tape. SETPOSITION calls WAITZONE first, in order to halt processing. Then it positions the medium to the desired block. For example,

```
SETPOSITION(MT0,3,8);
```

positions the tape to the eighth block of the third file on the tape on tape unit 0. (Here we use the word "file" in its physical sense.)

The command

```
MOVE(string name, from n+1st byte, to string name, from n+1st
      byte, for number of bytes)
```

that can be used for operator communication can also be used for data manipulation, for example,

```
MOVE(IN|,1,OUT|,0,LENGTH);
```

will take the current record of file IN, starting with the second byte, and move it into the current record from file OUT, starting with the first byte, until LENGTH number of bytes have been moved. Note that if LENGTH is too big, there will be no error message.

MOVE cannot be used to move bytes around within a single string.

```
CONVERT(string name, string name, table name, length);
```

can be used to convert between media. It is best explained by an example.

Consider the expression

```
CONVERT(MT0|,OUT|,TABLE1,OUT.ZLENGTH);
```

This command would take the current record of file MT0 and convert it according to the table defined as TABLE1, and put the result into the current record of file OUT, doing this for as many bytes of the first record as is represented by OUT.ZLENGTH, which is precisely the length of the current record of file OUT.

"Length" can be an expression, a number, or a variable name.

```
INSERT(byte name, record name, place);
```

This instruction can insert a byte into a place in a string. For example,

```
INSERT(VALUE,OUT|,2);
```

places the 8 least significant bits of VALUE into the third byte position of the current record of file OUT.

### 3.11 Possible I/O Errors

It is impossible to list all the things that might go wrong where I/O commands occur. In the list that follows, we have tried to list the most obvious places where the programmer might take extra care.

- 1) In setting up a GIVEUP mask remember to start counting bits from zero and be aware of what each bit signifies for the specific device you are using.
- 2) Try to provide for restart capacity in all operational situations where operator error might occur, such as his forgetting to put a line printer on line.
- 3) Check carefully for correct punctuation in record and file definitions.
- 4) Make sure you have up-to-date information on device names and kinds.
- 5) For difficult decisions about buffer size and number, you may consult Regnecentralen for advice.
- 6) For difficult GIVEUP procedures, particularly in communications programs, you may consult Regnecentralen.
- 7) Be particularly careful in using the GETREC and PUTREC commands. They are a common source of error.
- 8) Be particularly careful when using the primitive procedures. They are very sensitive.
- 9) Remember that MOVE cannot be used to move bytes within a string.

### **3.12 Example of a MUSIL Program**



## MUSIL COMPILER/2

0000 !  
 0001  
 0002  
 0003  
 0004  
 0005  
 0006  
 0007  
 0008  
 0009  
 0010  
 0011  
 0012  
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 0057  
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 0059  
 0060  
 0061

RCSL: 43-GL140

AUTHOR: CT

EDITED: 74.08.12

PROGRAM RC36-00007.00

MUS CARDS TO TAPE

KEYWORDS: MUSIL,CONVERSION,CDR,MTA,LISTING

ABSTRACT: THIS PROGRAM HANDLES 80-COLUMN CARDS IN EBCDIC CODE AND GENERATES FIXED OR VARIABLE LENGTH FORMAT RECORDS WHICH MAY BE WRITTEN WITH A SPECIFIED NUMBER OF RECORDS IN EACH OUTPUT BLOCK. OUTPUT IS EBCDIC CODE IN BLOCKS OF UP TO 2000 BYTES ON NO LABEL TAPE WITH OR WITHOUT BLOCK AND RECORD LENGTH FIELDS DUE TO THE RECORD TYPE SPECIFICATION. THE PROGRAM MAY BE OPERATED FROM EITHER OCP OR TTY.

RCSL 43-GL141: ASCII SOURCE TAPE !

0062 1  
0063  
0064 TITLE: MUS CARDS TO TAPE.  
0065  
0066 ABSTRACT: THIS PROGRAM HANDLES 80-COLUMN CARDS IN EBCDIC CODE  
0067 AND GENERATES FIXED OR VARIABLE LENGTH FORMAT RE-  
0068 CORDS WHICH MAY BE WRITTEN WITH A SPECIFIED NUMBER  
0069 OF RECORDS IN EACH OUTPUT BLOCK.  
0070 OUTPUT IS EBCDIC CODE IN BLOCKS OF UP TO 2000 BYTES  
0071 ON NO LABEL TAPE WITH OR WITHOUT BLOCK AND RECORD  
0072 LENGTH FIELDS DUE TO THE RECORD TYPE SPECIFICATION.  
0073 THE PROGRAM MAY BE OPERATED FROM EITHER OCP OR TTY.  
0074  
0075 SIZE: 5564 BYTES. INCLUDING TWO 80 BYTES INPUT BUFFER  
0076 AND ONE 2000 BYTES OUTPUT BUFFER.  
0077  
0078 DATE: AUGUST 12TH 1974.  
0079  
0080 RUNTIME PARAMETERS:  
0081 BLOCK NO : 00001 NEXT BLOCK TO BE WRITTEN TO CURRENT FILE.  
0082 FILE NO : 00001 THE FILE IN WHICH THE BLOCK IS WRITTEN.  
0083 REWIND : + INDICATES IF REWIND OF TAPE AT END OF INPUT.  
0084 FIXRECS : + INDICATES OUTPUT RECORD FORMAT FIXED/VARIABLE.  
0085 NOTE: THIS PARAMETER CAN ONLY BE CHANGED BEFORE  
0086 THE PROGRAM IS STARTED AND AFTER END OF JOB.  
0087 MAXCOL : 00080 MAXIMUM NUMBER OF COLUMNS TRANSFERRED WHEN  
0088 VARIABLE LENGTH FORMAT OUTPUT OR NUMBER OF  
0089 COLUMNS WHEN FIXED LENGTH FORMAT OUTPUT.  
0090 MINCOL : 00080 MINIMUM NUMBER OF COLUMNS TRANSFERRED WHEN  
0091 VARIABLE LENGTH FORMAT OUTPUT.  
0092 BLOCKED : 00025 MAXIMUM NUMBER OF RECORDS IN EACH BLOCK.  
0093 OTHER OUTPUT MESSAGES:  
0094 CONTSTATE: +/- STATE OF CONTINUE SWITCH (TTY ONLY).  
0095 PROG NO : 7 PROGRAM EXECUTION IS STOPPED.  
0096 RUNNING PROGRAM EXECUTION IS STARTED.  
0097 SUSPENDED DRIVERS RELEASED, PROGRAM EXECUTION IS STOPPED.  
0098 LOAD CARD DECK CARD READER HOPPER EMPTY AND CONTINUE IS ON.  
0099 CR ERROR NNNNN CONSULT THE RC3600 OPERATORS MANUAL.  
0100 MT ERROR NNNNN CONSULT THE RC3600 OPERATORS MANUAL.  
0101 END JOB PROGRAM EXECUTION IS TERMINATED.  
0102  
0103 INPUT MESSAGES:  
0104 STOP STOPS EXECUTION WRITING PROG NO : 7.  
0105 SUSPEND STOPS EXECUTION RELEASING DRIVERS (TTY ONLY).  
0106 INT NEXT PARAMETER IS DISPLAYED  
0107 (ESCAPE BUTTON ON TTY HAS SAME EFFECT).  
0108 STATE ALL PARAMETERS ARE DISPLAYED (TTY ONLY).  
0109 "VALUE" CURRENTLY DISPLAYED PARAMETER IS CHANGED  
0110 TO "VALUE".  
0111 "TEXT"="VALUE" THE PARAMETER IDENTIFIED BY "TEXT" IS  
0112 CHANGED TO "VALUE"  
0113 CONT STATE OF CONTINUE SWITCH IS INVERTED.  
0114 START PROGRAM EXECUTION IS STARTED.  
0115 NOTE: AFTER CR ERROR START MEANS ACCEPTING  
0116 THE ERRONEOUS INPUT, AFTER MT ERROR START  
0117 MEANS REPEATING THE WRITE OPERATION.  
0118 SPECIAL REQUIREMENTS: NONE  
0119





```

0165
0166 CRTABLE=          ! CR CONTROLLER FORMAT TO ERCDIC
0167          0      1      2      3      4      5      6      7 !
0168 #
0169 !   0 !   64 241 242 243 244 245 246 247
0170 !   8 !  249 49  50  51  52  53  54  55
0171 !  16 !  248 121 122 123 124 125 126 127
0172 !  24 !   56 57  58  59  60  61  62  63
0173 !  32 !  240 97 226 227 228 229 230 231
0174 !  40 !  233 33  34  35  36  37  38  39
0175 !  48 !  232 105 224 107 108 109 110 111
0176 !  56 !   40 41  42  43  44  45  46  47
0177 !  64 !   96 209 210 211 212 213 214 215
0178 !  72 !  217 17  18  19  20  21  22  23
0179 !  80 !  216 89  90  91  92  93  94  95
0180 !  88 !   24 25  26  27  28  29  30  31
0181 !  96 !  208 161 162 163 164 165 166 167
0182 ! 104 !  169 225  98  99 100 101 102 103
0183 ! 112 !  168 160 170 171 172 173 174 175
0184 ! 120 !  104 32 234 235 236 237 238 239
0185 ! 128 !   80 193 194 195 196 197 198 199
0186 ! 136 !  201  1  2  3  4  5  6  7
0187 ! 144 !  200 73  74  75  76  77  78  79
0188 ! 152 !   8  9  10  11  12  13  14  15
0189 ! 160 !  192 129 130 131 132 133 134 135
0190 ! 168 !  137 65  66  67  68  69  70  71
0191 ! 176 !  136 128 138 139 140 141 142 143
0192 ! 184 !   72  0 202 203 204 205 206 207
0193 ! 192 !  106 145 146 147 148 149 150 151
0194 ! 200 !  153 81  82  83  84  85  86  87
0195 ! 208 !  152 144 154 155 156 157 158 159
0196 ! 216 !   88  16 218 219 220 221 222 223
0197 ! 224 !  112 177 178 179 180 181 182 183
0198 ! 232 !  185 113 114 115 116 117 118 119
0199 ! 240 !  184 176 186 187 188 189 190 191
0200 ! 248 !  120  48 250 251 252 253 254 255
0201 #;
0202

```

```
0203
0204 VAR
0205
0206 OPDUMMY:          STRING(2);          ! RUNTIME PARAMETERS !
0207 PROGNO:          INTEGER;
0208 BLOCKNO:         INTEGER;
0209 FILENO:          INTEGER;
0210 REWIND:           INTEGER;
0211 FIXRECS:          INTEGER;
0212 MAXCOL:           INTEGER;
0213 MINCOL:           INTEGER;
0214 BLOCKED:          INTEGER;
0215
0216 OPTEXT:           STRING(20);         ! COMMUNICATION AREA !
0217 OPSTRING:         STRING(20);
0218 OPDEC:            STRING(10);
0219
0220 OPCONT:           STRING(2);          ! INTERNAL VARIABLES !
0221 NEXTCONT:         STRING(1);
0222 GLCONT:           STRING(1);
0223 WBLOCKED:         INTEGER;
0224 ERRORNO:         INTEGER;
0225 MASK:             INTEGER;
0226 TOM:              INTEGER;
0227 SIGN:             INTEGER;
0228 Q:                INTEGER;
0229 PAR:              INTEGER;
0230 LENGTH:           INTEGER;
0231 OPENED:           INTEGER;
0232 P1:               INTEGER;
0233 P2:               INTEGER;
0234 P3:               INTEGER;
0235 S1:               STRING(2);
0236 S2:               STRING(2);
0237 NEXTMT:           INTEGER;
0238 INLENGTH:         INTEGER;
0239 OUTLENGTH:        INTEGER;
0240 CARDSREAD:        INTEGER;
0241 SAVEDSUSPEND:     INTEGER;
0242
```

```

0243
0244
0245 IN:   FILE           ! INPUT FILE DESCRIPTION !
0246      'CDR',         ! NAME OF INPUT DRIVER !
0247      2,             ! KIND= BLOCKED !
0248      2,             ! BUFFERS !
0249      80,            ! SHARESIZE !
0250      U;             ! UNDEFINED !
0251
0252      CONV
0253      CRTABLE;       ! CONVERSION TABLE !
0254
0255      GIVEUP
0256      CRERROR,       ! CR ERROR PROCEDURE !
0257      2'0110001011110110 ! GIVE UP MASK !
0258
0259      OF STRING(80); ! RECORD STRUCTURE !
0260
0261
0262
0263 OUT:  FILE           ! OUTPUT FILE DESCRIPTION !
0264      'MTO',         ! NAME OF OUTPUT DRIVER !
0265      14,            ! KIND= REPEATABLE, !
0266                        !     POSITIONABLE, !
0267                        !     BLOCKED. !
0268      1,             ! BUFFERS !
0269      2000,          ! SHARESIZE !
0270      FB;           ! FIXED(BLOCKED)/VARIABLE(BLOCKED) !
0271
0272      GIVEUP
0273      MTOUTERROR,    ! MT ERROR PROCEDURE !
0274      2'0110011111011011 ! GIVE UP MASK !
0275
0276      OF STRING(80); ! RECORD STRUCTURE !
0277

```

```
0278
0279 PROCEDURE INITPOSITION;
0280 BEGIN
0281     IF FIXRECS= -1 THEN
0282     BEGIN
0283         IF IN.ZMODE=33 THEN FIXRECS:=-2;
0284         IF IN.ZMODE=0 THEN OPEN(IN,21);
0285     END;
0286     IF FIXRECS= -2 THEN
0287     BEGIN
0288         IF IN.ZMODE=21 THEN FIXRECS:=-1;
0289         IF IN.ZMODE=0 THEN OPEN(IN,33);
0290     END;
0291     IF FIXRECS=-1 THEN OUT.ZFORM:=3;
0292     IF FIXRECS=-2 THEN OUT.ZFORM:=5;
0293     IF OUT.ZMODE=0 THEN OPEN(OUT,3);
0294     IF BLOCKNO=OUT.ZBLOCK THEN
0295     IF FILENO=OUT.ZFILE THEN
0296     GOTO 999;
0297     WAITZONE(IN);
0298     SETPOSITION(OUT,FILENO,BLOCKNO);
0299 999:
0300 END;
0301
0302 PROCEDURE CONTINUE;
0303 BEGIN
0304     GLCONT:=OPCONT;
0305     OPCONT:=NEXTCONT;
0306     NEXTCONT:=GLCONT;
0307     OPMESS(OPCONT);
0308 END;
0309
```

```

0310
0311
0312 PROCEDURE DIRECTUPDATE;
0313 BEGIN
0314     P1:=0;           ! INDEX IN INPUT STRING !
0315     P2:=0;           ! INDEX IN CONSTANT STRING !
0316     P3:=1;          ! PARAMETER NUMBER IN CONSTANT STRING !
0317     REPEAT BEGIN
0318         MOVE(OPTTEXT,P1,S1,0,1);
0319         MOVE(OPTXTS,P2,S2,0,1);
0320         WHILE BYTE S1 <> BYTE S2 DO
0321             BEGIN
0322                 IF BYTE S2 = 0 THEN P3:=P3+1;
0323                 P2:=P2+1;
0324                 MOVE(OPTXTS,P2,S2,0,1);
0325                 IF P3>NOQ THEN S2:=S1;
0326             END;
0327             IF P3<=NOQ THEN
0328                 BEGIN
0329                     WHILE BYTE S1 = BYTE S2 DO
0330                         BEGIN
0331                             P1:=P1+1;
0332                             P2:=P2+1;
0333                             MOVE(OPTTEXT,P1,S1,0,1);
0334                             MOVE(OPTXTS,P2,S2,0,1);
0335                             IF BYTE S1 = 61 THEN
0336                                 BEGIN
0337                                     MOVE(OPTTEXT,P1+1,OPTTEXT,0,10);
0338                                     LENGTH:=LENGTH-P1-1;
0339                                     Q:=P3;
0340                                     MOVE(OPDUMMY,Q*2,OPDUMMY,0,2);
0341                                     PAR:= WORD OPDUMMY;
0342                                     P3:=NOQ;
0343                                 END;
0344                             END;
0345                             P2:=P2-P1+1;
0346                             P1:=0;
0347                         END;
0348                     END UNTIL P3>=NOQ;
0349                 END;
0350

```

```

0351
0352 PROCEDURE OPCDM;
0353 BEGIN
0354 1000: Q:=0;
0355 1010: REPEAT BEGIN
0356 IF OPTTEXT=STATE THEN
0357 BEGIN Q:=1; OPMESS(CONTSTATE); IF OPCONT=FIVE THEN
0358 OPMESS(PLUS); IF OPCONT=FIFTEEN THEN
0359 OPMESS(MINUS); GOTO 1040;
0360 END;
0361 1015: Q:=Q+1;
0362 1020: OPSTATUS(1 SHIFT(16-Q),OPTXTS); IF Q<>1 THEN BEGIN
0363 MOVE(OPDUMMY,Q*2,OPDUMMY,0,2);
0364 PAR:= WORD OPDUMMY;
0365 IF PAR = -1 THEN OPMESS(PLUS);
0366 IF PAR = -2 THEN OPMESS(MINUS);
0367 IF PAR >= 0 THEN
0368 BEGIN BINDEC(PAR,OPDEC); OPMESS(OPDEC); END; END;
0369 IF OPTTEXT=STATE THEN GOTO 1060;
0370 1040: OPMESS(ENDLINE);
0371 OPWAIT(LENGTH);
0372 OPTTEXT:=OPSTRING;
0373 OPIN(OPSTRING);
0374 IF OPTTEXT=STATE THEN BEGIN Q:=0; GOTO 1015; END;
0375 IF OPTTEXT = SUSPEND THEN
0376 BEGIN
0377 SAVEDSUSPEND:= 1;
0378 GOTO 1040;
0379 END;
0380 IF LENGTH > 6 THEN DIRECTUPDATE;
0381 IF LENGTH > 6 THEN GOTO 1020;
0382 IF OPTTEXT = START THEN GOTO 1070;
0383 IF OPTTEXT = STOP THEN GOTO 1000;
0384 IF OPTTEXT = CONT THEN
0385 BEGIN CONTINUE; GOTO 1040; END;
0386 IF OPTTEXT = INT THEN GOTO 1060;
0387 IF OPTTEXT = NEXTPARAM THEN GOTO 1060;
0388 IF OPTTEXT = NL THEN GOTO 1020;
0389 IF OPTTEXT = ENDLINE THEN GOTO 1020;
0390 IF OPTTEXT = RETURN THEN GOTO 1020;
0391 SIGN:=0;
0392 IF OPTTEXT = MINUS THEN SIGN:=-1;
0393 IF OPTTEXT = PLUS THEN SIGN:=+1;
0394 IF SIGN <> 0 THEN INSERT(48,OPTTEXT,0);
0395 DECBIN(OPTTEXT,TOM);
0396 IF PAR < 0 THEN
0397 BEGIN IF SIGN=0 THEN GOTO 1020; PAR:=-2;
0398 IF SIGN=1 THEN PAR:=-1; GOTO 1050;
0399 END;
0400 IF SIGN=0 THEN
0401 BEGIN SIGN:=1; PAR:=0; END;
0402 PAR:=PAR+TOM*SIGN;
0403 IF PAR<0 THEN GOTO 1020;
0404 1050: INSERT(PAR SHIFT(-8),OPDUMMY,0);
0405 INSERT(PAR, OPDUMMY,1);
0406 MOVE(OPDUMMY,0,OPDUMMY,Q*2,2);
0407 IF OPTTEST <> 0 THEN GOTO 1040;
0408 GOTO 1020;
0409 1060: IF OPTTEXT=STATE THEN IF Q<NOQ THEN GOTO 1015;
0410 END UNTIL Q>=NOQ; GOTO 1000;
0411 1070: OPMESS(RUNTXT);
0412 END;
0413

```

```

0414
0415 PROCEDURE OPSTOP;
0416 BEGIN
0417     OPWAIT(LENGTH);
0418     OPTEXT:=OPSTRING;
0419     OPIN(OPSTRING);
0420     IF OPTEXT=CONT THEN CONTINUE;
0421     IF OPTEXT=STOP THEN GOTO 1;
0422     IF OPTEXT = SUSPEND THEN SAVEDSUSPEND:= 1;
0423 END;
0424
0425 PROCEDURE SHOWERROR;
0426 BEGIN
0427     ERRORNO:=20;
0428     WHILE MASK>0 DO
0429     BEGIN
0430         MASK:=MASK SHIFT 1;
0431         ERRORNO:=ERRORNO+1
0432     END;
0433     BINDEC(ERRORNO,OPTEXT);
0434     OPMESS(OPTEXT); OPMESS(ENDLINE);
0435 END;
0436
0437 PROCEDURE CRERROR;
0438 BEGIN
0439     IF IN.Z0 AND 2'10000 <> 0 THEN
0440     BEGIN
0441         SAVEDSUSPEND:= -SAVEDSUSPEND;
0442         GOTO 9;
0443     END;
0444     IF CARDSREAD=0 THEN GOTO 9;
0445     IF IN.Z0<>0 THEN
0446     BEGIN
0447         OPMESS(CRTXT);
0448         MASK:=IN.Z0;
0449         SHOWERROR;
0450     END;
0451     REPEAT OPSTOP UNTIL OPTEXT=START;
0452     OPMESS(RUNTXT);
0453     IF IN.Z0 AND 8'040000 <>0 THEN !CARD READER OFFLINE!
0454     REPEATSHARE(IN);
0455 END;
0456
0457 PROCEDURE MTOUTERROR;
0458 BEGIN
0459     IF OUT.Z0 AND 8'043000 = 0 THEN BLOCKNO:=OUT.ZBLOCK;
0460     NEXTMT:= OUT.Z0 AND 8'000020;
0461     OUT.Z0:=OUT.Z0 - NEXTMT;
0462     IF OUT.Z0 SHIFT 1 < 0 THEN OPMESS(MTMOUNTTAPE);
0463     IF OUT.Z0 SHIFT 1 >=0 THEN
0464     IF OUT.Z0 <> 0 THEN
0465     BEGIN
0466         OPMESS(MTTXT);
0467         MASK:=OUT.Z0;
0468         SHOWERROR;
0469     END;
0470     IF OUT.Z0<>0 THEN
0471     BEGIN
0472         REPEAT OPSTOP UNTIL OPTEXT=START;
0473         OPMESS(RUNTXT);
0474         IF OUT.Z0 AND 8'063352 <> 0 THEN
0475         REPEATSHARE(OUT);
0476     END;
0477 END;
0478

```



```
0479
0480
0481 PROCEDURE CHANGETAPE;
0482 BEGIN
0483     OPMESS(ENDTAPE);
0484 99:   OPWAIT(LENGTH);
0485     OPTEXT:=OPSTRING;
0486     OPIN(OPSTRING);
0487     IF OPTEXT=CONT THEN
0488     BEGIN
0489         CONTINUE;
0490         GOTO 99;
0491     END;
0492     IF OPTEXT=STOP THEN
0493     BEGIN
0494         NEXTMT:=0;
0495         CLOSE(OUT,1);
0496         OPMESS(MTMOUNTTAPE);
0497         FILENO:=1;
0498         BLOCKNO:=1;
0499         REPEAT OPSTOP UNTIL OPTEXT=START;
0500         OPEN(OUT,3);
0501         SETPOSITION(OUT,1,1);
0502     END;
0503 END;
0504
```

```

0505
0506 BEGIN
0507 BLOCKNO:=1; FILENO:=1; REWIND:=-1; NEXTMT:=0;
0508 FIXRECS:=-1; MAXCOL:=80; MINCOL:=80; BLOKED:=25;
0509 OPCONT:=FIFTEEN; NEXTCONT:=FIVE; WBLOCKED:=0; CARDSREAD:=0;
0510 OUT.ZFILE:=1; OUT.ZBLOCK:=1; OPIN(OPSTRING); SAVEDSUSPEND:= 0;
0511
0512 1: OPCOM; INITPOSITION;
0513
0514 2: IF OPTTEST<>0 THEN OPSTOP;
0515
0516 3: GETREC(IN,INLENGTH);
0517 IF FIXRECS=-1 THEN
0518 BEGIN
0519     INLENGTH:=MAXCOL; OUTLENGTH:=MAXCOL;
0520     GOTO 4;
0521 END;
0522 IF INLENGTH<=MINCOL THEN OUTLENGTH:=MINCOL;
0523 IF INLENGTH>MINCOL THEN
0524 BEGIN
0525     OUTLENGTH:=INLENGTH;
0526     IF OUTLENGTH>MAXCOL THEN OUTLENGTH:=MAXCOL;
0527 END;
0528
0529 4: PUTREC(OUT,OUTLENGTH); CARDSREAD:=CARDSREAD+1;
0530 MOVE(IN|,0,OUT|,0,INLENGTH); IF INLENGTH<OUTLENGTH THEN
0531 MOVE(SPACES,0,OUT|,INLENGTH,OUTLENGTH-INLENGTH);
0532 IF BLOCKNO<>OUT.ZBLOCK THEN
0533 BEGIN
0534     BLOCKNO:=OUT.ZBLOCK;
0535     WBLOCKED:=0;
0536     IF NEXTMT<>0 THEN CHANGETAPE;
0537     GOTO 2;
0538 END;
0539 WBLOCKED:=WBLOCKED+1;
0540 IF WBLOCKED<BLOKED THEN GOTO 3;
0541 OUTBLOCK(OUT);
0542 BLOCKNO:=OUT.ZBLOCK;
0543 WBLOCKED:=0;
0544 IF NEXTMT<>0 THEN CHANGETAPE;
0545 GOTO 2;
0546
0547 9: IF SAVEDSUSPEND = -1 THEN
0548 BEGIN
0549     OUTBLOCK(OUT);
0550     BLOCKNO:= OUT.ZBLOCK;
0551     CLOSE(IN,1);
0552     CLOSE(OUT,1);
0553     CARDSREAD:= 0;
0554     WBLOCKED:= 0;
0555     SAVEDSUSPEND:= 0;
0556     OPMESS(SUSTXT);
0557     GOTO 12;
0558 END;
0559 IF CARDSREAD=0 THEN GOTO 10;
0560 IF OPCONT = FIVE THEN GOTO 10;
0561 GOTO 11;
0562 10: OPMESS(CRMOUNTDECK);
0563 GOTO 12;
0564 11: CLOSE(IN,1);
0565     CLOSE(OUT,REWIND+2);
0566     BLOCKNO:=1; FILENO:=FILENO+1;
0567     IF REWIND=-1 THEN FILENO:=1;
0568     IF FILENO=1 THEN NEXTMT:=0; CARDSREAD:=0; WBLOCKED:=0;
0569     OPMESS(EOJTXT);
0570 12: REPEAT OPSTOP UNTIL OPTTEXT=START;
0571     INITPOSITION; OPMESS(RUNTXT); GOTO 2;
0572 END;

```

**Part Four**  
**Appendix**



## 4.1 The Type Section

A MUSIL program may have a Type Section between the Constant Section and the Variable Section. The purpose of the Type Section is to provide a shorthand for defining types, or categories, of variables. It is most important to remember that the definitions in the Type Section are not substitutes for definitions in the Variable Section. Type Section definitions are used to define a structure of a variable type. For example, if one has two or more file definitions in the Variable Section, and both definitions have the same structure, then one can save some time and effort by defining the structure as a type within the Type Section and referring to this type in the Variable Section.

### 4.1.1 Section Structure

The Type Section begins with the word

```
TYPE
```

not followed by any punctuation. It ends with a semicolon and its statements are separated from one another by semicolons.

### 4.1.2 Integer and String Types

The format for integer type definition is

```
name = INTEGER;
```

and for a string variable type, it is

```
name = STRING(n);
```

where n is the string length in bytes, for example,

```
I = INTEGER;
```

```
LINE = STRING(132);
```

Once such definitions have been set up in the Type Section, we may use them to define variables in the Variable Section, thus:

```
VAR
```

```
    L,M,N:I;
```

```
    D,E:LINE;
```

which defines the variables L, M, and N as INTEGER and D and E as STRING(132).

As can easily be seen, the Type Section is not terribly useful in defining integer or string variables. It is much more useful, however, in defining file and record variables.

#### 4.1.3 Record and File Types

Record and file types are defined in the Type Section in a way very reminiscent of the definition of records and files in the Variable Section. The only difference is that the colon used to assign a name to the record or file is replaced by an equals sign, for example,

```

TYPE
    PLINE = RECORD
        L1:STRING(20);
        L2:STRING(50)
    END;
    IN    = FILE
        'MT0', 14, 1, 600, FB
    OF PLINE;

```

Once these definitions have been made of a record and a file structure, then in the Variable Section we can write, for example,

```

VAR
    IN1, IN2:IN;
    LINE1, LINE2, LINE3:PLINE;

```

which defines the file variables IN1 and IN2 as having the same structure as file type IN, and the record variables LINE1, LINE2, and LINE3 as having the structure of type PLINE.

#### 4.1.4 Possible Errors

In general one may make the same errors in the Type Section as one can make in the Variable Section. One can also forget to use the equals sign properly. But the most frequent error is to forget to define the appropriate variables in the Variable Section and to try to use type definitions as variable definitions.

## 4.2 Reference List of MUSIL Commands, Operators, and Symbols

### 4.2.1 MUSIL Commands

COMMAND	DESCRIPTION
BINDEC(binary value name, decimal value name)	Convert binary to decimal
CLOSE(filename, release)	Close file
CONVERT(stringname, stringname, tablename, length)	Code conversion
DECBIN(decimal value name, binary value name)	Convert decimal to binary
GETREC(filename, variable name)	Get next record for processing
GOTO label	Unconditional branching
IF relation THEN statement	Conditional branching
IF relation THEN statement ELSE statement	Conditional branching
INBLOCK(filename)	Prepare a block for input
INCHAR(filename, integer variable name)	Get next character for processing
INSERT(integer value or name, string variable name, integer value or name)	Insert byte value in string at place designated
MOVE(string name, from n+1th byte, to string name, from n+1th byte, for number of bytes)	Move bytes from one string to another
OPEN(filename, mode)	Open file
OPIN(string variable name)	Input string from console
OPMESS(string variable name)	Output string to console
OPSTATUS(filename.Z0, string constant)	Display status word
OPTEST	Test for successful input from console
OPWAIT(integer variable name)	Wait for operator input
OUTBLOCK(filename)	Prepare a block for output
OUTCHAR(filename, constant)	Prepare a character for output
OUTTEXT(filename, stringname)	Insert text into output file
PUTREC(filename, integer value or name)	Prepare next record for output
REPEAT statement UNTIL relation	Repeat command
REPEATSHARE(filename)	Restart command for a GIVEUP procedure
SETPOSITION(filename, filenumber, block number)	Position medium
TRANSFER(filename, length, mode)	Output or input data
WAITTRANSFER(filename)	Wait for completion of output or input
WAITZONE(filename)	Pause
WHILE relation DO statement	Repeat command
\$COPY	Copy code from a second source and place it here
\$END	Stop copying from this source
PROCEDURE name (parameters)	Define a code procedure
CODEBODY external identification;	

4.2.2 MUSIL Operators and Symbols

! text !	Comment
#list#	Table of numbers
<number>	Byte value
"text"	Text string of ASCII characters or byte values
,	Constant section separator
:	Variable definition
=	Constant declaration, Type definition
:=	Value assignment
;	Statement separator
BYTE	Byte value
WORD	Word value
nnnnn:	Label
BEGIN ... END	Compound statement
( )	Parentheses
+	Addition
-	Subtraction
*	Multiplication
/	Division
<	Less than
>	Greater than
<=	Less than or equal to
>=	Greater than or equal to
=	Equal to
<>	Not equal to
m AND n	Logical AND
m SHIFT n	Shift m left n bytes
m SHIFT (-n)	Shift m right n bytes
m EXTRACT n	Extract n bits of m
U	Undefined and unblocked
UB	Undefined and blocked
F	Fixed and unblocked
FB	Fixed and blocked
V	Variable and unblocked (IBM V format)
VB	Variable and blocked (IBM VB format)



## 4.3 ASCII Code Table

Decimal Representation	7-Bit Octal Code	Character	Decimal Representation	7-Bit Octal Code	Character	Decimal Representation	7-Bit Octal Code	Character
0	000	NUL	43	053	+	86	126	V
1	001	SOH	44	054	,	87	127	W
2	002	STX	45	055	-	88	130	X
3	003	ETX	46	056	.	89	131	Y
4	004	EOT	47	057	/	90	132	Z
5	005	ENQ	48	060	0	91	133	***
6	006	ACK	49	061	1	92	134	
7	007	BEL	50	062	2	93	135	
8	010	BS	51	063	3	94	136	↑
9	011	HT	52	064	4	95	137	←
10	012	LF	53	065	5	96	140	·
11	013	VT	54	066	6	97	141	a
12	014	FF	55	067	7	98	142	b
13	015	CR	56	070	8	99	143	c
14	016	SO	57	071	9	100	144	d
** 15	017	SI	58	072	:	101	145	e
16	020	DLE	59	073	;	102	146	f
17	021	DC1	60	074	<	103	147	g
18	022	DC2	61	075	=	104	150	h
19	023	DC3	62	076	>	105	151	i
20	024	DC4	63	077	?	106	152	j
21	025	NAK	64	100	@	107	153	k
22	026	SYN	65	101	A	108	154	l
23	027	ETB	66	102	B	109	155	m
24	030	CAN	67	103	C	110	156	n
25	031	EM	68	104	D	111	157	o
26	032	SUB	68	105	E	112	160	p
27	033	ESC	70	106	F	113	161	q
28	034	FS	71	107	G	114	162	r
29	035	GS	72	110	H	115	163	s
30	036	RS	73	111	I	116	164	t
31	037	US	74	112	J	117	165	u
32	040	SP	75	113	K	118	166	v
33	041	!	76	114	L	119	167	w
34	042	"	77	115	M	129	170	x
35	043	#	78	116	N	121	171	y
36	044	\$	79	117	O	122	172	z
37	045	%	80	120	P	123	173	***
38	046	&	81	121	Q	124	174	
39	047	'	82	122	R	125	175	
40	050	(	83	123	S	126	176	~
41	051	)	84	124	T	127	177	DEL
42	052	*	85	125	U			

\*\* Special control characters.  
Will be interpreted in accordance with actual device specifications.

\*\*\* Reserved for national characters.

## 4.4 Device Reference Tables

Device error numbers are explained in the RC 3600 Operator's Manual and on the RC 3600 Data Conversion System Operator's Reference Card. Device information of use to the MUSIL programmer is as follows:

### 4.4.1 Kind Table

bit 15	set if device is character-oriented
14	set if full blocks should be transferred
13	set if positioning has any effect
12	set if an operation may be repeated
11	set if the device is a cataloged disc file

#### Examples:

14 =	1110	Magnetic tape station
1 =	0001	Line printer
3 =	0011	Line printer
2 =	0010	Card reader
1 =	0001	Teletype
1 =	0001	Paper tape punch
1 =	0001	Paper tape reader

### 4.4.2 Operation Mode Table

#### Paper tape reader driver - PTR

1	binary, the input character is delivered
5	odd parity, the most significant bit is removed
9	even parity, the most significant bit is removed

#### Paper tape punch driver - PTP

3	binary, the converted character is output
7	odd parity, the converted character is augmented by the complement of its parity in the most significant position
11	even parity

## Line printer driver - LPn

- 3 the converted characters are output
- 7 the first byte of output is interpreted as a carriage control word

## Magnetic tape driver - MTn

- 1 read packed, byte limit = 18
- 5 read packed, byte limit = 0
- 9 read unpacked, byte limit = 18
- 13 read unpacked, byte limit = 0
- 3 write

Concerning the Magnetic tape driver: when using 7 track tape, if 4096 is added to any of the operation modes, then the reading or writing will be done in even parity. If the number 8192 is added to any of the mode numbers, then the resulting number will cause reading or writing to be done in the tape's lower density.

For example, 4099 signifies write with even parity, for  $4099 = 3 + 4096$ .

## Card reader driver - CDR

- 1 read binary byte
- 5 read decimal punched cards
- 21 read decimal punched cards
- 33 read decimal punched cards and skip trailing blank columns (a minimum of ten columns are read)

## Card reader punch driver - RDP

- 1 read binary bytes
- 5 read decimal word
- 9 read binary word
- 21 read decimal bytes, skip trailing blanks (a minimum of ten columns is read)
- 11 punch decimal byte
- 19 print decimal byte
- 27 punch and print decimal byte
- 47 punch binary word

If 256 is added to any of the read mode numbers then the resultant sum used as an operation mode causes Hopper #2 to be selected. Adding 64 causes the card to be released and a new card to be fed.

Adding 256 to any of the punch mode numbers causes Stacker #2 to be selected. Adding 64 causes a card to be fed before the operation is performed. For example, 257 = 1 + 256 means Read binary bytes from a card in the second hopper.

Plotter driver - PLT

3 write byte

Flexible disc driver - FDn

1 read  
3 write  
5 read, non-skip  
7 write and check read

Charaband printer driver - CP0

3 print converted characters  
7 interpret first byte as carriage control  
15 output to load Direct Access Vertical Format Unit at 6 lpi  
31 output to the DAVF unit at 8 lpi

Cassette tape unit driver - CTn

CT0

1 read one block per ECMA 34 version 2  
3 write one block per ECMA 34 version 2  
+4 read data blocks continuously  
9 read one block per ECMA 34 version 1  
17 read one block without checking

CT4

1 read one block per ECMA 34 version 2  
3 write per ECMA 34 version 2  
5 continuous read per ECMA 34 version 2  
7 write with control read per ECMA 34 version 2  
9 read one block per ECMA 34 version 1  
11 write per ECMA 34 version 1  
13 continuous read per ECMA 34 version 1  
15 write with control read per ECMA 34 version 1  
17 read one block without check  
21 continuous read without check

#### Serial printer driver - SPT

- 3 write converted character
- 7 interpret first byte as carriage control

#### Cartridge disc driver - DKPn

- 1 read
- 3 write

#### Communications equipment

See terminal reference cards.

NB: The Operation Code, which represents whether the operation is an output or an input operation, consists of the two least significant bits of the Operation Mode.

## 4.5 Program Production

Once one has written a MUSIL program down on paper, one may enter it onto some medium, such as cards or tape, or one may key it directly into the machine. For the latter choice, it is convenient to use MUSIL Text Editor, which takes in the MUSIL source code as data to itself. Text Editor can also be used to change programs that did not compile properly or to up-date job programs.

Instructions for using MUSIL Text Editor can be found in the manual MUSIL Text Editor Version Two. A list of Text Editor commands and error messages can also be found on the RC 3600 MUSIL Programmer's Reference Card.

### 4.5.1 Compilation

Once the MUSIL source code has been written and debugged, it must be compiled into MUSIL object code for execution by the MUSIL Interpreter. On the RC 3600 MUSIL Programmer's Reference Card can be found the procedures and commands for compilation.

#### 4.5.2 MUSIL Compiler Error Messages

During compilation errors give rise to the following numbers on the listing or on the operator's console:

- 020202    Number overflow, a numeric constant exceeds 65535, or 16 bits.
- 020301    Illegal character in input.
- 030102    < appearing within a string is not followed by a numeric literal.
- 030202    The construct < number is not followed by a >.
- 030302    The number between < and > exceeds an 8-bit byte value.
- 030403    Core overflow, produced code exceeds available space.
- 030503    Core overflow, code contains too many relocation bits.
- 040105    Name conflict in Constant Section.
- 040205    Name conflict in Type Section.
- 040302    Syntax in Type Section, no = following an identifier.
- 040405    Name conflict in Variable Section.
- 040506    File variable with 0 buffers.
- 040602    Procedure head not followed by ,
- 040702    Procedure without legal identifier or with name conflict.
- 050102    Type is no identifier.
- 050202    ( is missing after string.
- 050302    Length undefined for string.
- 050402    String with length 255 declared.
- 050502    ) is missing after string.
- 050604    Undefined type identifier. Note that no forward declarations are allowed.
- 050702    Improper termination of type specification.
- 051002    Field of type different from string.
- 051102    Incorrect use of FROM.
- 051205    Name conflict in GIVEUP procedure.
- 051304    Conversion table undeclared.
- 051406    Conversion table type error.
- 060206    Double defined label.
- 060302    Variable is no identifier. Or undeclared.
- 060402    . is not followed by identifier or by undeclared field.
- 060504    Identifier undeclared.
- 060606    Type error with BYTE or WORD.
- 060702    Relational operator missing.
- 061002    Procedure statement with missing )

- 061102 Type error in procedure parameter.
- 061306 Illegal number of parameters.
- 061406 Type error with operator.
- 061506 Overflow of work registers. Expression too complex.

Error Messages which cause skipping of program parts:

- 000040 Syntax in section delimiter.
- 000041 Syntax in constant declaration.
- 000042 Syntax in table declaration.
- 000043 Type specification incorrectly terminated.
- 000044 Variable declaration incorrectly terminated.
- 000045
- 000046
- 000051 Syntax in field list.
- 000052 Syntax in file declaration.
- 000063 Incomprehensible statement.
- 000064 Incorrect label declaration.
- 000065 Incomprehensible expression.

#### 4.5.3 Copying MUSIL Program Parts

At the time of compilation parts of one MUSIL source may be copied into a place within another MUSIL source. The sources must, of course, not be located on the same device. To do the copying the command

\$COPY

is inserted in the program into which the code is to be copied at the location where the code copied is to be placed. The command

\$END

is inserted into the source whose code is being copied and is placed after the last instruction to be copied. Copying begins from the first statement of the copied program.

## 4.6 List of Reserved MUSIL Words

To produce a tape or card deck or disc with a full run, including systems software and one or more applications programs in MUSIL object code, one must use the program generator supplied as part of one's program production package.

### 4.6 List of reserved MUSIL words

AND	GOTO	RECORD
		REMOVEENTRY
BEGIN	IF	REPEAT
BINDEC	INBLOCK	REPEATSHARE
BYTE	INCHAR	
	INITCAT	SETPOSITION
CHANGEENTRY	INSERT	SHIFT
CLOSE	INTEGER	STRING
CODEBODY		
CONV	LOOKUPENTRY	THEN
CONVERT		TRANSFER
CONST	MOVE	TRANSLATE
CREATEENTRY		TYPE
	OF	
DECBIN	OPEN	UNTIL
DO	OPIN	
	OPMESS	VAR
ELSE	OPSTATUS	
END	OPTEST	WAITTRANSFER
EXTRACT	OPWAIT	WAITZONE
	OUTBLOCK	WHILE
FILE	OUTCHAR	WORD
FROM	OUTTEXT	
GETREC	PROCEDURE	
GIVEUP	PUTREC	



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RC SL 42 - i 0344

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