

# **COURSE OVERVIEW**

- **Day 2 - Subsystems; System Management**
  - **Subsystems**
    - \* **Overview**
    - \* **Definitions**
    - \* **Structure**
    - \* **View operations**
    - \* **Code views**
    - \* **Compatibility**
    - \* **Multi-machine management**
    - \* **Target builder**
    - \* **Reservations and check-in/out**
    - \* **Work orders**
  - **Users, Groups, Access Control, and Accounting**
    - \* **Users and groups**
    - \* **Access control rules**
    - \* **Job identity**
    - \* **Privileged mode**
    - \* **Universe Acls**
    - \* **Problems and diagnosis**

# **COURSE OVERVIEW**

- **Day 3 - Networking; Disks and Deamons**
  - **Networking**
  - **Disk daemon management**
  - **Garbage problems**
  - **Daemons**
  - **Error log**
  - **System availability**

# **COURSE OVERVIEW**

- **Day 4 - System Diagnosis and Recovery**
  - **System structure**
  - **DFS**
  - **Crash dumps**
  - **Kernel and EEDB**
  - **Diagnostic Tools**
  - **Disk errors**
  - **Questions and answers**

# OBJECTS - DEFINITION

- **Defined by R1000 architecture**
- **Low-level, execution-oriented**
  - **The term "object" is also used to refer to environment objects such as files and Ada units**
- **Consist of a "type part" and a "value part"**
- **"Types" are special kinds of objects with a null value part**
- **The "value part" is some number of bits of data**
- **The "type part" indicates the type and structure of the value. Part of the type part of an object is a type descriptor. The format of the type descriptor is defined by the R1000 architecture.**

# **R1000 GURU COURSE**

- **Objectives**
  - **Understand the R1000 system implementation**
  - **Be able to diagnose system problems**
    - \* **find solutions**
    - \* **know what information to send**
    - \* **know when to get help**
  - **Be more effective at giving customers advice and training their system experts**
  - **Be a more worldly person**

# **COURSE OVERVIEW**

- **Day 1 - Basic Structures and Ada**
  - **Objects**
  - **R1000 storage organization**
  - **Instruction set**
  - **Storage hierarchy**
  - **Execution environment**
  - **Exceptions, unchecked deallocation, unchecked conversion**
  - **Program structure**
  - **Tasking and task operation**
  - **Virtual processors and jobs**
  - **Object management system**

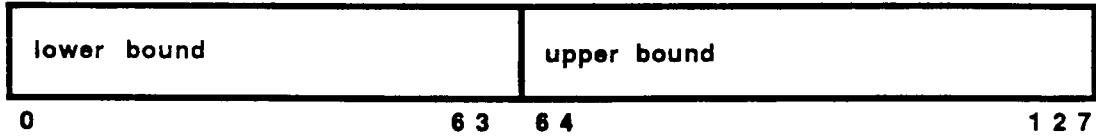
# OBJECTS - DEFINITION (CONT'D)

- Part of the type information is the object kind 7 bits 128 kinds  
~56 used
- The kind divides objects into classes that the R1000 knows about. For example:
  - Discrete (includes integer, fixed, enumerations, etc.)
  - Float
  - Access
  - Entry
  - Package/Task
  - Array
  - Record
  - Etc.
- The type identifies the specific Ada type and includes information about the data format
- Types follow Ada semantics with respect to type compatibility
- Type descriptors describe a type by storing all information relevant to interpreting the bits that represent the value

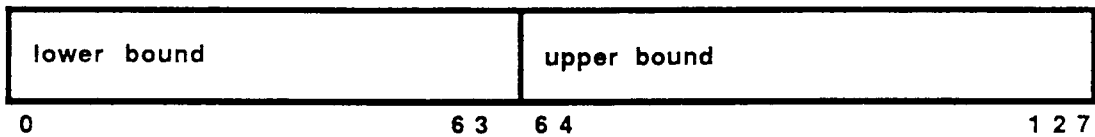
# OBJECT - TYPE DESCRIPTORS

- Examples of type descriptors

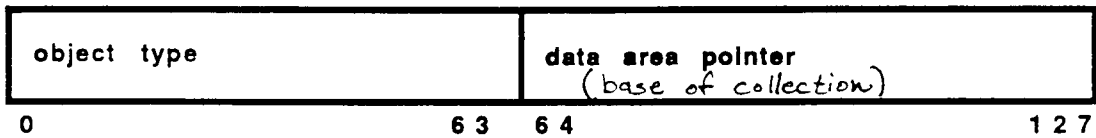
For discrete types:



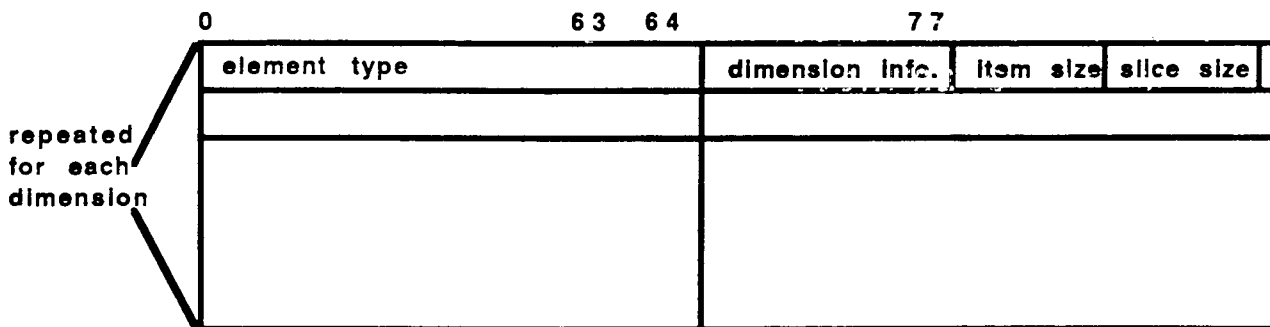
For floating types:



For access types:



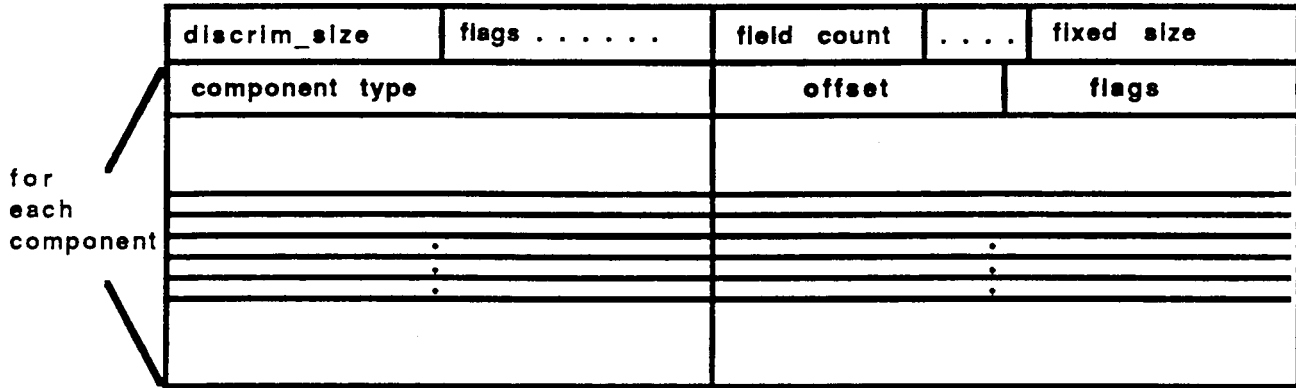
For array types:





# OBJECTS - TYPE DESCRIPTORS (CONT'D)

- For a record type:



- Each reference to a type is a pointer to one of these descriptors.
- Most Ada type declarations result in the creation of one (or more) of these type descriptors.
- Many operations on objects require that the descriptor be examined.

# OBJECTS - TYPE DESCRIPTORS (CONT'D)

- Sometimes some type information is associated with an individual object rather than in the type descriptor.
- Variant record objects are more complicated and have more information associated with them. They are slower to create and access.
- Variant objects can be constrained or unconstrained. An unconstrained object can assume any of its variants. A constrained object cannot - it can have only one specific variant.

# OBJECTS - ADA EXAMPLES

```
type Array_Type is array (integer range <>) of character;
```

```
type T (x: integer:=10; b:boolean:=false) is record
```

```
    A: Array_Type (1..x);
```

```
    case B is
```

```
        when true => F: integer
```

```
        when false => F2: Array_type (1..500)
```

```
    end case
```

```
end record;
```

```
V: T (1000,false);      -- big constrained object  
W: T (1,true);         -- small constrained object  
X: T;                  -- unconstrained, giant object
```

- Default initial values versus constraints

```
subtype T1 is T (100,true);      -- constraints  
V: T (1000,false);             -- constraints
```

# OBJECTS - REPRESENTATION OF VALUES

- For discrete and other "simple" objects, the representation is simple binary. Only the number of bits needed to represent the value are used. Values can be unsigned or two's complement (the type indicates which).

- Examples:

X: integer;	32 bits
b: boolean;	1 bit unsigned
type E is (A,B,C);	
V:E	2 bits unsigned
C: character;	8 bits - special case
l: long integer	64 bits

- Records and arrays are fully packed. (More details on this later)
- "Scalar" objects have simple, fixed representations
- Records with no discriminants are also simply represented by the values of each of the fields

# OBJECTS - REPRESENTATION OF VALUES (CONT'D)

- **Array and variant record objects may require additional descriptive information**
  - There is an indicator of whether the value is constrained (the "is\_constrained bit")
  - There is a "variant clause index" indicating which specific variant is active
  - There may be array bounds information for array objects ("bounds with object")
- **See handout #1 - Note the following:**
  - Objects of constrained <sup>array</sup>(sub)types do not have extra bits
  - Unconstrained values have descriptors that indicate their current configuration
  - Constants also have such descriptors

```

subtype Short_Int is Integer range 0 .. 20; 5+1sign = 6 bits
type T1 (D1 : Boolean := False; D2 : Short_Int := 20) is
  record
    F1 : Natural := 100;
    case D1 is
      when True =>
        F2 : Boolean := True;
      when False =>
        F4 : String (1 .. D2);
    end case;
  end record;

```

```
V1 : T1;
```

```

type T1 size      = 334 bits
Object V1 size   = 334 bits.
Object fits in 42 bytes.
0128000001900000 0100000003800000 0004000000500000 0000000000000000
0000000000000000 0002

```

```

subtype T2 is T1 (D1 => True, D2 => 3);
V2 : T2;

```

```

type T2 size      = 47 bits
Object V2 size   = 47 bits.
Object fits in 6 bytes.
80C7FFFFFFFFFE

```

```

subtype T3 is String (1 .. 5);
V3 : T3 := "abcde";

```

```

type String size = 4294967359
type T3 size     = 40

```

```

Object "abc" size = 88 bits.
Object fits in 11 bytes.
0000000100000003 616263

```

*64 bits of bound, 24 bits of data (for constants)*

```

Object V3 size   = 40 bits.
Object fits in 5 bytes.
6162636465

```

# OBJECTS - OTHER KINDS

- **Not all R1000 architectural objects are Ada data objects. Others include:**
  - **Subprograms**
  - **Packages and tasks**
  - **Exceptions**
  - **"Select" objects**
  - **Task entries**

**See handout #2 for a full list**

The Kind field is 7 bits.

0 Discrete\_Var,  
1 Control\_State,  
2 Subprogram\_Ref\_For\_Call,  
3 Word3\_Flag,  
4 Discrete\_Ref,  
5 Module\_Key,  
6 Subprogram\_For\_Call,  
7 Mark\_Word\_Flag,  
8 Float\_Var,  
A Variable\_Ref,  
C Float\_Ref,  
D Deletion\_Key,  
E Entry\_Var,  
10 Access\_Var,  
11 Static\_Connection,  
12 Subprogram\_Ref\_For\_Call\_Elaborated,  
14 Access\_Ref,  
15 Interface\_Key,  
16 Subprogram\_For\_Call\_Elaborated,  
18 Task\_Var,  
19 Dependence\_Link,  
1C Task\_Ref,  
1E Select\_Var,  
1F Auxiliary\_Mark,  
21 Micro\_State1,  
22 Subprogram\_Ref\_For\_Call\_Visible,  
24 Subvector\_Var,  
26 Subprogram\_For\_Call\_Visible,  
29 Micro\_State2,  
2C Subarray\_Var,  
2E Family\_Var,  
32 Subprogram\_Ref\_For\_Call\_Visible\_Elaborated,  
36 Subprogram\_For\_Call\_Visible\_Elaborated,  
38 Heap\_Access\_Var,  
3C Heap\_Access\_Ref,  
3E Default\_Var,  
3F Activation\_Link,  
41 Control\_Allocation,  
42 Accept\_Subprogram\_Ref,  
44 Record\_Var,  
46 Accept\_Subprogram,  
49 Scheduling\_Allocation,  
4C Variant\_Record\_Var,  
4E Delay\_Alternative,  
52 Interface\_Subprogram\_Ref,  
56 Interface\_Subprogram,  
58 Package\_Var,  
5F Accept\_Link,  
6C Vector\_Var,  
6E Family\_Alternative,  
74 Matrix\_Var,  
76 Null\_Subprogram,  
7C Array\_Var,  
7E Exception\_Var,  
7F Activation\_State



# OBJECTS - BASIC MACHINE PREDEFINED TYPES

<u>NAME</u>	<u>USE</u>
Discrete	scalar types (integer, fixed point, enumeration)
Float	floating point
Access	pointers (access types)
Package	Ada packages
Task	Ada tasks
Array	arrays
Record	records
Variant_Record	variant records
Select	internal state for select statement
Entry	task entry
Family	family of task entries
Exception	Ada exceptions
Subprogram	represent machine callable entities

- These are called "operand classes". If an instruction tries to operate on the wrong kind of object, the exception "operand\_class\_error" is raised.

# OBJECTS - TASK/PACKAGE OBJECTS

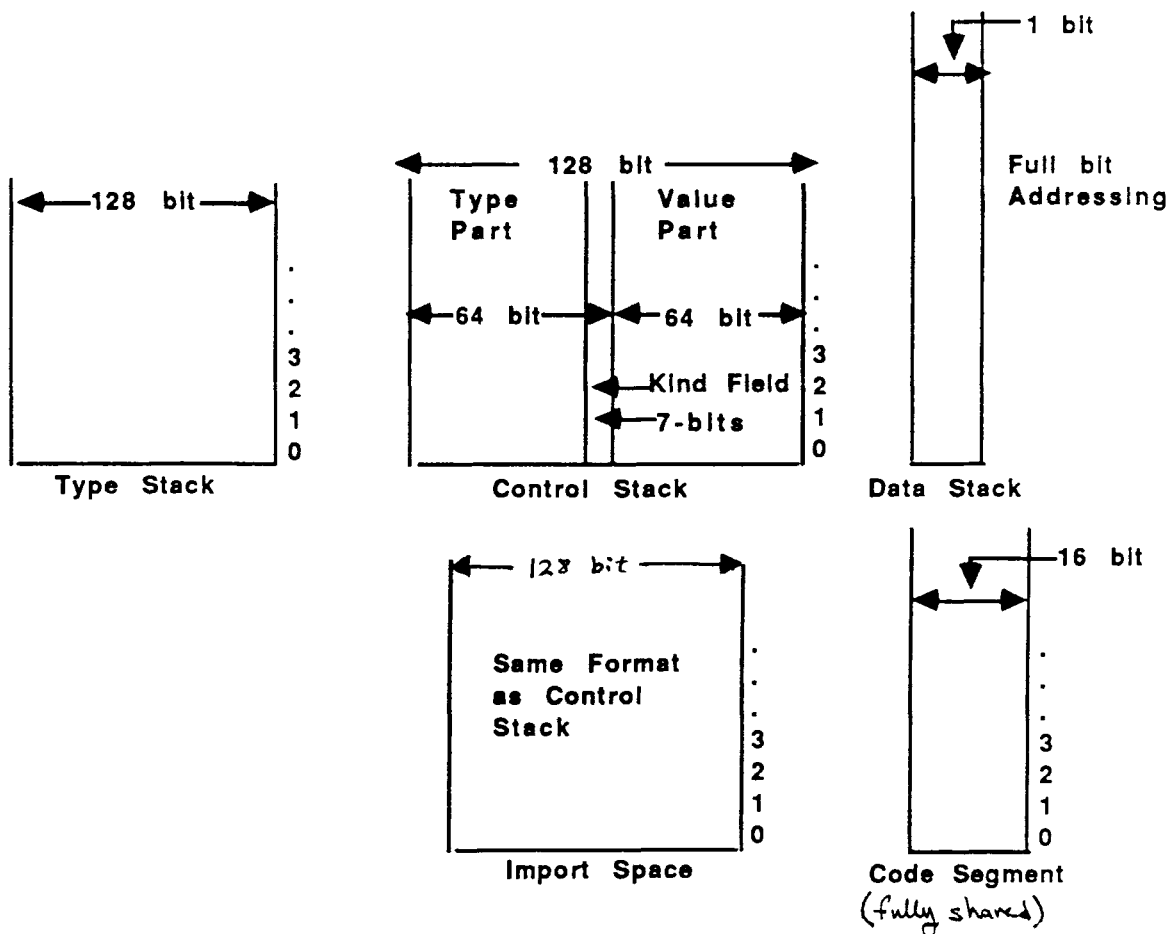
- **Tasks and packages are treated almost identically in the R1000 architecture**
- **While a package is elaborating it runs on its own "thread" of control**
- **Each package and task has its own id which is unique on that machine**
- **The task id (or more generally, module id) is the "value" part of a task or package object**
- **Each module takes several pages, minimum**  
*task or package*
- **Packages declare visible types, subprograms and data objects. Each of these exists at a specific offset within the package.** *marked visible or not*
- **Tasks are similar, but can only declare visible entries to be called**
- **Operations on tasks and packages include calling the entries of a task or the visible subprograms of a package**
- **More on this later**

# R1000 STORAGE ORGANIZATION - INTRODUCTION

- **A module (package or task) is the basic container of information**
- **Each module consists of several memory segments**  
**A segment is a contiguous region of virtual memory**
- **The basic segments for a module are:**
  - **Type stack - contains type descriptions for types declared in the module**
  - **Control stack - contains task control block (TCB), visible objects, and procedure call stack** *for packages also*
  - **Data stack - contains value part of arrays and records. Operates in parallel with control stack.**
  - **Import space - used to address non-local objects. More later.**
  - **Queue space - used to store entry queue information (tasks only)**

# R1000 STORAGE ORGANIZATION - SEGMENTS

- In addition to Control, Type, Data, Import and Queue segments, there are also:
  - Code segments - which contain instructions to be executed
  - System segments - which are used for internal machine state
- Segments (memory in general) are bit addressable. Memory address bus addresses are always bit addresses.
- From a software standpoint, some segments have common word addressing structure.



# OBJECTS - TYPE OBJECTS

- **Same as a normal object, but value part is null**
  - **Zero for discrete and float objects**
  - **Null pointer for record and array objects**
  - **Null task id for module objects**
- **Generic packages and task types are examples of module types**
- **Each package has such a type (even for non-generic packages)**

# OBJECTS - REVIEW

**Type part and value part**

**Type descriptions**

**Variant records - constrained vs unconstrained**

**Basic machine types**

**Packages and task objects**

**Objects of type "type"**

**Key areas of impact:**

- **Performance: costs in size and access time**
- **Unchecked\_Conversion: bounds and constrained bits are visible**
- **Debugging: misuse of unconstrained objects is a common customer mistake and has big impact on system performance and disk use**



# R1000 STORAGE ORGANIZATION - SEGMENTS (CONT'D)

- Each segment has a limit to its size:

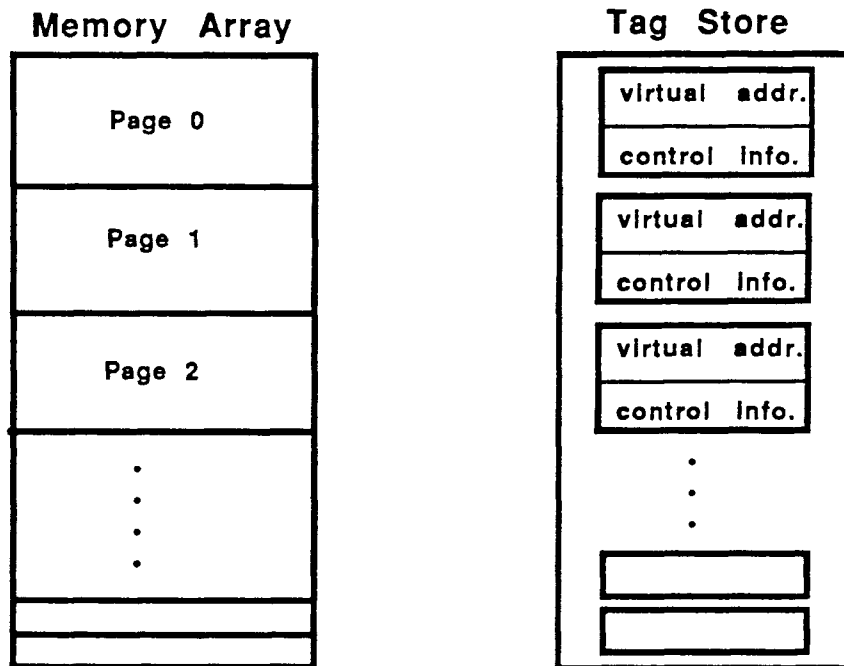
<u>Segment</u>	<u>Unit</u>	<u>Unit Limit</u>	<u>Byte Limit</u>
Control stack	128 bit word	$2^{20}$ (1M) words	16 MB
Type stack	128 bit word	$2^{20}$ (1M) words	16 MB
Data stack	Bit	$2^{32}$ (4G) bits	512 MB
Import	128 bit word	$2^{20}$ (1M) words	16 MB
Code	16 bit instr	32K instructions	64 KB

- Control stack depth limits recursive calls. An infinite recursion will consume 16 MB, then raise storage error

\* Debug tables can extend up to 128 KB  
*in Code segment*

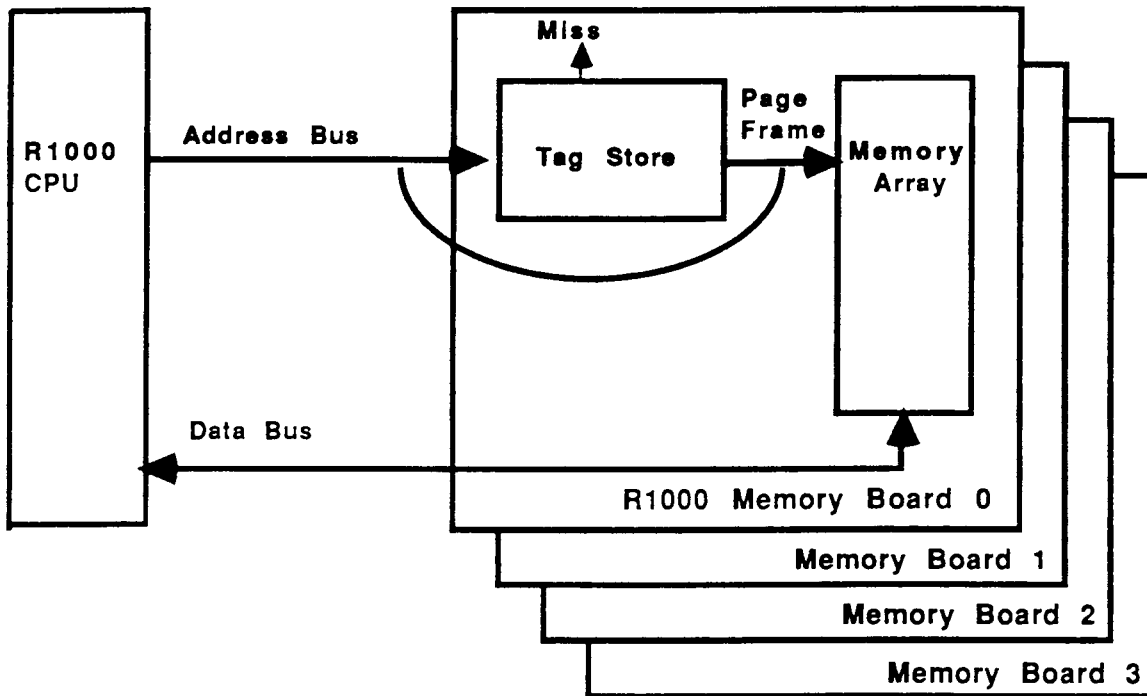
# R1000 STORAGE ORGANIZATION - PHYSICAL MEMORY

- R1000 physical memory consists of four 8 MB memory arrays.
- Memory is accessed in words of 128 bits (16 bytes) ECC is done on this unit.
- Memory modules respond directly to virtual addresses; there is no virtual -> physical address map.
- A Tag Store on each board identifies which virtual locations are present. *Associative*
- Memory is further divided into 1024 byte pages.





# R1000 STORAGE ORGANIZATION - PHYSICAL MEMORY (CONT'D)



- On a memory access, if all four boards generate a "miss" then the virtual address is not present in the physical memory and a page fault is signalled.
- The physical memory operates as a "cache" for virtual pages from disk.
- All tags on all boards are searched in parallel for fast access.

# R1000 STORAGE ORGANIZATION - REVIEW

- 7 segment <sup>types</sup> ~~bytes~~, each for different purpose
- Physical memory organization: tag store, ECC, 1K pages *128 bit bus*
- Segment size limits

# INSTRUCTION SET - INTRODUCTION

- **Code stored in code segments**
- **16 bit basic instructions; some 32 bit**
- **Stack oriented instructions**
  - **load (push)**
  - **load (push)**
  - **Add – no addresses in operator instruction**
  - **Store (pop)**
- **Basic instruction breakdown**
  - **Declarative instructions (4) - create types, variable, subprograms**
  - **Stack instructions (9) - push, pop**
  - **Subprogram call (3) - call, exit, end rendezvous**
  - **Classed action (2) - operate on typed operands**
  - **Unclassed action (1) - operate on untyped operands**
  - **Branch (6) - jump, case branch, counting loop**
  - **Literal (2) - load constants**
  - **Special (5) - markers, not really instructions**

# INSTRUCTION SET - CODE SEGMENT STRUCTURE

- Code segment consists of a header, a series of subprograms, and a debug table
- The header is the first 8 words (16 bit words) and has the following structure:

	0	4	8	15
0	version of machine code			
1	Diana version	RCG major version	RCG minor version	
2	0			
3	offset of debug table			
4	default exception handler - raise instruction			
5	module termination instruction - signal completion			
6	offset to segment table (only in elab segments)			
7	0	wired	# pages in seg - 1	

- Each subprogram has a 3 word header and begins on a 128 bit word boundary. There may be a few unused words at the end of a subprogram so that the next starts at the proper boundary.

The header is:

XXX0 offset of begin (first statement)

XXX1 offset of exception handler

XXX2 number of locals

XXX3 first instruction (entry point) of subprogram

# INSTRUCTION SET - GETTING CODE LISTINGS

- **Controlled by switches**

R1000_CG. Seg_Listing	- controls code segment listing
R1000_CG. Elab_Order_Listing	- shows order of module elaboration
R1000_CG. ASM_Listing	- produces earlier code generating listing - less useful
• <i>TTY_Echo</i>	

- **Don't forget to do Switches.Associate when you create a new switch file**
- **When the unit is coded, files will be created as children of the Ada unit:**

Foo	: C Proc_Spec
Foo	: C Main_Body
.<Code_Listing>	: File
.<Elaboration_Order_Listing>	: File
.<Elaboration_Code_Listing>	: File

# INSTRUCTION SET - SAMPLE

Short_Literal n	- push n on the control stack
Load_Encached n	- push contents of "special" register on the control stack
Declare_Type . . .	- create a new type (pushes a control word)
Complete_Type . . .	- fill in more info. about a type
Declare_Variable. . .	- create a new object (pushes a control word)
Declare subprogram	- create a new subprogram (32 bit instruction)
Load lex, offset	- push a copy of the specified word on the stack
Store lex, offset	- pop a word and store into the specified location
Execute class,field_write_op	-- store into a record, array, or package field
Exit_Subprogram n	- return and pop n words off the control stack
Action accept_activation	- finished package is visible part of elaboration
Action signal_activation	- finished package elaboration, starting body statements
Action break_optional	- debugger statement marker <i>inserted only if no other way for debugger to find statement boundary</i>

# INSTRUCTION SET - OVERVIEW

- **Declarative instructions create types, objects and subprograms. There are many flavors of these instructions**

- **Declare\_subprogram**

- **Declare\_variable**            **<class>**            **<options>**

- **Declare\_type**                **<class>**                **<options>**

- **Complete\_type**               **<class>**               **<options>**

- **Notes**

**<class>** refers to the object class. For example: record, array, task, etc.

# INSTRUCTION SET - OVERVIEW (CONT'D)

- **Stack Instructions**

**Load lex level, offset**

**Store lex level, offset**

**Store\_unchecked lex level, offset**                   – no constraint check

**Reference lex level, offset**                   – push a reference  
rather than a value

**Store\_top <class> <offset>**                   – tos relative address

**Store\_top\_unchecked <offset>**

**Load\_top <offset>**                   – tos relative address

**Pop\_control count**

**Load\_encached <register#>**                   – magic values

\* **tos = Top of Stack**

*continued following next page*



## Machine type summary

## Operand\_Class

-----

Discrete_Class	+, -, *, /, =, /=, etc
Float_Class	+, -, *, /, =, /=, etc
Access_Class	
Heap_Access_Class	
Package_Class	Declare, Activate, etc.
Task_Class	Declare, Abort, etc.
Module_Class	
Array_Class	Field_Read/Write, Get_Bounds, etc.
Vector_Class	Field_Read/Write, Get_Bounds, etc.
Matrix_Class	Field_Read/Write, Get_Bounds, etc.
Subarray_Class	Field_Read/Write, Get_Bounds, etc.
Subvector_Class	Field_Read/Write, Get_Bounds, etc.
Record_Class	Field_Read/Write
Variant_Record_Class	Field_Read/Write
Select_Class	call
Entry_Class	Rendezvous
Family_Class	Rendezvous
Exception_Class	Raise, Test_Value

# INSTRUCTION SET - OVERVIEW (CONT'D)

- **Subprogram call**

Call	lex level, offset
Exit_Subprogram	pop count, exit options
End_Rendezvous	parm count

also

Execute, Package, Field\_Execute - - call to another package

- **Unclassed actions**

Action	<operation>
	signal_activated
	swap_control
	initiate_delay
	etc.

General miscellaneous operations (about 56 of them).

- **Classed Actions - cover most operations on typed operands**

Execute	<class>	<operation>
	float	is_unsigned
	select	time_guard_write
	variant_record	field_read_dynamic
	module	check_elaborated
	array	concatenate
	discrete	not_equal
	etc.	

- **Note: Vector, Subvector, Subarray, and Matrix are special kinds of Array classes.**

```

package Machine_Code is

  pragma Subsystem (Ada_Base);

  pragma Module_Name (4, 67);

  Version : constant := 15;

  type Operand_Class is (Discrete_Class, Float_Class, Access_Class,
    Heap_Access_Class, Package_Class,
    Task_Class, Module_Class, Array_Class,
    Vector_Class, Matrix_Class, Subarray_Class,
    Subvector_Class, Record_Class,
    Variant_Record_Class, Select_Class, Entry_Class,
    Family_Class, Exception_Class, Any_Class);

  type Operation is (Equal_Op, Not_Equal_Op, Greater_Op, Less_Op,
    Greater_Equal_Op, Less_Equal_Op, Case_Compare_Op,
    Below_Bound_Op, Above_Bound_Op, In_Range_Op,
    Not_In_Range_Op, Case_In_Range_Op, Bounds_Check_Op,
    Reverse_Bounds_Check_Op, Bounds_Op, Reverse_Bounds_Op,
    First_Op, Last_Op, Set_Value_Op, Set_Value_Unchecked_Op,
    Set_Value_Visible_Op, Set_Value_Visible_Unchecked_Op,
    Test_And_Set_Previous_Op, Test_And_Set_Next_Op,
    Successor_Op, Predecessor_Op, And_Op, Or_Op, Xor_Op,
    Complement_Op, Unary_Minus_Op, Absolute_Value_Op,
    Plus_Op, Minus_Op, Times_Op, Divide_Op, Modulo_Op,
    Remainder_Op, Minimum_Op, Maximum_Op, Exponentiate_Op,
    Binary_Scale_Op, Multiply_And_Scale_Op,
    Divide_And_Scale_Op, Partial_Plus_Op, Partial_Minus_Op,
    Arithmetic_Shift_Op, Logical_Shift_Op,
    Rotate_Op, Insert_Bits_Op, Extract_Bits_Op,
    Is_Unsigned_Op, Count_Nonzero_Bits_Op,
    Count_Leading_Zeros_Op, Count_Trailing_Zeros_Op,
    Round_To_Discrete_Op, Truncate_To_Discrete_Op,
    Convert_From_Discrete_Op, Equal_Zero_Op,
    Not_Equal_Zero_Op, Greater_Zero_Op, Less_Zero_Op,
    Greater_Equal_Zero_Op, Less_Equal_Zero_Op, Is_Null_Op,
    Not_Null_Op, Set_Null_Op, Element_Type_Op, All_Read_Op,
    All_Write_Op, All_Reference_Op, Convert_Reference_Op,
    Allow_Deallocate_Op, Deallocate_Op, Hash_Op,
    Construct_Segment_Op, Get_Segment_Op, Get_Offset_Op,
    Diana_Tree_Kind_Op, Diana_Map_Kind_To_Vci_Op,
    Diana_Arity_For_Kind_Op, Diana_Allocate_Tree_Node_Op,
    Diana_Put_Node_On_Seq_Type_Op,
    Diana_Seq_Type_Get_Head_Op,
    Diana_Find_Permanent_Attribute_Op, Diana_Spare0_Op,
    Diana_Spare1_Op, Diana_Spare2_Op,
    Adaptive_Balanced_Tree_Lookup_Op,
    Elaborate_Op, Check_Elaborated_Op,
    Augment_Imports_Op, Activate_Op, Abort_Op,
    Abort_Multiple_Op, Rendezvous_Op, Count_Op,
    Is_Callable_Op, Is_Terminated_Op,
    Timed_Duration_Write_Op, Terminate_Guard_Write_Op,
    Timed_Guard_Write_Op, Length_Op,
    Structure_Write_Op, Structure_Clear_Op,
    Slice_Read_Op, Slice_Write_Op,

```

```

    Slice_Reference_Op, Catenate_Op,
    Append_Op, Prepend_Op, Subarray_Op,
    Component_Offset_Op, Is_Constrained_Op,
    Is_Constrained_Object_Op, Make_Constrained_Op,
    Indirects_Appended_Op, Structure_Query_Op,
    Reference_Makes_Copy_Op, Read_Variant_Op,
    Read_Discriminant_Constraint_Op,
    Field_Read_Op, Field_Write_Op,
    Field_Reference_Op, Field_Append_Op,
    Field_Type_Op, Field_Constrain_Op,
    Set_Bounds_Op, Set_Variant_Op,
    Field_Execute_Op, Entry_Call_Op,
    Conditional_Call_Op, Timed_Call_Op,
    Family_Call_Op, Family_Timed_Op,
    Family_Conditional_Op, Guard_Write_Op,
    Member_Write_Op, Write_Unchecked_Op,
    Field_Read_Dynamic_Op, Field_Write_Dynamic_Op,
    Field_Type_Dynamic_Op, Field_Execute_Dynamic_Op,
    Field_Reference_Dynamic_Op,
    Reference_Lex_1_Op, Convert_Op,
    Convert_Unchecked_Op, Convert_To_Formal_Op,
    Check_In_Type_Op, Check_In_Formal_Type_Op,
    Check_In_Integer_Op, In_Type_Op,
    Not_In_Type_Op, Address_Op, Size_Op,
    Address_Of_Type_Op, Make_Root_Type_Op,
    Set_Constraint_Op, Is_Value_Op,
    Is_Scalar_Op, Is_Default_Op,
    Make_Visible_Op, Make_Aligned_Op,
    Run_Utility_Op, Change_Utility_Op,
    Run_Initialization_Utility_Op,
    Has_Default_Initialization_Op,
    Has_Repeated_Initialization_Op,
    Is_Initialization_Repeated_Op, Raise_Op, Reraise_Op,
    Get_Name_Op, Is_Constraint_Error_Op,
    Is_Numeric_Error_Op, Is_Program_Error_Op,
    Is_Tasking_Error_Op, Is_Storage_Error_Op,
    Is_Instruction_Error_Op, Instruction_Read_Op,
    Spare14_Op, Spare15_Op, Spare16_Op,
    Spare17_Op, Spare18_Op, Spare19_Op,
    Spare20_Op, Spare21_Op, Spare22_Op,
    Spare23_Op, Spare24_Op, Spare25_Op);

```

```

type Type_Sort is (Defined, Constrained, Incomplete,
  Defined_Incomplete, Constrained_Incomplete);

```

```

type Type_Option_Set is
  record
    Is_Visible : Boolean;
    Unsigned : Boolean;
    With_Size : Boolean;
    Accesses_Protected : Boolean;
    Values_Relative : Boolean;
    Bounds_With_Object : Boolean;
    Not_Elaborated : Boolean;
  end record;

```

```

type Type_Completion_Mode is

```

(By\_Defining, By\_Renaming, By\_Constraining, By\_Constraining\_Incomplete,  
By\_Completing\_Constraint, By\_Component\_Completion);

```
type Variable_Option_Set is
  record
    Is_Visible : Boolean;
    Duplicate : Boolean;
    By_Allocation : Boolean;
    With_Constraint : Boolean;
    With_Subtype : Boolean;
    With_Value : Boolean;
    As_Component : Boolean;
    On_Processor : Boolean;
    Choice_Open : Boolean;
  end record;
```

```
type Subprogram_Sort is (For_Call, For_Outer_Call,
  For_Accept, Null_Subprogram);
```

```
type Subprogram_Option_Set is
  record
    Is_Visible : Boolean;
    Not_Elaborated : Boolean;
    With_Address : Boolean;
  end record;
```

```
type Exit_Option_Set is
  record
    With_Result : Boolean;
    From_Utility : Boolean;
  end record;
```

```
Field_Index_Size : constant Integer := 8;
```

```
type Field_Index is new Integer range 0 .. 2 ** Field_Index_Size - 1;
subtype Variant_Record_Index is Field_Index range 1 .. Field_Index'Last;
```

```
No_Variants : constant Variant_Record_Index := Field_Index'Last;
```

```
type Field_Sort is (Fixed, Variant);
type Field_Mode is (Direct, Indirect);
```

```
subtype Field_Op is Operation range Field_Read_Op .. Member_Write_Op;
subtype Component_Op is Field_Op range Field_Read_Op .. Field_Append_Op;
```

```
type Access_Spec (Op : Operation := Equal_Op) is
  record
    case Op is
      when Component_Op =>
        Kind : Field_Sort;
        Mode : Field_Mode;

      when others =>
```

```
    null;
  end case;
end record;
```

```
type Field_Spec (Class : Operand_Class := Discrete_Class;
  Op : Operation := Equal_Op) is
```

```
  record
    case Class is
      when Package_Class | Task_Class =>
        Offset : Field_Index;

      when Record_Class | Select_Class =>
        Number : Field_Index;

      when Variant_Record_Class =>
        Index : Field_Index;
        Component : Access_Spec (Op);

      when others =>
        null;
    end case;
  end record;
```

```
type Operator_Spec (Class : Operand_Class := Discrete_Class;
  Op : Operation := Equal_Op) is
```

```
  record
    case Op is
      when Field_Op =>
        Field : Field_Spec (Class, Op);

      when others =>
        null;
    end case;
  end record;
```

```
type Lexical_Level is new Integer range 0 .. 15;
type Lexical_Delta is new Integer range -256 .. 511;
```

```
subtype Scope_Delta is Lexical_Delta range 0 .. 511;
subtype Frame_Delta is Lexical_Delta range -256 .. 255;
```

```
type Object_Reference is
  record
    Level : Lexical_Level;
    Offset : Lexical_Delta;
  end record;
```

```
type Inner_Frame_Delta is new Integer range 0 .. 2 ** 8 - 1;
```

```
type Parameter_Count is new Integer range 0 .. 2 ** 8 - 1;
```

```
type Pc_Offset is new Integer range -2 ** 10 .. 2 ** 10 - 1;
```

```
subtype Loop_Offset is Pc_Offset range -2 ** 9 .. -1;
```

```
type Case_Maximum is new Integer range 0 .. 2 ** 9 - 1;
```

```

type Stack_Top_Offset is new Integer range -6 .. 0;
type Stack_Pop_Count is new Integer range 1 .. 7;
type Encached_Object_Number is new Integer range 0 .. 31;
type Symbolic_Label is new Integer;

type Segment_Index is new Integer range 0 .. 2 ** 3 - 1;
type Segment_Displacement is new Integer range 0 .. 2 ** 12 - 1;

type Segment_Reference is
  record
    Offset : Segment_Displacement;
    Index : Segment_Index;
  end record;

type Short_Literal_Value is new Integer range -2 ** 10 .. 2 ** 10 - 1;
type Partial_Structure is new Long_Integer range 0 .. 2 ** 32 - 1;
type Structure_Literal_Value is
  record
    High : Partial_Structure;
    Low : Partial_Structure;
  end record;

type Literal (Of_Kind : Operand_Class := Discrete_Class) is
  record
    case Of_Kind is
      when Discrete_Class =>
        Discrete_Literal : Long_Integer;

      when Float_Class =>
        Float_Literal : Float;

      when Any_Class =>
        Structure_Literal : Structure_Literal_Value;

      when others =>
        null;
    end case;
  end record;

type Unsigned_Immediate_Value is new Integer range 0 .. 255;
type Signed_Immediate_Value is new Integer range -128 .. 127;

type Immediate_Value (Is_Signed : Boolean := True) is
  record
    case Is_Signed is
      when True =>
        Signed_Value : Signed_Immediate_Value;

      when False =>
        Unsigned_Value : Unsigned_Immediate_Value;
    end case;
  end record;

```

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```

    end case;
  end record;

type Extended_Literal is new Integer range 0 .. 2 ** 16 - 1;
type Extension_Kind is (Is_Literal, Is_Location);
type Instruction_Extension (Of_Kind : Extension_Kind := Is_Literal) is
  record
    case Of_Kind is
      when Is_Literal =>
        Value : Extended_Literal;

      when Is_Location =>
        Location : Segment_Reference;
    end case;
  end record;

type Unclassed_Action is
  (Illegal, Idle,

  Elaborate_Subprogram, Check_Subprogram_Elaborated,

  Accept_Activation, Activate_Tasks, Activate_Heap_Tasks,
  Signal_Activated, Signal_Completion, Propagate_Abort,
  Set_Priority, Increase_Priority, Get_Priority,
  Initiate_Delay, Exit_Nullary_Function, Set_Block_Start,

  Make_Default, Make_Self, Make_Scope, Make_Parent, Name_Partner,

  Swap_Control, Mark_Auxiliary, Pop_Auxiliary,
  Pop_Auxiliary_Loop, Pop_Auxiliary_Range,

  Pop_Block, Pop_Block_With_Result,

  Break_Unconditional, Break_Optional, Exit_Break,
  Query_Break_Cause, Query_Break_Address, Query_Break_Mask,
  Alter_Break_Mask, Query_Frame, Establish_Frame,

  Loop_Increasing_Extended, Loop_Decreasing_Extended,
  Push_Discrete_Extended, Push_Float_Extended,
  Push_Structure_Extended, Push_String_Extended,
  Push_String_Extended_Indexed, Store_String_Extended,

  Jump_Extended, Jump_Zero_Extended, Jump_Nonzero_Extended,
  Jump_Dynamic, Jump_Zero_Dynamic, Jump_Nonzero_Dynamic,

  Load_Dynamic, Store_Dynamic, Call_Dynamic, Reference_Dynamic,

  Spare6_Action, Spare7_Action, Spare8_Action,
  Spare9_Action, Spare10_Action, Spare11_Action);

type Op_Code is (Declare_Type, Complete_Type, Declare_Variable,
  Declare_Subprogram, Load, Store, Store_Unchecked,

```

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```
Reference, Call, Exit_Subprogram, End_Rendezvous, Execute,
Action, Execute_Immediate, Jump, Jump_Zero, Jump_Nonzero,
Jump_Case, Loop_Increasing, Loop_Decreasing, Short_Literal,
Indirect_Literal, Pop_Control, Load_Top, Store_Top,
Store_Top_Unchecked, Load_Encached, Extension,
Literal_Value, Block_Begin, Block_Handler, End_Locals);
```

```
type Instruction (For_Op : Op_Code := Action) is
record
  case For_Op is
    when Declare_Type =>
      Type_Class : Operand_Class;
      Type_Kind : Type_Sort;
      Type_Options : Type_Option_Set;

    when Complete_Type =>
      Completion_Class : Operand_Class;
      Completion_Mode : Type_Completion_Mode;

    when Declare_Variable =>
      Variable_Class : Operand_Class;
      Variable_Options : Variable_Option_Set;

    when Declare_Subprogram =>
      Subprogram_Kind : Subprogram_Sort;
      Subprogram_Options : Subprogram_Option_Set;

    when Load | Store | Store_Unchecked | Reference | Call =>
      Object : Object_Reference;

    when Exit_Subprogram =>
      New_Top_Offset : Inner_Frame_Delta;
      Exit_Options : Exit_Option_Set;

    when End_Rendezvous =>
      Return_Count : Parameter_Count;

    when Execute =>
      Operator : Operator_Spec;

    when Action =>
      To_Perform : Unclassed_Action;

    when Execute_Immediate =>
      Discrete_Op : Operation;
      With_Value : Immediate_Value;

    when Jump | Jump_Zero | Jump_Nonzero |
      Loop_Increasing | Loop_Decreasing =>
      Relative : Pc_Offset;

    when Jump_Case =>
      Case_Max : Case_Maximum;

    when Pop_Control =>
      Pop_Count : Stack_Pop_Count;

    when Load_Top =>
      At_Offset : Stack_Top_Offset;
```

```
when Store_Top | Store_Top_Unchecked =>
  Target_Class : Operand_Class;
  Target_Offset : Stack_Top_Offset;

when Load_Encached =>
  With_Number : Encached_Object_Number;

when Short_Literal =>
  Short_Value : Short_Literal_Value;

when Indirect_Literal =>
  Value_Class : Operand_Class;
  Value_Relative : Pc_Offset;

when Literal_Value =>
  Value : Literal;

when Extension =>
  Argument : Instruction_Extension;

when Block_Begin | Block_Handler =>
  Location : Segment_Reference;

when End_Locals =>
  Offset : Extended_Literal;
end case;
end record;
```

-- indices into the encaching registers for declarations from Standard

```
Standard_Cache_Index : constant Encached_Object_Number := 0;
Boolean_Cache_Index : constant Encached_Object_Number := 1;
Integer_Cache_Index : constant Encached_Object_Number := 2;
Natural_Cache_Index : constant Encached_Object_Number := 3;
Positive_Cache_Index : constant Encached_Object_Number := 4;
Long_Integer_Cache_Index : constant Encached_Object_Number := 5;
Float_Cache_Index : constant Encached_Object_Number := 6;
Duration_Cache_Index : constant Encached_Object_Number := 7;
Character_Cache_Index : constant Encached_Object_Number := 8;
String_Cache_Index : constant Encached_Object_Number := 9;
Null_String_Cache_Index : constant Encached_Object_Number := 10;
```

-- indices into the encaching registers for declarations from Diana

```
Diana_Cache_Index : constant Encached_Object_Number := 11;
Diana_Tree_Cache_Index : constant Encached_Object_Number := 12;
Diana_Symbol_Rep_Cache_Index : constant Encached_Object_Number := 13;
Diana_Seq_Type_Cache_Index : constant Encached_Object_Number := 14;
Diana_Sequence_Cache_Index : constant Encached_Object_Number := 15;
```

-- indices into the encaching registers for declarations from System

```
Segment_Cache_Index : constant Encached_Object_Number := 16;
```

-- more indices into the encaching registers for declarations from Diana

```
Diana_Temp_Seq_Index : constant Encached_Object_Number := 17;
Diana_Attr_List_Index : constant Encached_Object_Number := 18;
Diana_Tree_Node_Index : constant Encached_Object_Number := 19;
Diana_Seq_Type_Node_Index : constant Encached_Object_Number := 20;
```

-- unused indices

```
Unused_Index_21 : constant Encached_Object_Number := 21;  
Unused_Index_22 : constant Encached_Object_Number := 22;  
Unused_Index_23 : constant Encached_Object_Number := 23;  
Unused_Index_24 : constant Encached_Object_Number := 24;  
Unused_Index_25 : constant Encached_Object_Number := 25;  
Unused_Index_26 : constant Encached_Object_Number := 26;  
Unused_Index_27 : constant Encached_Object_Number := 27;  
Unused_Index_28 : constant Encached_Object_Number := 28;  
Unused_Index_29 : constant Encached_Object_Number := 29;  
Unused_Index_30 : constant Encached_Object_Number := 30;  
Unused_Index_31 : constant Encached_Object_Number := 31;
```

end Machine\_Code;

# INSTRUCTION SET - DEBUGGER

- You can display code and monitor PC values with the R1000 debugger.

`Memory_Display ("seg offset", count, "code")`

`Memory_Display ("17634 250", 30, "code");`

↑  
hex segment name and offset

- **Enable (Addresses)**
  - Causes debugger to display PC and other machine information in debugger displays.



```

with Io;
with Job_Segment;

procedure Architecture_Example is
  package P1 is
    procedure Op (X : String);
  end P1;
  package body P1 is

    Local_Value : Integer := 33;
    procedure Local_Procedure (S : String) is
    begin
      Io.Put_Line (S & Integer'Image (Local_Value));
    end Local_Procedure;

    procedure Op (X : String) is
      Local_Variable : Integer;
      type A is array (1 .. 10) of Boolean;
      Ar : A;
      type Pa is access A;
      type Sa is access A;
      pragma Segmented_Heap (Sa);
      P1 : Pa := new A'(others => True);
      P2 : Sa := new A'(others => True);
      pragma Heap (Job_Segment.Get);

      procedure Nested (Parm : Natural) is
        Nested_Local : Integer;
      begin
        Local_Variable := Parm;
        Nested_Local := Local_Variable;
        Local_Value := Parm;
        P2 (3) := False;
        null;
      end Nested;

      begin
        Local_Procedure (X & "hi");
        Nested (23);
      end Op;

    end P1;
  -- hi
  begin
    P1.Op ("this is a test");
    null;
  end Architecture_Example;
pragma Main;

```

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1

Code for segment 1033472

```
0 0000: 000F
1 0001: 5800
2 0002: 0000
3 0003: 00B8
4 0004: 0100
5 0005: 00BB
6 0006: 0000
7 0007: 0000
```

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2

```
8 0008: 0010 BLOCK_BEGIN      2, 0
9 0009: 0004 BLOCK_HANDLER    0, 4
10 000A: 0001 END_LOCALS      1
```

;;; module ARCHITECTURE EXAMPLE

```
11 000B: 029A DECLARE_SUBPROGRAM FOR_OUTER_CALL, IS_VISIBLE,
                                         NOT_ELABORATED
12 000C: 001B EXTENSION        3, 3      ;; ARCHITECTURE_EXAMPLE
13 000D: 00BF ACTION            ACCEPT_ACTIVATION
14 000E: 1D0D EXECUTE_IMMEDIATE REFERENCE_LEX_1_OP, 13
15 000F: 00C7 ACTION            ELABORATE_SUBPROGRAM
16 0010: 00BC ACTION            SIGNAL_ACTIVATED
17 0011: 00BB ACTION            SIGNAL_COMPLETION
18 0012: 0000
19 0013: 0000
20 0014: 0000
21 0015: 0000
22 0016: 0000
23 0017: 0000
```

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```

24 0018: 0029 BLOCK BEGIN      5, 1
25 0019: 0004 BLOCK_HANDLER  0, 4
26 001A: 0002 END_LOCALS      2

;;; subprogram ARCHITECTURE_EXAMPLE
;;; package P1 is
27 001B: 4800 SHORT_LITERAL  0
28 001C: 0093 EXTENSION      147
29 001D: 0033 EXTENSION      6, 3    ;;; push full address of a location in
                                         current code segment
30 001E: 02A0 DECLARE_SUBPROGRAM NULL SUBPROGRAM
31 001F: 038E DECLARE_TYPE    PACKAGE_CLASS, DEFINED    ;;; P1
32 0020: 0387 DECLARE_VARIABLE PACKAGE_CLASS    ;;; P1

;;; package body P1 is
33 0021: E001 LOAD            0, 1    ;;; DEFAULT
34 0022: E002 LOAD            0, 2    ;;; JOB_SEGMENT
35 0023: E003 LOAD            0, 3    ;;; IO
36 0024: 4803 SHORT_LITERAL    3
37 0025: E402 LOAD            2, 2    ;;; P1 ;;; P1
38 0026: 020E EXECUTE         MODULE_CLASS, AUGMENT_IMPORTS_OP
39 0027: 00D8 LOAD_TOP        0 ;;; ( 2, 2) ;;; P1
40 0028: 020F EXECUTE         MODULE_CLASS, ACTIVATE_OP

;;; P1.OP ("this is a test");
;;; push parameters
41 0029: 0092 ACTION          PUSH_STRING_EXTENDED
42 002A: 00A8 EXTENSION        21, 0
43 002B: E402 LOAD            2, 2    ;;; P1
44 002C: 180D EXECUTE         PACKAGE_CLASS, FIELD_EXECUTE_OP, 13    ;;; OP

;;; null;
45 002D: 4501 EXIT_SUBPROGRAM  1
46 002E: 0000
47 002F: 0000

```

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```

48 0030: 003C BLOCK BEGIN      7, 4
49 0031: 0004 BLOCK_HANDLER  0, 4
50 0032: 000F END_LOCALS      15

;;; module P1
;;; procedure OP (X : STRING);
51 0033: 029A DECLARE_SUBPROGRAM FOR_OUTER_CALL, IS_VISIBLE,
                                         NOT_ELABORATED
52 0034: 0043 EXTENSION        8, 3    ;;; OP
53 0035: 00BF ACTION           ACCEPT_ACTIVATION

;;; LOCAL VALUE : INTEGER := 33;
54 0036: 00E2 LOAD_ENCACHED    2 ;;; INTEGER
55 0037: 0621 EXECUTE_IMMEDIATE SET_VALUE_UNCHECKED_OP, 33

;;; procedure LOCAL_PROCEDURE (S : STRING) is
56 0038: 029D DECLARE_SUBPROGRAM FOR_OUTER_CALL
57 0039: 008B EXTENSION        17, 3    ;;; LOCAL_PROCEDURE

;;; procedure OP (X : STRING) is
58 003A: 1D0D EXECUTE_IMMEDIATE REFERENCE_LEX_1_OP, 13
59 003B: 00C7 ACTION           ELABORATE_SUBPROGRAM
60 003C: 00BC ACTION           SIGNAL_ACTIVATED
61 003D: 00BB ACTION           SIGNAL_COMPLETION
62 003E: 0000
63 003F: 0000

```

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#4b

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5

```

64 0040: 0079 BLOCK_BEGIN      15, 1
65 0041: 0004 BLOCK_HANDLER    0, 4
66 0042: 0009 END_LOCALS       9

;;; subprogram OP
;;; LOCAL_VARIABLE : INTEGER;
67 0043: 00E2 LOAD_ENCACHED    2      ;;; INTEGER

;;; type A is array (1 .. 10) of BOOLEAN;
68 0044: 4801 SHORT_LITERAL    1      ;;; 1
69 0045: 480A SHORT_LITERAL    10     ;;; 10
70 0046: 00E1 LOAD_ENCACHED    1      ;;; BOOLEAN
71 0047: 02A0 DECLARE SUBPROGRAM NULL_SUBPROGRAM
72 0048: 4801 SHORT_LITERAL    1
73 0049: 035D DECLARE_TYPE     ARRAY_CLASS, DEFINED      ;;; A

;;; AR : A;
74 004A: 00D8 LOAD_TOP         0      ;;; ( 2, 3)      ;;; A
75 004B: 0337 DECLARE_VARIABLE ARRAY_CLASS

;;; type PA is access A;
76 004C: E403 LOAD             2, 3    ;;; A
77 004D: 02A0 DECLARE SUBPROGRAM NULL_SUBPROGRAM
78 004E: 4818 SHORT_LITERAL    24
79 004F: 03D5 DECLARE_TYPE     ACCESS_CLASS, DEFINED      ;;; PA

;;; type SA is access A;
80 0050: E403 LOAD             2, 3    ;;; A
81 0051: 02A0 DECLARE SUBPROGRAM NULL_SUBPROGRAM
82 0052: 03AD DECLARE_TYPE     HEAP_ACCESS_CLASS, DEFINED ;;; SA

;;; pragma SEGMENTED HEAP (SA);
83 0053: 0007 ACTION          BREAK_OPTIONAL

;;; P1 : PA := new A' (others => TRUE);
84 0054: E405 LOAD             2, 5    ;;; PA
85 0055: E403 LOAD             2, 3    ;;; A
86 0056: 0337 DECLARE_VARIABLE ARRAY_CLASS
87 0057: 480A SHORT_LITERAL    10
88 0058: 4801 SHORT_LITERAL    1
89 0059: 4801 SHORT_LITERAL    1      ;;; TRUE
90 005A: E40A LOAD             2, 10
91 005B: E408 LOAD             2, 8
92 005C: 01D6 EXECUTE          VECTOR_CLASS, FIELD_WRITE_OP
93 005D: 3FFB LOOP_INCREASING  -5
94 005E: E403 LOAD             2, 3    ;;; A
95 005F: 01C3 EXECUTE          VECTOR_CLASS, CHECK_IN_TYPE_OP ;;; A' (others =>
TRUE)

96 0060: E405 LOAD             2, 5    ;;; PA
97 0061: 03B6 DECLARE_VARIABLE ACCESS_CLASS, BY_ALLOCATION, WITH_VALUE
;;; new A' (others => TRUE)

98 0062: C407 STORE            2, 7

;;; P2 : SA := new A' (others => TRUE);
99 0063: E406 LOAD             2, 6    ;;; SA
100 0064: E403 LOAD             2, 3    ;;; A
101 0065: 0337 DECLARE_VARIABLE ARRAY_CLASS
102 0066: 480A SHORT_LITERAL    10
103 0067: 4801 SHORT_LITERAL    1

```

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6

```

104 0068: 4801 SHORT_LITERAL    1      ;;; TRUE
105 0069: E40B LOAD             2, 11
106 006A: E409 LOAD             2, 9
107 006B: 01D6 EXECUTE          VECTOR_CLASS, FIELD_WRITE_OP
108 006C: 3FFB LOOP_INCREASING  -5
109 006D: E403 LOAD             2, 3    ;;; A
110 006E: 01C3 EXECUTE          VECTOR_CLASS, CHECK_IN_TYPE_OP ;;; A' (others =>
TRUE)

111 006F: E406 LOAD             2, 6    ;;; SA

;;; push parameters
112 0070: E001 LOAD             0, 1    ;;; DEFAULT
113 0071: 1811 EXECUTE          PACKAGE_CLASS, FIELD_EXECUTE_OP, 17 ;;; PROCESS
114 0072: E002 LOAD             0, 2    ;;; JOB_SEGMENT
115 0073: 1811 EXECUTE          PACKAGE_CLASS, FIELD_EXECUTE_OP, 17 ;;; GET
116 0074: 0396 DECLARE_VARIABLE HEAP_ACCESS_CLASS, BY_ALLOCATION,
WITH_VALUE ;;; new A' (others => TRUE)

117 0075: C408 STORE            2, 8

;;; pragma HEAP (JOB_SEGMENT.GET);
118 0076: 0007 ACTION          BREAK_OPTIONAL

;;; procedure NESTED (PARM : NATURAL) is
119 0077: 029F DECLARE SUBPROGRAM FOR_CALL
120 0078: 009B EXTENSION       19, 3    ;;; NESTED

;;; LOCAL_PROCEDURE (X & "h1");
;;; push parameters
121 0079: E5FF LOAD             2, -1   ;;; X
122 007A: 0092 ACTION          PUSH_STRING_EXTENDED
123 007B: 00A9 EXTENSION       21, 1
124 007C: 01CC EXECUTE          VECTOR_CLASS, CATENATE_OP
125 007D: 820F CALL             1, 15   ;;; LOCAL_PROCEDURE

;;; NESTED (23);
;;; push parameters
126 007E: 4817 SHORT_LITERAL    23     ;;; 23
127 007F: 8409 CALL             2, 9    ;;; NESTED
128 0080: 4502 EXIT_SUBPROGRAM  2
129 0081: 0000
130 0082: 0000
131 0083: 0000
132 0084: 0000
133 0085: 0000
134 0086: 0000
135 0087: 0000

```

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```

136 0088: 008B BLOCK_BEGIN      17, 3
137 0089: 0004 BLOCK_HANDLER    0, 4
138 008A: 0001 END_LOCALS      1

;;; subprogram LOCAL_PROCEDURE
;;; IO.PUT_LINE (S & INTEGER' IMAGE (LOCAL_VALUE));
;;; push parameters
139 008B: E5FF LOAD      2, -1    ;;; S
140 008C: E20E LOAD      1, 14    ;;; LOCAL_VALUE
141 008D: 00E0 LOAD_ENCACHED 0      ;;; Standard
142 008E: 1818 EXECUTE    PACKAGE_CLASS, FIELD_EXECUTE_OP, 24
143 008F: 01CC EXECUTE    VECTOR_CLASS, CATENATE_OP
144 0090: E003 LOAD      0, 3      ;;; IO
145 0091: 1864 EXECUTE    PACKAGE_CLASS, FIELD_EXECUTE_OP, 100 ;;;
                                                                    PUT_LINE

146 0092: 4502 EXIT_SUBPROGRAM  2
147 0093: 0000
148 0094: 0000
149 0095: 0000
150 0096: 0000
151 0097: 0000
    
```

```

152 0098: 009C BLOCK_BEGIN      19, 4
153 0099: 0004 BLOCK_HANDLER    0, 4
154 009A: 0002 END_LOCALS      2

;;; subprogram NESTED
;;; NESTED_LOCAL : INTEGER;
155 009B: 00E2 LOAD_ENCACHED  2      ;;; INTEGER

;;; LOCAL_VARIABLE := PARM;
156 009C: E7FF LOAD      3, -1    ;;; PARM
157 009D: A402 STORE_UNCHECKED 2, 2  ;;; LOCAL_VARIABLE

;;; NESTED_LOCAL := LOCAL_VARIABLE;
158 009E: E402 LOAD      2, 2      ;;; LOCAL_VARIABLE
159 009F: A602 STORE_UNCHECKED 3, 2  ;;; NESTED_LOCAL

;;; LOCAL_VALUE := PARM;
160 00A0: E7FF LOAD      3, -1    ;;; PARM
161 00A1: A20E STORE_UNCHECKED 1, 14 ;;; LOCAL_VALUE

;;; P2.all (3) := FALSE;
162 00A2: 4800 SHORT_LITERAL 0      ;;; FALSE
163 00A3: 4803 SHORT_LITERAL 3      ;;; 3
164 00A4: E408 LOAD      2, 8      ;;; P2
165 00A5: 0217 EXECUTE    HEAP_ACCESS_CLASS, ALL_REFERENCE_OP ;;; P2.all
166 00A6: 01D6 EXECUTE    VECTOR_CLASS, FIELD_WRITE_OP ;;; P2.all (3)

;;; null;
167 00A7: 4502 EXIT_SUBPROGRAM  2
    
```

#4b

!USERS.PHIL.GURU COURSE.ARCHITECTURE EXAMPLE'BODY.<CODE LISTING>'V(2)

9

```
;;; String Pointer table
168 00A8: 0006 0012 0012

;;; String Literal Pool
171 00AB: 7468 6973 2069 7320 6120 7465 7374 6869 {this is a test!}
179 00B3: 0000
180 00B4: 0000
181 00B5: 0000
182 00B6: 0000
183 00B7: 0000
```

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!USERS.PHIL.GURU COURSE.ARCHITECTURE EXAMPLE'BODY.<CODE LISTING>'V(2)

10

```
;;; debug table start
184 00B8: 0004 0000 5397 0000 0001 0007 000B 0000
192 00C0: 8000 0000 001B 00DA 8001 0000 0033 00E0
200 00C8: 0000 0000 0036 00E3 0000 0000 0043 00E8
208 00D0: 0001 0000 008B 00F6 0001 0000 009B 00F9
216 00D8: 0001 0000 0002 0002 001B 0021 0029 002D
224 00E0: 0001 0000 0033 0003 0000 0036 0038 003A
232 00E8: 000A 0002 0043 0044 004A 004C 0050 0053
240 00F0: 0054 0063 0076 0077 0079 007E 0000 0001
248 00F8: 008B 0001 0005 009B 009C 009E 00A0 00A2
256 0100: 00A7
```

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RI000 Ada Debugger (Version 9.1.6 - 7/11/87)  
 Beginning to debug: !USERS.PHIL.TEST\_AREA & !USERS.PHIL.TEST\_AREA.TASKY  
 Stop at: .command\_procedure, Root task: [Task : ROOT\_TASK, #794DD].

Enable (ADDRESSES, TRUE);  
 The ADDRESSES flag has been set to TRUE.

Break (".TASKY.JUNK.CHUCK.18", 1, "all", "TRUE");  
 The breakpoint has been created and activated:  
 Active Permanent Break 1 at .TASKY.JUNK.CHUCK.18 [any task]

Execute ("all");  
 Break 1: .TASKY.JUNK.CHUCK.18 [Task : #7C8DD].  
 Pc = #21D03, #9B.  
 Break 1: .TASKY.JUNK.CHUCK.18 [Task : #7C4DD].  
 Pc = #21D03, #9B.  
 Break 1: .TASKY.JUNK.CHUCK.18 [Task : #7CCDD].  
 Pc = #21D03, #9B.

Stack ("\*#7CCDD", 0, 0);  
 Stack of task #7CCDD:  
 -1: CHUCK.18  
   Pc = #21D03, #9B, Frame = #7CCDD, #18, Lex = 2, Outer = #7C0DD, #0  
 -2: BUSY2.2s  
   Pc = #21D03, #8D, Frame = #7CCDD, #12, Lex = 2, Outer = #7C0DD, #0  
 -3: BUSY.6s  
   Pc = #21D03, #B9, Frame = #7CCDD, #0, Lex = 1, Outer = #7CCDD, #0

Information (Space, "all");

Job: 221, Root task: #794DD  
 Space use for task #7C4DD:  
 Current control space: 1632, Max control space: 3056 (bytes).  
 Current data space: 514, Max data space: 2047 (bytes).  
 Address space sizes (in bits):

segment	current	max
Control	13056	24448
Data	4112	16383
Type	0	0
Queue		0

Space use for task ROOT\_TASK, #794DD:  
 Current control space: 752, Max control space: 8160 (bytes).  
 Current data space: 24, Max data space: 2138111 (bytes).  
 Address space sizes (in bits):

segment	current	max
Control	3328	57216
Data	193	17104895
Type	2688	8064
Queue		0

Space use for task #7C8DD:  
 Current control space: 1632, Max control space: 3056 (bytes).  
 Current data space: 514, Max data space: 2047 (bytes).  
 Address space sizes (in bits):

segment	current	max
Control	13056	24448
Data	4112	16383
Type	0	0
Queue		0

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Space use for task #7CCDD:  
 Current control space: 1632, Max control space: 3056 (bytes).  
 Current data space: 514, Max data space: 2047 (bytes).  
 Address space sizes (in bits):

segment	current	max
Control	13056	24448
Data	4112	16383
Type	0	0
Queue		0

Task\_Display ("all", ALL\_TASKS);

Job: 221, Root task: #794DD  
 #7C4DD (.TASKY.BUSY): At a breakpoint at .TASKY.JUNK.CHUCK.18 [Pri = 1,  
 Memory cond = 0]  
 Stopped in debugger.  
 ROOT\_TASK, #794DD: Running, waiting for delay [Pri = 1, Memory cond = 0]  
 #7C8DD (.TASKY.BUSY): At a breakpoint at .TASKY.JUNK.CHUCK.18 [Pri = 1,  
 Memory cond = 0]  
 Stopped in debugger.  
 #7CCDD (.TASKY.BUSY): At a breakpoint at .TASKY.JUNK.CHUCK.18 [Pri = 1,  
 Memory cond = 0]  
 Stopped in debugger.

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# **INSTRUCTION SET - REVIEW**

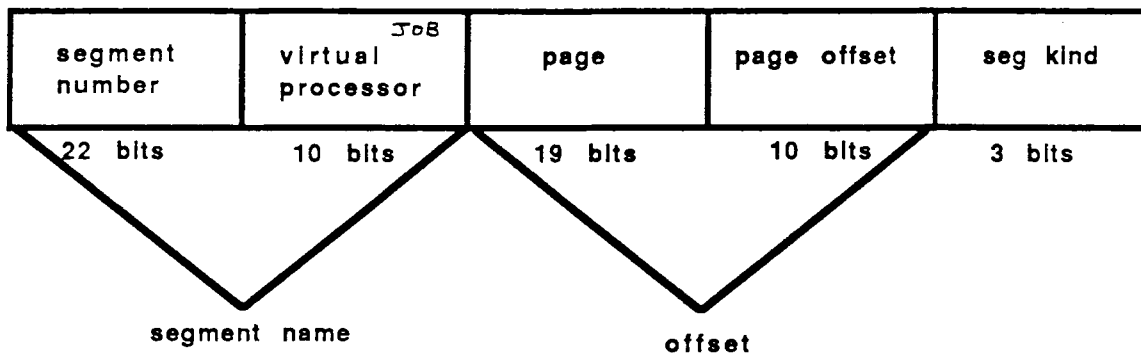
- **Code Segments**
- **Basic flavor of instructions**
- **Getting code segment listings for a unit**
- **Using the R1000 debugger to look at code segments**



# STORAGE HIERARCHY

- Cache memory (RAM) and disk interact closely for ALL memory use by the R1000 CPU.
- No separation of program execution virtual space and file space.
- If the disks get full, you cannot execute instructions on the machine, and you are hosed.
- Full virtual address

*page address*



- The segment kind specifies which of control, type, data, import, code, queue, and system segments are to be referenced.
- The virtual processor is part of the segment name. More later.
- Segments are divided into 1 Kbyte pages.

# **STORAGE HIERARCHY - KERNEL**

- **The kernel subsystem manages segments on the disk.**
- **The microcode manages segments in the main memory (cache).**
- **Segments can be created, used, and deleted without ever reaching disk.**
- **Tasks cannot take page faults until the virtual memory system is up.**

# **STORAGE HIERARCHY - SEGMENT NAMES**

- **Each module (package or task) has a name.**
- **This name is the segment name of the module's control, type and data stacks.**
- **The import space for a module has a different name.**
- **Code segments also have different names.**
- **Independent data objects (files and Ada units, for example) exist in their own data stacks. Each of these segments has a name.**

# EXECUTION ENVIRONMENT

- Each active thread of control (task or package during elaboration) has a control stack.
- The control stack contains a Task Control Block (TCB) and the procedure call stack.
- The stack is divided into frames. Each frame contains the information for a particular procedure activation.
- The bottom frame represents the task or package outer block.

# EXECUTION ENVIRONMENT - ADA ADDRESSING

With Text\_IO;

package body P is

    A: Natural

    procedure Proc ( x: Natural) is

        y: Natural

        procedure Proc\_Inner ( Z: Natural) is

            w: Natural

        begin

            if Z > 0 then

                y: =Z;

                Text\_IO.Put (A);

                Proc\_Inner (Z-1);

            end if;

        end Proc\_Inner;

    begin

        Proc\_Inner (X);

    end Proc;

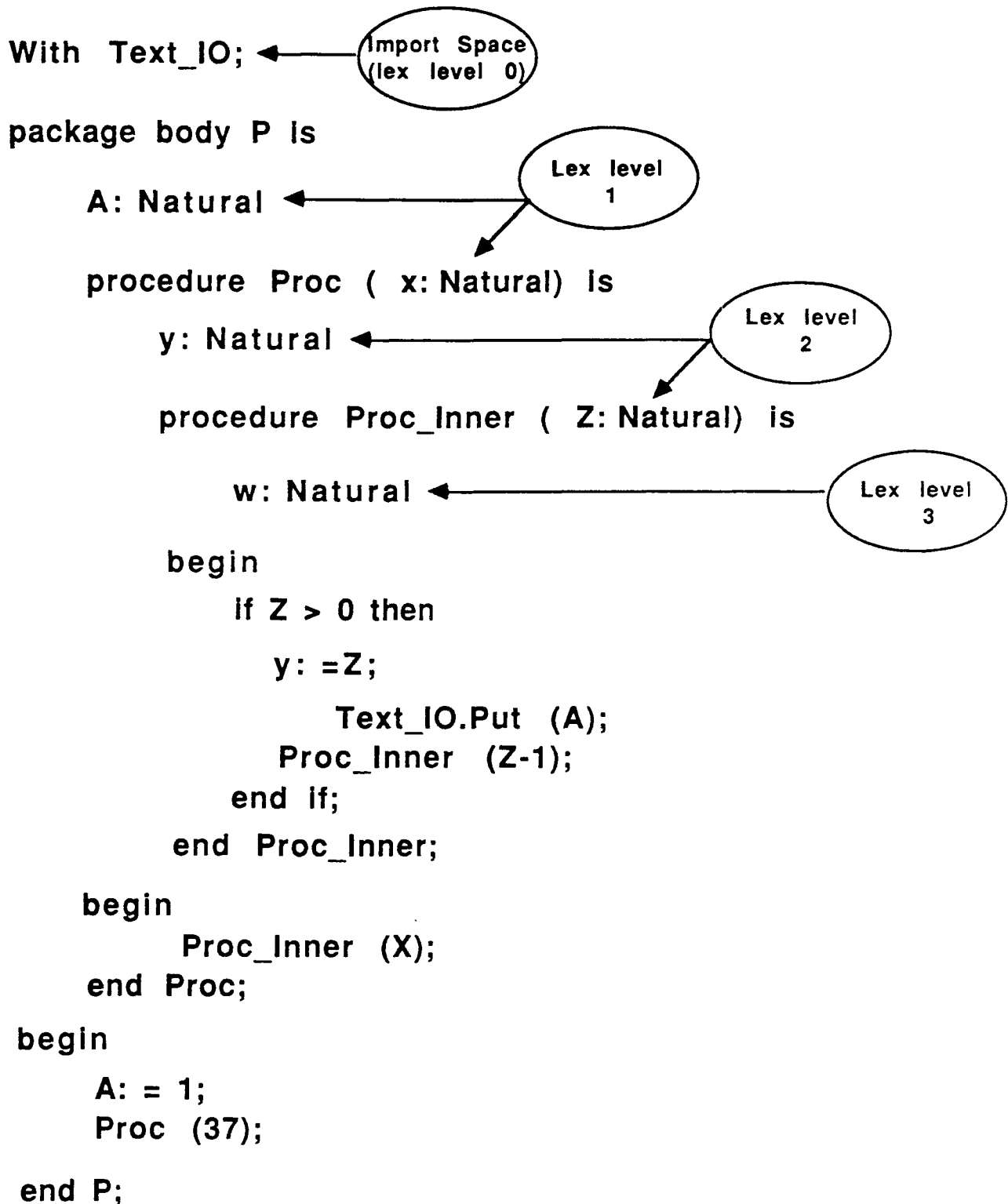
begin

    A: = 1;

    Proc (37);

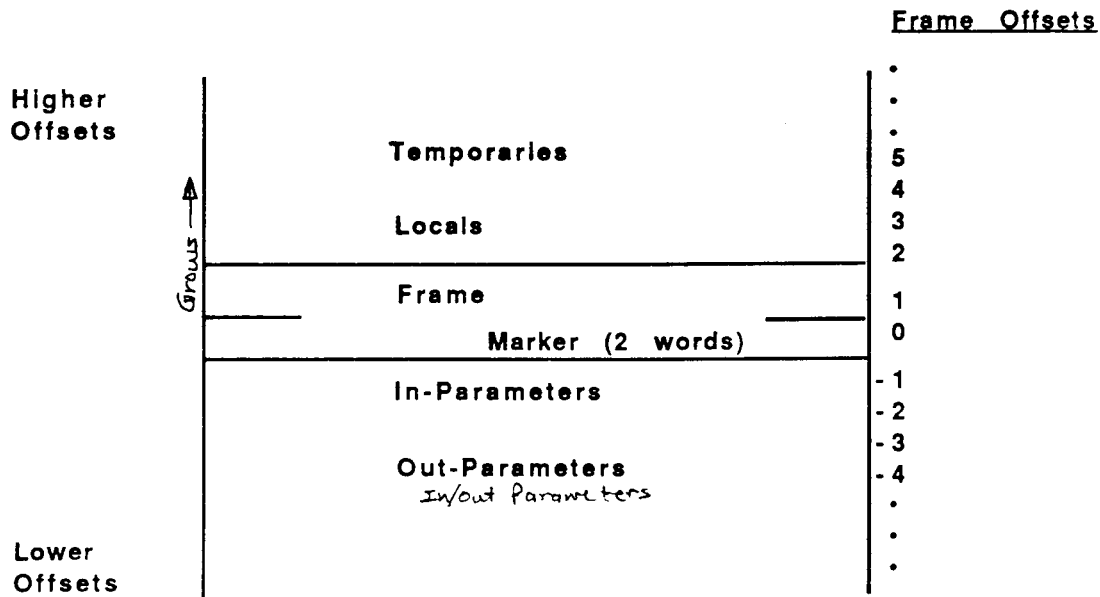
end P;

# EXECUTION ENVIRONMENT - ADA ADDRESSING

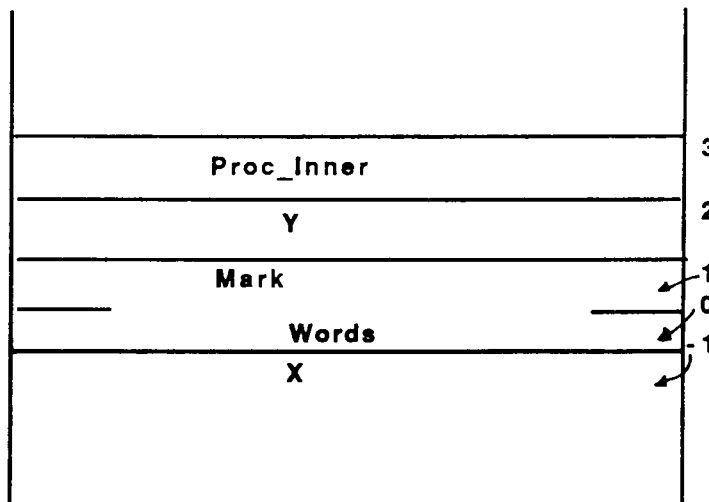


# EXECUTION ENVIRONMENT - ADA ADDRESSING (CONT'D.)

- Stack frame for a subprogram invocation.



- Stack frame for procedure Proc. Lex Level = 2



- Thus the call to Proc\_Inner is:

```
Load 2,-1 -- push X
Call 2, 3 -- invoke Proc_Inner
```

# EXECUTION ENVIRONMENT - ADA ADDRESSING (CONT'D.)

	<u>LEX LEVEL</u>	<u>OFFSET</u>	
P.A	1,	13	12 for TCB
P.Proc	1,	14	
Proc.X	2,	-1	
Proc.Y	2,	2	841 mark words
Proc.Proc_Inner.Z	3,	-1	below mark words
Proc.Proc_Inner.W	3,	2	
Text_IO	0, <i>import space</i>	1 (probably)	

R1φφφ

① use Display

② dynamic links

-once, then put into display



# EXECUTION ENVIRONMENT - ADA ADDRESSING (CONT'D.)

- **Lexical Level** is the static level of nesting of an object from its containing package.
- Machine addressing limits lex level to 4 bits - max procedure static nesting is 15 levels.
- Lex level 0 refers to import space.
- Each module is a closed scope. Can only address internal objects and objects imported to it by its parent.

*library subprograms  
get package to surround  
them (needs module)*

# EXECUTION ENVIRONMENT - FRAME INFORMATION

- **Return address**
- **Data and type stack marks**
- **"Outer frame" pointer - name of enclosing package**
- **Lex level**
- **Flag**
  - utility
  - rendezvous
- **Children**
  - tasks, packages

# EXECUTION ENVIRONMENT - RETURNING VALUES

- Out parameters are on the stack
- Function return is a control word left on the top-of-stack after return
- Data and type stack are popped to the pre-call values when the subprogram returns except:
  - \* - functions returning structured value: data stack not popped because return value is up there.
  - \* - Implications: Returning from a subprogram that consumes a lot of data stack space may not free that space. Loops with such calls could be trouble spots causing Storage\_Errors.
- The switch R1000\_CG.Reclaim\_Space cause the code generator to emit code in each loop to mark and pop the data stack to avoid run-away allocation. This switch defaults to true and is invisible unless set to false.

*doesn't appear in switch file  
- can't be edited until set false*
- This applies only to loops.

# EXECUTION ENVIRONMENT - RETURNING VALUES (CONT'D.)

- Thus, In

(A)       for i in 1..5 loop  
              F (G(i));  
          end loop;

and

(B)       F(G(1));  
          F(G(2));  
          F(G(3));  
          F(G(4));  
          F(G(5));

If G is a function returning a structure, (B) consumes 5 times the data stack space as (A).

# EXECUTION ENVIRONMENT - DEBUGGER

- **A number of debugger operations are available to display the execution environment:**

<b>Enable (Addresses)</b>	<b>causes display of machine information by a number of commands</b>
<b>Memory_Display</b>	<b>shows contents of the different kinds of memory segments</b>
<b>Stack</b>	<b>shows a module's call stack with frame addresses and information</b>
<b>Information (Space)</b>	<b>shows sizes of various memory segments</b>
<b>Task_Display</b>	<b>shows state of various tasks in the program</b>

*should follow section*

# **EXECUTION ENVIRONMENT - REVIEW**

- **Control stack frames**
- **Lex level, offset addressing**
- **Import spaces**
- **Function returns and data stack**
- **Debugger and kernel commands**
- **Package addressing**

# EXECUTION ENVIRONMENT - KERNEL COMMAND INTERPRETER

<b>Show_Task_States</b>	<b>displays task names and state based on various queries</b>
<b>LMR</b>	<b>Logical Memory Read. Hex display of memory segments. Especially useful if debugger is not runnable.</b>
<b>LMW</b>	<b>Logical Memory Write</b>
<b>Jobs</b>	<b>show information about jobs. Covered more later.</b>
<b>Job</b>	<b>Same as Jobs, but for one job</b>

# EXECUTION ENVIRONMENT - OTHER COMMANDS

**Show\_Tasks**

**displays stack dumps of  
specified task or jobs.  
More later.**



# EXECUTION ENVIRONMENT - PACKAGE ADDRESSING

package P is

Type T is private;

x: Natural;

Subtype S is integer;

function F returns T;

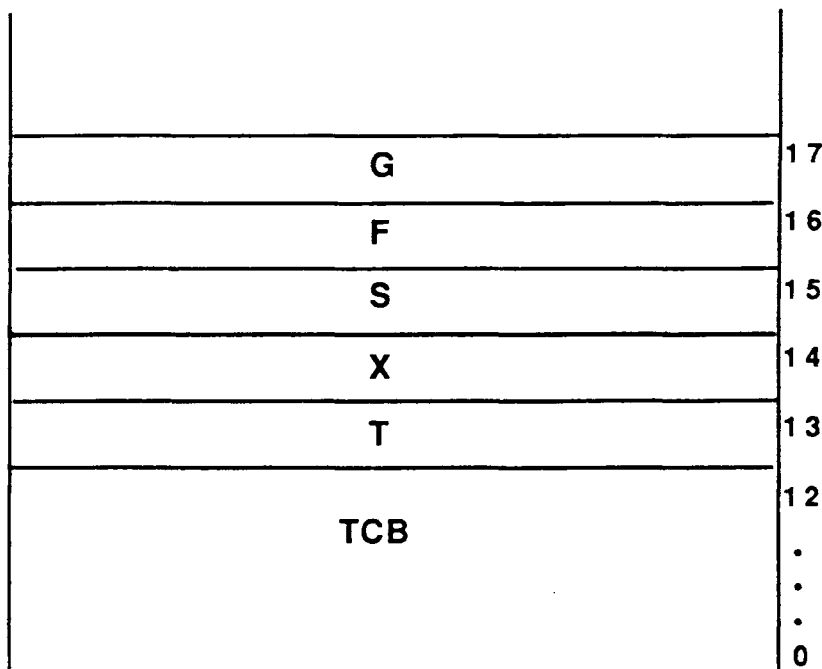
package IO renames Text\_IO; *- no offset for package renames*

function G renames F; *- generates offset*

private

. . .

end P;



# EXECUTION ENVIRONMENT - PACKAGE ADDRESSING (CONT'D.)

to get P.X from outside P

LOAD 0,4 -- suppose P is at import space  
offset 4

Execute Package, Field\_Read, 14 -- X

Similarly to call P.F

LOAD 0,4

Execute Package, Field\_Execute, 16

- Each visible object has an offset in the module.
- Visible objects are marked specially as visible.
- Package renames occupy no runtime offset.
- Some constants also occupy no runtime offset.

REVIEW

# ADA FEATURES - EXCEPTIONS

- Exceptions are represented by a 48 bit number.
- An exception can be raised explicitly by loading the exception number and executing:

**Execute, Exception, Raise**

- An exception can be raised implicitly when some Ada rule is violated. This is built into the R1000 instruction set. Checks are done in parallel with the operation.
- Exception propagation is handled largely by the microcode. The header of the subprogram points to the exception handler. If there is no handler, it points to location 4 in the code segment. *— instruction there will propagate the exception handler*
- Once raised, an exception variable is pushed on the stack. This control word contains the exception number and raise PC.
- Functions are available in the environment to get an image of the exception variable

**Debug\_Tools.Get\_Exception\_Name**

returns a string. Parameters control display of raise PC and machine representation of the exception number.

# ADA FEATURES - EXCEPTIONS (CONT'D)

- **Some values of the exception number are pre-defined. Others are encodings of the Diana pointer to the exception declaration.**
- **Some values are grouped into a contiguous range that represents different flavors of the same exception. This matters only in the image function which displays the flavor:**

**Constraint\_Error            (type range)**

**Constraint\_Error            (null access)**

**etc.**

- **See handout for ranges.**

```

0 => "Unknown (0)";
1 => "Constraint_Error (Type Range)";
2 => "Constraint_Error (Case Range)";
3 => "Constraint_Error (Exponent)";
4 => "Constraint_Error (Null Access)";
5 => "Constraint_Error (Array Index)";
6 => "Constraint_Error (Length Error)";
7 => "Constraint_Error (Discriminant)";
8 => "Constraint_Error (Variant)";
9 => "Constraint_Error (Entry Family)";
10 .. 31 => "Constraint_Error";
32 => "Numeric_Error (zero divide)";
33 => "Numeric_Error (overflow)";
34 .. 47 => "Numeric_Error";
48 => "Program_Error (elaboration order)";
49 => "Program_Error (function exit)";
50 => "Program_Error (select)";
51 => "Program_Error (prompt executed)";
52 .. 63 => "Program_Error";
64 => "Storage_Error (control)";
65 => "Storage_Error (type)";
66 => "Storage_Error (data)";
67 => "Storage_Error (import)";
68 => "Storage_Error (queue)";
69 => "Storage_Error (program)";
70 => "Storage_Error (allocation)";
71 => "Storage_Error (job page limit)";
72 => "Storage_Error (oversize object)";
73 => "Storage_Error (name)";
74 .. 79 => "Storage_Error";
80 => "Tasking_Error (activation)";
81 => "Tasking_Error (completed task)";
82 => "Tasking_Error (abnormal task)";
83 .. 95 => "Tasking_Error";
96 => "!Lrm.System.Operand_Class_Error";
97 => "!Lrm.System.Type_Error";
98 => "!Lrm.System.Visibility_Error";
99 => "!Lrm.System.Capability_Error";
100 => "!Lrm.System.Machine_Restriction";
101 => "!Lrm.System.Illegal_Instruction";
102 => "!Lrm.System.Illegal_Reference";
103 => "!Lrm.System.Illegal_Frame_Exit";
104 => "!Lrm.System.Record_Field_Error";
105 => "!Lrm.System.Utility_Error";
106 => "!Lrm.System.Unsupported_Feature";
107 => "!Lrm.System.Illegal_Heap_Access";
108 => "!Lrm.System.Select_Use_Error";
109 .. 127 => "Instruction_Error";
128 => "Task_Aborted";
129 => "!Lrm.System.Frame_Establish_Error";
130 => "Unimplemented_Microcode_Error";
131 => "!Lrm.System.Nonexistent_Space_Error";
132 => "!Lrm.System.Nonexistent_Page_Error";
133 => "!Lrm.System.Write_To_Read_Only_Page";
134 => "!Lrm.System.Heap_Pointer_Copy_Error";
135 => "!Lrm.System.Assertion_Error";
136 => "!Lrm.System.Microcode_Assist_Error";
137 .. 223 => not used
224 => "Constraint_Error";
225 => "Numeric_Error";

```

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```

226 => "Program_Error";
227 => "Storage_Error";
228 => "Tasking_Error";
229 .. 255 => not used
256 => "!Io.Io_Exceptions.Status_Error";
257 => "!Io.Io_Exceptions.Status_Error (already open)";
258 => "!Io.Io_Exceptions.Status_Error (not open)";
259 .. 271 => "Io_Exceptions.Status_Error";
272 => "!Io.Io_Exceptions.Mode_Error";
273 => "!Io.Io_Exceptions.Mode_Error (illegal on infile)";
274 => "!Io.Io_Exceptions.Mode_Error (illegal on outfile)";
275 .. 287 => "!Io.Io_Exceptions.Mode_Error";
288 => "!Io.Io_Exceptions.Name_Error";
289 => "!Io.Io_Exceptions.Name_Error (illformed name)";
290 => "!Io.Io_Exceptions.Name_Error (nonexistent directory)";
291 => "!Io.Io_Exceptions.Name_Error (nonexistent object)";
292 => "!Io.Io_Exceptions.Name_Error (nonexistent version)";
293 => "!Io.Io_Exceptions.Name_Error (ambiguous name)";
294 .. 303 => "!Io.Io_Exceptions.Name_Error";
304 => "!Io.Io_Exceptions.Use_Error";
305 => "!Io.Io_Exceptions.Use_Error (access restriction)";
306 => "!Io.Io_Exceptions.Use_Error (not checked out)";
307 => "!Io.Io_Exceptions.Use_Error (wrong class)";
308 => "!Io.Io_Exceptions.Use_Error (object frozen)";
309 => "!Io.Io_Exceptions.Use_Error (lock error)";
310 => "!Io.Io_Exceptions.Use_Error (file/device full)";
311 => "!Io.Io_Exceptions.Use_Error (line/page length exceeded)";
312 => "!Io.Io_Exceptions.Use_Error (reset failed)";
313 => "!Io.Io_Exceptions.Use_Error (operation unsupported)";
314 .. 319 => "!Io.Io_Exceptions.Use_Error";
320 => "!Io.Io_Exceptions.Device_Error";
321 => "!Io.Io_Exceptions.Device_Error (nonexistent page)";
322 => "!Io.Io_Exceptions.Device_Error (write to read-only page)";
323 => "!Io.Io_Exceptions.Device_Error (illegal reference)";
324 => "!Io.Io_Exceptions.Device_Error (illegal heap access)";
325 => "!Io.Io_Exceptions.Device_Error (device data error)";
326 .. 335 => "!Io.Io_Exceptions.Use_Error";
336 .. 351 => "!Io.Io_Exceptions.End_Error";
352 => "!Io.Io_Exceptions.Data_Error";
353 => "!Io.Io_Exceptions.Data_Error (bad input syntax)";
354 => "!Io.Io_Exceptions.Data_Error (bad input value)";
355 => "!Io.Io_Exceptions.Data_Error (bad output value)";
356 => "!Io.Io_Exceptions.Data_Error (bad output type)";
357 .. 367 => "!Io.Io_Exceptions.Data_Error";
368 => "!Io.Io_Exceptions.Layout_Error";
369 => "!Io.Io_Exceptions.Layout_Error (bad column)";
370 => "!Io.Io_Exceptions.Layout_Error (bad position)";
371 => "!Io.Io_Exceptions.Layout_Error (bad item length)";
372 .. 383 => "!Io.Io_Exceptions.Layout_Error";

```

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```

package Debug_Tools is

  procedure Debug_On;
  procedure Debug_Off;

  -- Enable or disable debugging for the calling task's job. When enabled,
  -- only tasks that are descendents of the caller can be debugged.
  -- When debugging is disabled, the task is released to execute, and all
  -- active debugger "hooks" are deactivated (eg, breakpoints, etc).

  function Debugging return Boolean;

  -- return true if calling task is being debugged.

  procedure Message (Info : String);

  -- Print the message string in the debugger window. No operation if
  -- the debugger is not activated

  procedure User_Break (Info : String);

  -- "Break" in the debugger. The calling task stops as though it
  -- encountered a breakpoint. If the debugger is not active, no action
  -- is performed. Otherwise, the task remains stopped until the
  -- debugger user explicitly continues its execution.

  procedure Set_Task_Name (Name : String);
  function Get_Task_Name return String;

  -- Set or retrieve a string "synonym" for the calling task. This name
  -- is used within the debugger to make identifying task easier.
  -- It is also useful for multiple instances of the same task type
  -- to distinguish themselves in the debugger.
  -- No operation if the debugger is not activate.

  function Ada_Location (Frame_Number : Natural := 0;
                        Fully_Qualify : Boolean := True;
                        Machine_Info : Boolean := False) return String;

  -- Return a string name for the Ada location of execution in the
  -- specified stack frame. Frame_Number = 0 refers to the caller
  -- of Ada_Location. Frame_Number = 1 refers to its caller, and so
  -- on. The null string is returned if the frame is nonexistent or
  -- its location cannot be found for some other serious reason.

  -- This procedure works independent of whether there is an active
  -- debugger for the calling tasks, but it may return less information
  -- if there is not.

  function Get_Exception_Name (Fully_Qualify : Boolean := True;
                               Machine_Info : Boolean := False) return String;

  -- return a string representation of the exception most recently
  -- executed by the calling task. Get_Exception_Name must be called
  -- either directly or indirectly from an exception handler. If
  -- no exception is found, the null string is returned.

  function Get_Raise_Location (Fully_Qualify : Boolean := True;

```

```

                               Machine_Info : Boolean := False) return String;

  -- return a string representation of the location of the exception
  -- most recently executed by the calling task. Get_Raise_Location
  -- must be called either directly or indirectly from an exception
  -- handler. If no exception is found or other problems encountered,
  -- the null string is returned.

  generic
    type T is limited private;
    -- Type that will be displayed using Image procedure as part
    -- of debugger's object display procedure.

    with function Image (Value : T;
                        Level : Natural;
                        Prefix : String;
                        Expand_Pointers : Boolean) return String;

    -- Given the Value, return its image. Level specifies the number
    -- of levels to detail to be displayed. Expand_Pointers indicates
    -- that internal pointers should be expanded in the image.
    -- Level = 0 => entire object should be elided.

    -- If any ASCII.LF's are emitted, the following line should
    -- start with Prefix as an "indent" value.

  procedure Register;
  -- Make this special display procedure known to the debugger
  -- of the session making the call.

  generic
    type T is limited private;
  procedure Un_Register;
  -- Remove the display procedure from the calling session's
  -- database of special types for the type T given.

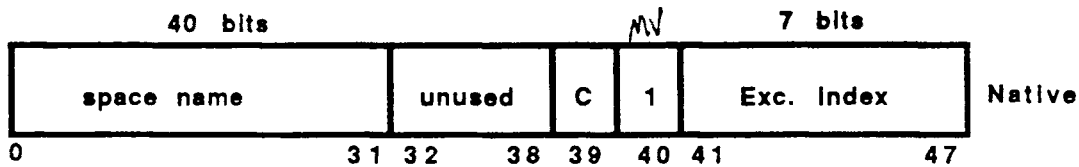
  pragma Subsystem (Native_Debugger);
  pragma Module_Name (4, 3802);

end Debug_Tools;

```

# ADA FEATURES - EXCEPTIONS (CONT'D.)

- Exception numbers representing native exceptions are interpreted differently from those representing exception in cross-compiled (MV) code.



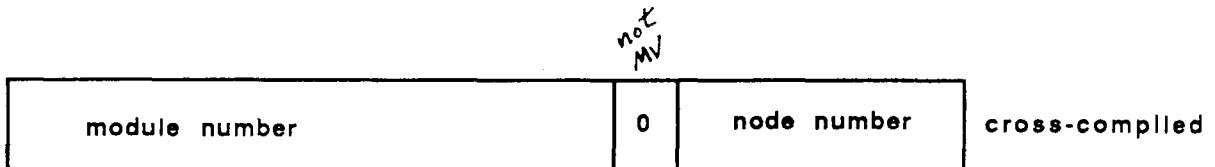
*R18815  
Native*

<Space = 19230979, Index = 2>

*- included bit at 40?*

*MV*

<Module = 3174, Ord = 46>



C - Is\_Code\_Archived bit

# ADA FEATURES - EXCEPTIONS

- **Getting exception names**
  - **Cross-compiled exceptions: exceptions that are defined in the universe are built into a special table and the names of these are, thus, known. When the name of another cross-compiled exception is requested, the MV-cross debugger is called if one is connected. If not connected, or the answer takes more than 20 seconds, the numeric form**  
  
**<module = n, ord = n >**  
  
is returned.
  - **Native code exceptions: since the Diana pointer is encoded into the exception number, the tree can be accessed to get the exception name if the tree exists. Cases where it doesn't:**
    - \* loaded main programs
    - \* subsystem code views
    - \* exceptions defined in command windows



# ADA FEATURES - EXCEPTIONS

- **Getting exception information from the debugger**

**Information (Exceptions)**

- **If addresses is enabled, shows control offset of exception variable. Use Memory\_Display to look at this - It give the true information about the exception.**
- **When an exception is raised in the environment, a loaded main subprogram, or code view, the debugger provides the view that the exception was actually raised at the statement in Native code that called whoever actually raised the exception.**
- **The debugger keeps this information in a cache that can be displayed with:**

**Debug.Show (Exception\_Cache)**

- **In obscure cases, the cache information can be lost and the debugger will be vague about the raise location.**
- **If putting in a problem report about an environment exception flying out, It is helpful to include the true raise address if the problem is not easily reproducible.**

# ADA FEATURES - UNCHECKED CONVERSION

- **Ada generic function** (*of unchecked conversion*)
  - checks type for legality on every call. This makes it slow.
- **Package Unchecked\_Conversions**
  - package Unchecked\_Conversion\_Package checks type for legality when instantiation is elaborated. Faster, but requires a module (several pages) for each instantiation.
  - conversions to/from byte string; good performance and very useful

*Basic Rule*

- The bits you put in are the bits you get out.
- See handout

with System;

package Unchecked\_Conversions is

generic

type Source is limited private;

type Target is limited private;

package Unchecked\_Conversion\_Package is

function Convert (S : Source) return Target;

-- Package form of LRM Unchecked\_Conversion.

-- Type-specific calculations are made during package elaboration, making

-- calls to this convert faster than to an equivalent Unchecked\_Conversi\

on

-- instantiation. Speed improvement depends on the type involved.

end Unchecked\_Conversion\_Package;

generic

type Source is limited private;

function Convert\_To\_Byte\_String (S : Source) return System.Byte\_String;

-- Convert from Source to a byte string. The byte string may contain

-- more bits than the object.

generic

type Target is limited private;

function Convert\_From\_Byte\_String (S : System.Byte\_String) return Target;

-- Convert from a byte string to a Target type. The string should

-- have been produced by an instantiation of Convert\_To\_Byte\_String

-- with the same type. A constrained object is always returned.

pragma Subsystem (Miscellaneous);

pragma Module\_Name (4, 3543);

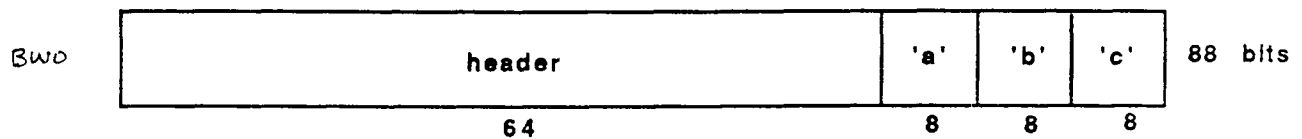
end Unchecked\_Conversions;

# ADA FEATURES - UNCHECKED CONVERSION (CONT'D.)

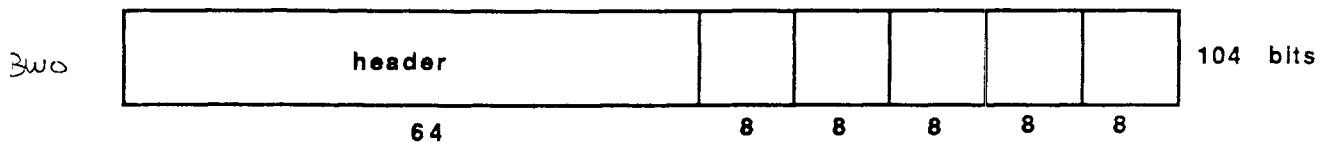
BWO Bounds with Object  
BWT Bounds with Type

- Some examples of object layout :

String constant: "abc"



S: string (1..5)



Subtype Fixed\_String is string (2..4);

S2: Fixed\_String



```

with Display; -- a simple generic that dumps an object in hex.
with System, Io;
with Unchecked_Conversion; -- the slow unchecked conversion

procedure Array_Example is

  type Block is array (0 .. 7) of System.Byte; -- This is the dest type
  Test_Block : Block;

  S : String (1 .. 5) := "abcde";
  subtype Fixed_String is String (2 .. 4);
  S2 : Fixed_String := "123";

  generic
    type E is private;
    type A is array (Positive range <>) of E;
  function Convert (Source : A) return Block;
  -- This generic converts "any" array type into a block. Raises
  -- nasty exceptions of the data doesn't fit. Only actual data is
  -- stored, not descriptor bits.

  function Convert (Source : A) return Block is
    Size : Natural := Source'Length * E'Size; -- Number of data bits

    type Bit_Vector is array (Integer range <>) of Boolean;
    subtype Bit_Vect is Bit_Vector (1 .. 1024); -- bit image of Source
    subtype Block_Bits is Bit_Vector (1 .. Block'Size); -- of Block

    Bits : Bit_Vect; -- Gets binary image of source object
    Real_Bits : Block_Bits; -- Gets binary image of value to return
    Data_Offset : Natural; -- Offset in Bits of start of real data

    function C1 is new Unchecked_Conversion (A, Bit_Vect);
    function C2 is new Unchecked_Conversion (Block_Bits, Block);
  begin
    Io.Put_Line ("source size = " & Natural'Image (Size));
    Io.Put_Line ("source obj size = " & Natural'Image (Source'Size));

    -- Form binary image of source parameter
    Bits := C1 (Source);

    -- Extract real data bits
    Data_Offset := Source'Size - Size + 1;
    Real_Bits (1 .. Size) := Bits (Data_Offset .. Source'Size);

    -- return final value
    return C2 (Real_Bits);
  end Convert;

  function Cvt is new Convert (Character, String);
  procedure D is new Display (Block);
begin
  Io.Put_Line ("String constant");
  Test_Block := Cvt ("abc");
  D (Test_Block);
  Io.New_Line;

  Io.Put_Line ("String variable");
  Test_Block := Cvt (S);

```

```

  D (Test_Block);
  Io.New_Line;

  Io.Put_Line ("String expression");
  Test_Block := Cvt ("A" & S & "A");
  D (Test_Block);
  Io.New_Line;

  Io.Put_Line ("Constrained string variable");
  Test_Block := Cvt (S2);
  D (Test_Block);
end Array_Example;

```

Output from this program:

```

String constant
source size = 24
source obj size = 88
Object size = 64 bits.
Object fits in 8 bytes.
616263AB0829AA09
Val:

```

```

String variable
source size = 40
source obj size = 104
Object size = 64 bits.
Object fits in 8 bytes.
6162636465090128

```

```

String expression
source size = 56
source obj size = 120
Object size = 64 bits.
Object fits in 8 bytes.
4161626364654138

```

```

Constrained string variable
source size = 24
source obj size = 24
Object size = 64 bits.
Object fits in 8 bytes.
31323307278F6374

```

```

!USERS.PHIL.GURU COURSE.UNCHECKED CONVERSION EXAMPLE'V(2)
with Unchecked_Conversion; -- Language-defined function from !LRM
with Unchecked_Conversions; -- package form from !tools
with String_Utilities, Io, System;

procedure Conversion is

  type Vstring (Max_Length : Positive) is
    record
      Length : Positive;
      Contents : String (1 .. Max_Length);
    end record;

  subtype V20 is Vstring (20); -- A constrained subtype.
  type Bits is array (0 .. 400) of Boolean;

  S : V20; -- We will be converting between these two objects
  B : Bits;

  function Convert1 is new Unchecked_Conversion (V20, Bits);
  package Convert2 is new Unchecked_Conversions.Unchecked_Conversion_Package
    (Bits, V20);

  generic
    type T is private;
  procedure Display (Object : T);
  -- This procedure produces a binary dump of the Object.
  procedure Display (Object : T) is
    function Convert_To_Bytes is
      new Unchecked_Conversions.Convert_To_Byte_String (Source => T);
    procedure Display_Bytes (Value : System.Byte_String) is
      Bytes_Output : Natural := 0;
    begin
      Io.Put_Line ("Object fits in " & Natural'Image (Value'Length) &
        " bytes.");
      for I in Value'First .. Value'Last loop
        Io.Put (String_Utilities.Number_To_String (Integer (Value (I)),
          Base => 16,
          Width => 2,
          Leading => '0'));

          Bytes_Output := Bytes_Output + 1;
          if Bytes_Output mod 32 = 0 then
            Io.New_Line;
          elsif Bytes_Output mod 8 = 0 then
            Io.Put (" ");
          end if;
        end loop;
      end Display_Bytes;
    begin
      Io.Put_Line ("Object size = " & Integer'Image (Object'Size) & " bits.");
      Display_Bytes (Convert_To_Bytes (Object));
      Io.New_Line;
    end Display;

  procedure Dump is new Display (V20); -- Instantiate dump for each
  procedure Dump is new Display (Bits); -- of the two types.

begin
  S.Length := 15; -- Initialize S

```

```

!USERS.PHIL.GURU COURSE.UNCHECKED CONVERSION EXAMPLE'V(2)
S.Contents (1 .. 15) := "000011112222abc";

-- Display initial object
Io.Put_Line ("initial VString object:");
Dump (S);

-- Convert and display result
Io.Put_Line ("converted bits object");
B := Convert1 (S);
Dump (B);

-- Now, modify the bits in
Io.Put_Line ("Modified original object");
B (256) := True; -- This changes one of the characters
S := Convert2.Convert (B); -- Convert bits back into record
Dump (S);
Io.Put_Line (S.Contents (1 .. 15));

end Conversion;

Generated output:

initial VString object:
Object size = 351 bits.
Object fits in 44 bytes.
800000140000001E 00000080000001C0 000000200000028 6060606062626262
64646464C2C4C600 00000000
converted bits object
Object size = 401 bits.
Object fits in 51 bytes.
800000140000001E 00000080000001C0 000000200000028 6060606062626262
64646464C2C4C600 0000000100000028 000000
Modified original object
Object size = 351 bits.
Object fits in 44 bytes.
800000140000001E 00000080000001C0 000000200000028 6060606062626262
E4646464C2C4C600 00000001
00001111r222abc

```

```
generic
  type T is private;
  procedure Display (Object : T);
  - This procedure produces a binary dump of the Object.

with Unchecked_Conversions;
with Io;
with System;
with String_Uilities;

procedure Display (Object : T) is
  function Convert_To_Bytes is
    new Unchecked_Conversions.Convert_To_Byte_String (Source => T);

  procedure Display_Bytes (Value : System.Byte_String) is
    Bytes_Output : Natural := 0;
  begin
    Io.Put_Line ("Object fits in " &
      Natural'Image (Value'Length) & " bytes.");
    for I in Value'First .. Value'Last loop
      Io.Put (String_Uilities.Number_To_String (Integer (Value (I)),
        Base => 16,
        Width => 2,
        Leading => '0'));

      Bytes_Output := Bytes_Output + 1;
      if Bytes_Output mod 32 = 0 then
        Io.New_Line;
      elsif Bytes_Output mod 8 = 0 then
        Io.Put (" ");
      end if;
    end loop;
  end Display_Bytes;

begin
  Io.Put_Line ("Object size = " & Integer'Image (Object'Size) & " bits.");
  Display_Bytes (Convert_To_Bytes (Object));
  Io.New_Line;

end Display;
```

```

with System;
with Unchecked_Conversion;
with Text_Io;
with Display;

procedure Unch is
  Base_B : constant := 31;
  Base : constant Long_Integer := 2 ** Base_B;
  subtype Universal_Digit is Long_Integer range -(Base - 1) .. (Base - 1);
  type Vector is array (Positive range <>) of Universal_Digit;
  X : Vector (1 .. 1) := (1 => 6);
  subtype Int is Vector (1 .. X'Length);
  Y : Int := X; -- make sure we have a real object of type Int

  subtype Str is System.Byte_String
    (1 .. X'Length * Universal_Digit'Size / System.Byte'Size);

  function Xbytes is new Unchecked_Conversion (Int, Str);
  Result : constant System.Byte_String := Str'(Xbytes (X));
  Result2 : constant System.Byte_String := Str'(Xbytes (Y));
  Result3 : constant System.Byte_String := Str'(Xbytes (X));

  procedure Dump is new Display (Int);
  procedure Dump is new Display (Str);

  procedure Show (Result : Str) is
  begin
    Text_Io.Put_Line ("result'first = " & Integer'Image (Result'First));
    Text_Io.Put_Line ("result'last = " & Integer'Image (Result'Last));
    for I in Result'Range loop
      Text_Io.Put (System.Byte'Image (Result (I)) & " ");
    end loop;
  end Show;
begin
  Dump (X);
  Dump (Y);
  Dump (Result);
  Show (Result);
  Dump (Result2);
  Show (Result2);
  Dump (Result3);
  Show (Result3);
end Unch;

```

Output from this program:

```

Object size = 96 bits.      -- Dump of X : Vector (1..1) := (1=>6)
Object fits in 12 bytes.
0000000100000001 00000006

Object size = 32 bits.     -- Dump of Y : Int := X; -- constrained
Object fits in 4 bytes.
00000006

Object size = 32 bits.     -- Result : conversion of Int' (X)

```

```

Object fits in 4 bytes.
00000001                -- ended up with array bounds bits!
result'first = 1
result'last = 4
0 0 0 1

Object size = 32 bits.   -- Result2 : conversion of Y
Object fits in 4 bytes.
00000006                -- correct answer
result'first = 1
result'last = 4
0 0 0 6

Object size = 32 bits.   -- Result3 : conversion of X
Object fits in 4 bytes.
00000001                -- gets array bounds bits
result'first = 1
result'last = 4
0 0 0 1

```



# UNCHECKED - DEALLOCATION

- Normally, calls to Instantiations of Unchecked\_Deallocation do not reclaim storage.
- For certain access types deallocation can be enabled and Unchecked\_Deallocation will actually reclaim storage.
- When deallocation is enabled, storage requirements in the collection are increased.
  - small amount on (a per page basis) for headers
  - twice the 'size of the access type per allocated object (used to maintain the free list).
- Enabling deallocation for an access type, enables deallocation for all of its subtypes and all types derived from it (since they all share the same collection).

*see next page*

# ENABLING UNCHECKED\_ DEALLOCATION

- **Pragma Enable\_Deallocation (Access\_Type\_Name)**
  - Enables deallocation, if possible, only for the named access type. Warnings are produced when the unit is coded if the compiler can determine that deallocation is not possible for this type.
  - This pragma may be applied to generic formal access types and the compiler will produce warnings for instantiations that provide actual access types for which deallocation is enabled.
- **Library switch R1000\_Cg.Enable\_Deallocation:=True**
  - Enables deallocation for all access types for which deallocation is possible. Deallocation may be disabled for specific access types using pragma **Disable\_Deallocation (Access\_Type\_Name)**.

# ALLOCATION STRATEGY

- **First fit**
- **Items that are larger than required are split, if possible (remaining fragment must be big enough).**
- **Deallocation returns items to the head of the free list.**

# RESTRICTIONS ON UNCHECKED\_ DEALLOCATION

- Designated type may not be a task or contain task components
- Designated type may not be access to task or contain access to task components
- Access type may not be a segmented heap pointer
- Designated type may not be a segmented heap pointer or contain segmented heap pointers
- The compiler cannot always detect these cases at compile time to produce warnings. To determine if Unchecked\_Deallocation will actually reclaim storage, instantiate the following generic in !Tools:

```
Generic
    type Object is limited private;
    type Name is access object;
function Allows_Deallocation return Boolean;
```

# PROGRAM STRUCTURE

- **When a command is run or main program is built, information on the entire program is gathered.**
  - the unit state of all units is checked
  - the code for all units is located
  - an elaboration order for library units is computed
  - specs in spec views are matched with bodies in load views using the current activity
  - an elaboration module is generated
  
- **The elaboration module is like a package that contains all units in the program. It:**
  - 1) Declares each package type
  - 2) Sets-up import spaces for the modules
  - 3) Creates instances for each package
  - 4) Elaborates packages in order
  - 5) Calls the command procedure

# **PROGRAM STRUCTURE - ELABORATION SEGMENT**

- **When a main subprogram (pragma main) is promoted from installed to coded, its elaboration segment is built.**
- **This will take a long time if the closure is large. Thus, one shouldn't put code likely to change in such a unit.**
- **There is no Ada source, or Diana tree for the elaboration segment, and a minimal debug table.**
- **Occasionally an exception will be raised in an elaboration segment. It is difficult to pinpoint this with the debugger because there is no source. One must take the PC and look at the code to see if it is an elaboration segment. The segment listing can be helpful here.**
- **A special attribute called "units in program" is attached to the elaboration module. It lists all the Ada units in the program closure. See !Compiler\_Interface.Units\_In\_Program.**

Code for segment 1035520

```
0 0000: 000F
1 0001: 5800
2 0002: 0000
3 0003: 0068
4 0004: 0100
5 0005: 00BB
6 0006: 0048
7 0007: 0000
```

```
8 0008: 000E BLOCK_BEGIN 1, 6
9 0009: 0004 BLOCK_HANDLER 0, 4
10 000A: 0001 END_LOCALS 1
```

```
::: module Main
11 000B: 029C DECLARE_SUBPROGRAM FOR_OUTER_CALL, IS_VISIBLE
12 000C: 0013 EXTENSION 2, 3
13 000D: 00BF ACTION ACCEPT_ACTIVATION
14 000E: 00BC ACTION SIGNAL_ACTIVATED
15 000F: 00BB ACTION SIGNAL_COMPLETION
```

#9

```

16 0010: 0015 BLOCK_BEGIN      2, 5
17 0011: 001A BLOCK_HANDLER    3, 2
18 0012: 0002 END_LOCALS      2

;;; subprogram ARCHITECTURE_EXAMPLE
19 0013: 029F DECLARE_SUBPROGRAM FOR_CALL
20 0014: 0023 EXTENSION        4, 3
21 0015: 8402 CALL              2, 2
22 0016: 7801 JUMP              1
23 0017: 0003 EXTENSION        3
24 0018: 006F ACTION            BREAK_UNCONDITIONAL
25 0019: 4501 EXIT_SUBPROGRAM   1
26 001A: 7801 JUMP              1
27 001B: 0003 EXTENSION        3
28 001C: 006F ACTION            BREAK_UNCONDITIONAL
29 001D: 00D8 LOAD_TOP          0      ;;; ( 2, 3)
30 001E: 0100 EXECUTE           EXCEPTION_CLASS, RAISE_OP
31 001F: 0000

```

```

32 0020: 004A BLOCK_BEGIN      9, 2
33 0021: 0004 BLOCK_HANDLER    0, 4
34 0022: 0006 END_LOCALS      6

;;; subprogram ARCHITECTURE_EXAMPLE
35 0023: 7801 JUMP              1
36 0024: 0002 EXTENSION        2
37 0025: 006F ACTION            BREAK_UNCONDITIONAL
38 0026: 6039 INDIRECT_LITERAL DISCRETE_CLASS, 57      ;;; page limit -
                                                    actual value is patched
39 0027: 6034 INDIRECT_LITERAL DISCRETE_CLASS, 52      ;;; module name of
                                                    (Runtime_Uutilities)

40 0028: E002 LOAD              0, 2      ;;; Proto Package Type
41 0029: E001 LOAD              0, 1      ;;; Ucode Assist Package
42 002A: 1835 EXECUTE           PACKAGE_CLASS, FIELD_EXECUTE_OP, 53
43 002B: 180D EXECUTE           PACKAGE_CLASS, FIELD_EXECUTE_OP, 13
44 002C: 4800 SHORT_LITERAL    0
45 002D: 602A INDIRECT_LITERAL DISCRETE_CLASS, 42
46 002E: 02A0 DECLARE_SUBPROGRAM NULL_SUBPROGRAM
47 002F: 038E DECLARE_TYPE     PACKAGE_CLASS, DEFINED      ;;;
                                                    ARCHITECTURE_EXAMPLE

48 0030: 0387 DECLARE_VARIABLE PACKAGE_CLASS      ;;; ARCHITECTURE_EXAMPLE
49 0031: 00D8 LOAD_TOP          0
50 0032: 190D EXECUTE           PACKAGE_CLASS, FIELD_REFERENCE_OP, 13      ;;; Local
                                                    copy of lib unit subprogram ARCHITECTURE_EXAMPLE
51 0033: 6020 INDIRECT_LITERAL DISCRETE_CLASS, 32      ;;; module name of
                                                    IO

52 0034: E002 LOAD              0, 2      ;;; Proto Package Type
53 0035: E001 LOAD              0, 1      ;;; Ucode Assist Package
54 0036: 1835 EXECUTE           PACKAGE_CLASS, FIELD_EXECUTE_OP, 53
55 0037: 6018 INDIRECT_LITERAL DISCRETE_CLASS, 24      ;;; module name of
                                                    JOB_SEGMENT

56 0038: E002 LOAD              0, 2      ;;; Proto Package Type
57 0039: E001 LOAD              0, 1      ;;; Ucode Assist Package
58 003A: 1835 EXECUTE           PACKAGE_CLASS, FIELD_EXECUTE_OP, 53
59 003B: 6010 INDIRECT_LITERAL DISCRETE_CLASS, 16      ;;; module name of
                                                    DEFAULT

60 003C: E002 LOAD              0, 2      ;;; Proto Package Type
61 003D: E001 LOAD              0, 1      ;;; Ucode Assist Package
62 003E: 1835 EXECUTE           PACKAGE_CLASS, FIELD_EXECUTE_OP, 53
63 003F: 00D8 LOAD_TOP          0      ;;; ( 3, 6)      ;;; DEFAULT
64 0040: E605 LOAD              3, 5      ;;; JOB_SEGMENT
65 0041: E604 LOAD              3, 4      ;;; IO
66 0042: 4803 SHORT_LITERAL    3
67 0043: E602 LOAD              3, 2      ;;; ARCHITECTURE_EXAMPLE
68 0044: 020E EXECUTE           MODULE_CLASS, AUGMENT_IMPORTS_OP
69 0045: E602 LOAD              3, 2      ;;; ARCHITECTURE_EXAMPLE
70 0046: 020F EXECUTE           MODULE_CLASS, ACTIVATE_OP
71 0047: 7802 JUMP              2      ;;; code segment reference table start
72 0048: 002D EXTENSION        5, 5
73 0049: 0000 EXTENSION        0
74 004A: 8603 CALL              3, 3      ;;; ARCHITECTURE_EXAMPLE
75 004B: 4501 EXIT_SUBPROGRAM   1
76 004C: 0000 LITERAL_VALUE    DISCRETE_CLASS, 931844
77 004D: 0000
78 004E: 000E
79 004F: 3804
80 0050: 0000 LITERAL_VALUE    DISCRETE_CLASS, 935940
81 0051: 0000

```



```
82 0052: 000E
83 0053: 4804
84 0054: 0000 LITERAL_VALUE DISCRETE_CLASS, 3590148
85 0055: 0000
86 0056: 0036
87 0057: C804
88 0058: 000F LITERAL_VALUE DISCRETE_CLASS, 1033472
89 0059: C500
90 005A: 0000
91 005B: 00B0
92 005C: 0000 LITERAL_VALUE DISCRETE_CLASS, 1641476
93 005D: 0000
94 005E: 0019
95 005F: 0C04
96 0060: 0000 LITERAL_VALUE DISCRETE_CLASS, 0
97 0061: 0000
98 0062: 0000
99 0063: 0000
100 0064: 0000
101 0065: 0000
102 0066: 0000
103 0067: 0000
```

```
;;; debug table start
104 0068: 0004 0000 5397 0000 0001 0001 000B 0000
112 0070: 0009 0000
```

# **PROGRAM STRUCTURE - LIBRARY UNITS**

- **Each library package is a child of the elaboration module.**
- **Library unit subprograms are wrapped in packages. Thus, each library unit is a module.**
- **Consequently, many library subprograms occupy far more pages than one package with many subprograms. There may be closure size issues with this approach.**

# PROGRAM STRUCTURE - MAIN UNITS

- **Main subprograms serve as roofs over entire programs. When a main subprogram is called, a new set of library units are elaborated.**

```
package P is ... end p
```

```
with P; procedure M is ... end m; pragma main;
```

- **A command:**

```
M;
```

```
M;
```

```
M;
```

**will elaborate 3 independent copies of package P, not one shared copy. If M were not a main subprogram, only one shared copy of P would be elaborated.**

*\*  
Beware of package state!*

- **In addition to the semantic Implications, this also has debugger implications since there can be more than one elaborated copy of a package in a program.**
- **If any unit in the closure of a main unit is uncoded, the main unit will be demoted as well.**

# PROGRAM STRUCTURE - LOADED MAIN UNITS

- A loaded main unit is built by the Compilation.Load command and is a single object containing a full copy of the entire closure of the main unit.
- It does not include any trees or source and cannot be debugged. *<unknown location>*
- It is independent of the units from which it was built - they can be demoted or deleted independent of the loaded main unit.

# PROGRAM STRUCTURE - PAGE LIMITS

*intended to catch runaway jobs*

- When a page limit pragma appears, the limit value for the job is set to that value if that value is bigger.
- Every unit has a limit associated with it. The value from the switch file is used if there is no explicit pragma.
- The largest limit is set before program elaboration begins.
- An explicit call in the program to set the limit will set the limit independent of any value imposed by pragmas or switches.

`System_Uilities.Set_Job_Page_Limit`

# PROGRAM STRUCTURE - PAGE LIMITS (CONT'D.)

## EXAMPLE 1

Session Default: 8000  
Library Default: 9000  
Unit A: Pragma Page\_Limit (7000)  
Unit B: No Pragma  
Unit C: Pragma Page\_Limit (10000)

Result: Elaboration begins  
with limit 10000

## EXAMPLE 2

Session Default: 8000  
Library Default: 12000  
Unit A: Pragma Page\_Limit (7000)  
Unit B: No Pragma  
Unit C: Pragma Page\_Limit(10000)

Result: Elaboration begins  
with limit 12000

## EXAMPLE 3

Session Default: 8000  
Library Default: 8000  
Unit A: Pragma Page\_Limit (9000)  
Unit B: Main subprogram with its  
limit 12000

Result: Elaboration begins with  
limit 9000. Limit raised  
to 12000 when B is called.  
(N.B. - This currently does not work.)

*Stays at 12000  
when B returns*

*4-Aug-1987*

# **PROGRAM STRUCTURE - REVIEW**

- **Elaboration modules**
- **Library units**
- **Setting-up import spaces**
- **Main unit semantics**
- **Loaded main subprograms**
- **Page limits**

# TASKING - BASICS

- Each task has a name. This is the segment name of the task's control, type, and data stacks.
- Recall that each package is a "task", too.
- Tasks have state and other information the system keeps about them.
- Task creation, activation, termination, rendezvous, etc., are handled by the R1000 microcode.
- Major steps in a task:
  - creation of the task type
  - creation of the task
  - activation of the task - elaboration of declarative part
    - \* occurs when "begin" reached; Tasking\_Error if exception during activation; or when allocated
  - execution of the task; interactions with other tasks; delays
  - completion of the task                      -- T' callable = false
  - termination of the task                    -- T' terminated = true
  - aborts make tasks "abnormal"

*at next synchronization point, task will terminate*



# TASKING - TASK CONTROL BLOCK

- The first 13 words of each control stack are the Task Control Block (TCB).
- This has lots of state information concerning the task.
- See handout for format.
- Many debugger and diagnostic commands extract information from the TCB
- Summary

## Word

0	Outer block mark words
1	" " " "
* 2	Current PC, TOS, priority
3	Data TOS, Type TOS
* 4	Microcode temp storage
* 5	" " " "
* 6	Task type, import space, and declarer
7	Task dependence information
8	Size of control stack, break mask
9	Size of type, queue, and data spaces
A	Delay, scheduler, and debugger information
B	Debugger interface subprogram
C	Microcode temp storage for slice information

- \* These tend to be the most interesting during problem diagnosis

MAR and TAGSTORE

0	8	16	24	32	40	48	56	64	72	80	88	96	104	112	120
refresh interval		refresh window		flags	filu length reg	segment				vpid	page	word	bit		
flags : 32 - scavenger trap 33 - cs out of range 34 - page crossing 35 - cache miss				36 - fill mode 37 - physical last 38 - write last 39 - mar modified		40 - incomplete				dirty v					

Space kinds

- 1 - Control
- 2 - Type
- 3 - Queue
- 4 - Data
- 5 - Import
- 6 - Code
- 7 - system

segment	vpid	page	lru	pl	if	b
page state : 00 - invalid 01 - F/v 10 - F/O 11 - loading	flags : 58 - wired 59 - permanent 60 - writable					

SOME USEFUL TAGS

Discrete : 00 (80)	Record : 44 (C4)	Subprogram : 06 (86)	Wall Subprogram : 76	Seg Heap : 38 (B8)
Access : 10 (90)	Variant Record : 4C (CC)	(Elaborated) : 16 (96)	Utility : 68 (D8)	
Task : 18 (98)	Vector : 6C (EC)	(Visible) : 26 (A6)	Accept : 46 (C6)	
Package : 58 (D6)	Matrix : 74 (F4)	(Visible & Elaborated) : 36 (B6)	Interface : 56 (D6)	
Float : 08 (88)	Array : 7C (FC)		Exception Var : 7E (FE)	

BLOCKED STATES

Unblocked : 00	Terminable At End : 07	In FS Rendezvous : 0E	Blocking On Accept : 18
Declaring Module : 01	Blocking On Entry : 08	In Wait Svc : 0F	Blocking On Select : 19
Awaiting Activation : 02	Delaying On Entry : 09	Delay In Wait Svc : 10	Delaying On Select : 1A
Activating Module : 03	Attempting Entry : 0A	Blocking On Abort : 11	Await Children Select : 1B
Activating Tasks : 04	Delaying : 0B	Deleted : 12	Terminable In Select : 1C
Awaiting Task Activ : 05	Aborting Module : 0C	Aborted While In MTS : 13	
Awaiting Children : 06	Terminated : 0D	In_MTS_Rendezvous : 14	

offset

0	8	16	24	32	40	48	56	64	72	80	88	96	104	112	120		
C	slice stuff								(09)							60C	
B	debug interface subprogram								subprg var (36)							56C	
A	delay days (18 bits)	delay ticks (36 bits)			m f sched t l alloc s g  (49)		scheduling_group		debugging state						50C		
9	breakpoint scope			type extent		aux alloc (31)		queue extent		data extent				48C			
8	distributor's name			flgs	control extent		contrl alloc (41)		breakpoint mask						40C		
7	dependence site name			dependence site offset		depend link (19)								38C			
6	our type name			flgs	our type offset		static conn (11)		our import stack name		declarer's name				30C		
5							micro state2 (29)								28C		
4							micro state1 (21)								20C		
3	queue successor name			type tos		aux state (31)		inner frame		data tos						18C	
2	current slice time	blkd state	m pri d ori ty	flgs	control tos		contrl state (01)		current macro pc						1x 10C		
1	outer frame name			flgs	type frame		activ link (31)		control pred		block start		data frame				10C
0	enclosing frame name			flgs	enclosing frame offset		activ state (71)		return address segment		children start offset		return address offset/index				1x 00C
0	0	10	11	12	13	14	15	16	17	18	19	20	21	22	23		

# TASKING - STATES

- See handout *(next page)*
- Described in debugger manual
- In bad situations, a task may get "stuck" in an unusual state.
- Examples of this include:
  - unrecoverable disk error
  - microcode bug
  - I/O problems
- Diagnosis
  - Kernel Show\_Task\_States command. Normally shows "interesting" tasks.
  - Show\_Tasks environment command. This is in the System\_Maintenance subsystem.
  - Debugger Task\_Display command if job is running under the debugger
  - more shortly

## Handout # 12 - Task states

- **Unblocked—“running”**—The task is ready to execute or is currently executing. (Higher-priority tasks may prevent the task from running, but there is no condition associated with the task to keep it from running.)
- **Declaring\_Module—“waiting for child elaboration”**—The module is in the process of declaring a task or package.
- **Awaiting\_Activation—“waiting for parent elaboration”**—The module is waiting for a package or task it has declared to become activated.
- **Activating\_Module—“activating child packages”**—The module is waiting for children packages to become elaborated.
- **Activating\_Tasks—“activating child tasks”**—The module is telling its child tasks to become activated.
  - **Awaiting\_Task\_Activation—“waiting for task activation”**—The module has told its child tasks to become activated and is waiting for confirmation that each has been activated.
  - **Awaiting\_Children—“waiting for children”**—The task is waiting for its children tasks to be terminated.
  - **Terminable\_At\_End—“package completed”**—The task has reached its end and will terminate according to Ada rules when its parent decides it's time to terminate.
  - **Blocking\_On\_Entry—“waiting at entry for accept”**—The task is blocked waiting for an entry call to be accepted.
  - **Delaying\_On\_Entry—“waiting at timed entry for accept”**—The task is blocked at a timed entry call waiting for the called task to accept the call.
  - **Attempting\_Entry—“attempting entry call”**—Reserved for future use.
  - **Delaying—“waiting for delay”**—The task is blocked at a delay statement.
  - **Aborting\_Module—“begin aborted”**—The module is in the process of aborting a task it has declared.
  - **Terminated—“terminated”**—The task is terminated.
  - **In\_FS\_Rendezvous—“state presently unknown”**—The task is processing a page fault. This transient state is unlikely to be seen.
  - **In\_Wait\_Service—“in wait service”**—The task is blocked in an Environment service.
  - **Delaying\_In\_Wait\_Service—“delaying in wait service”**—The task is blocked in an Environment service and has a maximum wait time.
  - **Blocking\_in\_Abort—“being aborted”**—The module is in the process of making another task abnormal (see *Ada Language Reference Manual*, Section 9.10).
  - **Blocking\_On\_Accept—“waiting at accept for entry call”**—The task is blocked at an accept statement, waiting for an entry call.
  - **Blocking\_On\_Select—“waiting at select for entry call”**—The task is blocked at a select statement, waiting for an entry call.
  - **Delaying\_On\_Select—“waiting at select-delay for entry call”**—The task is blocked in a select with a delay alternative waiting for an entry call.
  - **Awaiting\_Children\_In\_Select—“waiting at select-terminate for entry call”**—The task is executing at a select, with a terminate alternative, and has dependent tasks.
  - **Terminable\_In\_Select—“waiting at select-terminate for entry call”**—The task is blocked in a select, with a terminate alternative, waiting for an entry call and can terminate.

# TASKING - KERNEL SHOW\_TASK\_STATES

- Kernel CI command
- By default, shows names and states of "interesting" tasks
- Simple use: on IOA console, type ↑Z until the kernel prompt appears, then enter:

Show\_Task\_States

*(example next page)*

```

EEDB: kernel
Kernel: show_task_states
CACHE/DISK [CACHE]:
16#4388004#; IN_WAIT_SERVICE X25_WAIT; PRI 1
16#27805#; UNBLOCKED; PRI 15
16#488D9#; IN_WAIT_SERVICE TCP_IP_INPUT_WAIT; PRI 14
16#127CD804#; DELAYING_IN_WAIT_SERVICE U031; PRI 3
16#1245A804#; IN_WAIT_SERVICE TCP_IP_INPUT_WAIT; PRI 3
16#5B8D3#; UNBLOCKED; PRI 14
16#5B4D3#; UNBLOCKED; PRI 14
16#52CEE#; UNBLOCKED; PRI 8
16#5BCD3#; UNBLOCKED; PRI 14
16#5C4D7#; IN_WAIT_SERVICE TCP_IP_INPUT_WAIT; PRI 4
16#7C8DD#; DELAYING_IN_WAIT_SERVICE NATIVE_DEBUGGING_WAIT; PRI 8
16#7C4DD#; DELAYING_IN_WAIT_SERVICE NATIVE_DEBUGGING_WAIT; PRI 8
16#7CCDD#; DELAYING_IN_WAIT_SERVICE NATIVE_DEBUGGING_WAIT; PRI 8
Kernel: enable_priv_cmds
You are enabling a set of commands which must be used
with extreme care. They should be used only by
knowledgeable support personnel. These commands can
easily crash/hang the machine; some can competely trash
the state of the machine such that you must recover the
machine from backup tapes.
Proceed [FALSE]: yes
Password: secret
*Kernel: lmr
KIND [1]:
NAME [0]: 16#794dd@
EXPECTED VALUES ARE: 0 .. 4294967295
NAME [0]: 16#794dd#
OFFSET [0]:
COUNT [0]: 5
  not WIRED  not IN_TRANSIT  DIRTY  not PERMANENT  READ_WRITE
16#00#  16#000794DD0000007F# 16#0000000000180001#
16#01#  16#000794DD9800003F# 16#FFFFF00200000000#
16#02#  16#FFB802D810000D01# 16#DD021D0300000422#
16#03#  16#0007C8DD00000ABF# 16#00014000000000C1#
16#04#  16#0000000000000021# 16#0000063E00000000#
*Kernel: job
VPID [4]: 16#dd#
  Job Pri Stat CPU%  ModCt Cache  Disk  PgLim DskWts D/S JSegSz WsSiz WsLim
  --- --- --- --- --- --- --- --- --- --- --- --- ---
  221  6 I,AT  0      8    63      0   8000    88  0    50    50    50
*Kernel: quit
EEDB:

```

# **TASKING - EXCEPTION PROPAGATION**

- **If an exception propagates out of a package that package terminates and the exception is re-raised at the point of declaration of the package. The raise PC is set to the point of declaration. This has debugging implications.**
- **If you did not catch the original raise in the debugger, it may be difficult to locate the source of the exception without re-running the program.**

# TASKING - WAIT SERVICE

- **Wait service is a wait/signal mechanism used by the environment.**
- **It allows a task to block until another task signals (rousts) it.**
- **The wait service is used as part of normal operation to wait for I/O transfers and resource availability.**
- **The wait service is also used when a serious error is detected. Then, the affected task may be blocked in the wait service so someone can look at it and evaluate the situation.**
- **The kernel Show\_Task\_States command shows stopped in the waiter service and displays the reason they are stopped**



**ENVIRONMENT\_DEBUGGING\_WAIT**

**is the reason that appears for such fatal conditions.**

- **Tasks stopped in the wait service can be manually roused by using a kernel command.**





# TASKING - WAIT SERVICE (CONT'D.)

- **This is generally dangerous and should only be done if you know what you are doing or have advice from someone who does. Permanent, fatal system damage can result if a task is roused at the wrong time, although this is rare**

**\*Kernel:        ROUST**

**Task id:        16#127FB04#**

- **Kernel always takes decimal numeric parameters. Used based notation (16#...#) for hex addresses and names.**

# TASKING - SHOW\_TASKS COMMAND

- **Displays stack of specific or group of tasks**
- **Useful to monitor progress or diagnose hangs**
- **Return dump of this in hung editor or session case**
- **When given job id, attempts to locate as many tasks as possible in that job**
- **!Commands.System\_Maintenance.Current\_Combined\_Mload provides source information for the environment itself. This is packages level information only.**
- **There are parameters to specify the type of stock display, and to filter the task search.**

# TASKING - SHOW\_TASKS COMMAND (CONT'D.)

- Parameters

**Task\_or\_Job\_ID**

**Show\_Stack**                    **if false, only list task names and states**

**Show\_Full\_Stack**            **show each stack word one way or another, not just frames**

**Show\_Hex\_Stack**            **show control words in hex**

**Show\_Packages**            **display packages as well as tasks**

**Name\_Filter**                **skip entries that don't include this as a name fragment**

**Max\_Tasks**                **stop after this many tasks have been found. You may need to increase this to get all tasks in a job.**

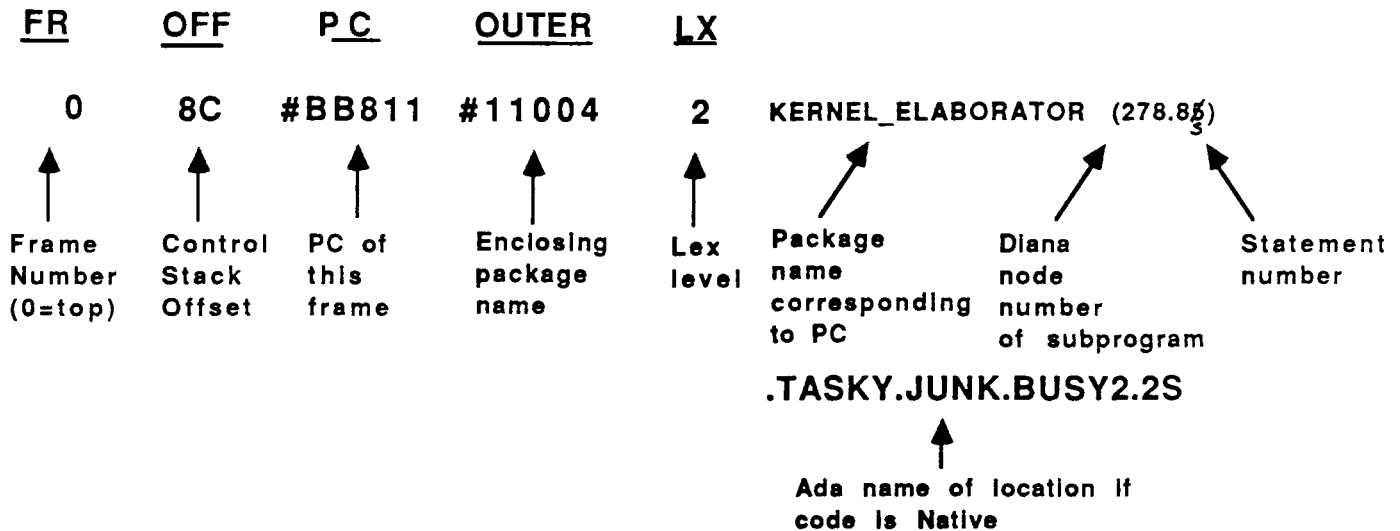
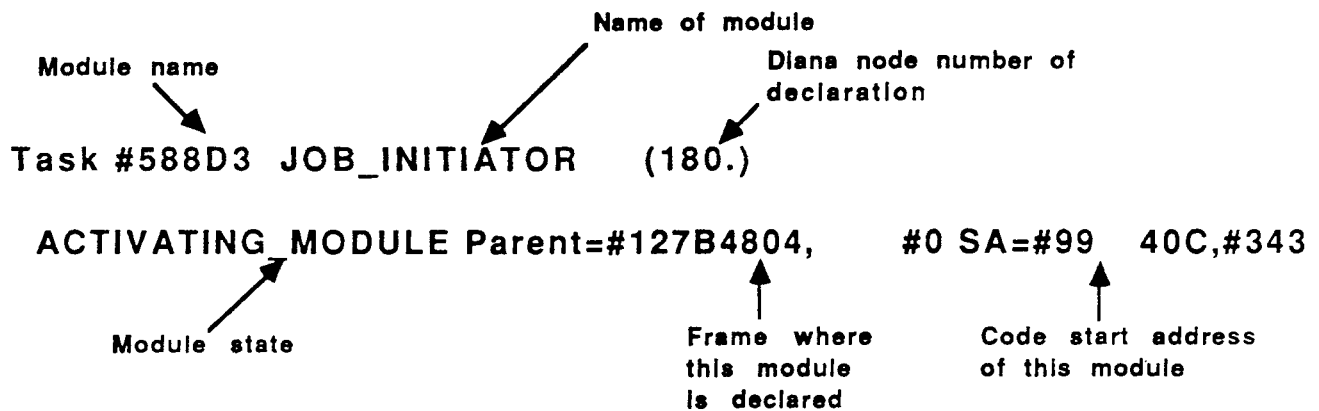
**Names\_File**                **file containing environment package/segment map**

**Debug\_File**                **not used**

# TASKING - SHOW\_TASKS COMMAND (CONT'D.)

- Output

- First, symbol segment files are read
- Next, stacks are scanned to locate as many modules as possible
- For each module



- No Ada names available for elaboration code, loaded main programs or code views, or environment code if file is missing.

User Status on August 2, 1987 at 2:13:54 PM

User	Line	Job	S	Time	Job Name
PHIL	242	223	IDLE	1:26:18	[PHIL.S 1 Command]
		238	RUN	1:26:27	[PHIL.S 1 Editor]
		251	RUN	4.382	!USERS.PHIL & WHAT.USERS

```
Show_Tasks (Task_Or_Job_Id => 238,
Show_Stack => True,
Show_Full_Stack => False,
Show_Hex_Stack => False,
Show_Packages => False,
Name_Filter => "",
Max_Tasks => 20,
Names_File => "!Commands.System_Maintenance.@_Mload",
Debug_File => "!Commands.System_Maintenance.@_Rdebug");
```

!USERS.PHIL.GURU COURSE & SHOW TASKS STARTED 2:15:36 PM

Processing file: !COMMANDS.SYSTEM\_MAINTENANCE.A\_9\_18\_0\_COMBINED\_MLOAD'V(1) ... 1326 defined.  
 Processing file: !COMMANDS.SYSTEM\_MAINTENANCE.LATEST\_COMBINED\_MLOAD'V(1) ... 1349 defined.  
 Processing file: !COMMANDS.SYSTEM\_MAINTENANCE.UOC931\_MLOAD'V(2) ... 33 defined.

Max\_Tasks reached; not all tasks will be displayed

Task #1A0EE USER (210.)  
 BLOCKING\_ON\_SELECT Parent=#198EE, #0 SA=#6B300B, #B

Fr Off	Pc	Outer	Lx
0	26	#6B1C0B, #293#12717804	2 ROUTER (225.2s)
1	1E	#6B300B, #136	#1A0EE 2 USER (643.3s)
2	0	#6B300B, #59	#1A0EE 1 USER (24.4s)

Task #198EE CORE EDITOR (444.)  
 BLOCKING\_ON\_SELECT Parent=#1F5404, #0 SA=#6B380B, #B

Fr Off	Pc	Outer	Lx
0	76	#6B380B, #E53	#198EE 2 CORE EDITOR (1785.2s)
1	0	#6B380B, #239	#198EE 1 CORE EDITOR (30.5s)

Task #1A4EE RUN MACRO (230.)  
 BLOCKING\_ON\_SELECT Parent=#198EE, #0 SA=#6B2C0B, #B

Fr Off	Pc	Outer	Lx
0	1A	#6B1C0B, #35A#12717804	2 ROUTER (319.1s)
1	15	#6B2C0B, #40	#1A4EE 2 RUN_MACRO (94.1s)
2	0	#6B2C0B, #28	#1A4EE 1 RUN_MACRO (10.4s)

Task #19CEE INPUT (190.)  
 IN\_WAIT\_SERVICE Parent=#198EE, #0 SA=#6B280B, #B

Fr Off	Pc	Outer	Lx
0	AF	#B82C11, #18F	#19C04 2 PROCESSOR_MANAGER (1272.5s)
1	AB	#B82C11, #143	#19C04 2 PROCESSOR_MANAGER (638.2s)
2	A7	#18F408, #4F	#67804 2 WAIT_SERVICE (85.1s)
3	A3	#715008, #1000	#6FC04 2 TCP_IP_DRIVER (541.2s)
4	97	#715008, #AEF	#6FC04 2 TCP_IP_DRIVER (295.16s)
5	8F	#546808, #23E#12459404	2 TCP_IP_COUPLER (338.1s)
6	88	#55B008, #37E	#273404 2 TRANSPORT (372.1s)
7	81	#718C08, #803	#272804 4 TELNET_PROTOCOL (2155.1s)
8	76	#718C08, #731	#272804 3 TELNET_PROTOCOL (2129.4s)
9	6D	#718C08, #572	#272804 2 TELNET_PROTOCOL (2072.2s)
10	63	#718C08, #6CD	#272804 3 TELNET_PROTOCOL (3216.2s)
11	5C	#718C08, #375	#272804 2 TELNET_PROTOCOL (384.9s)
12	4F	#71D808, #19E	#276404 2 TELNET_SERVER (121.11s)
13	46	#6EC09, #F8#12490804	2 TELNET_COUPLER (282.1s)
14	3D	#6C1008, #5F4	#36B004 3 TERMINAL_IO (2100.3s)
15	36	#6C1008, #326	#36B004 2 TERMINAL_IO (300.7s)
16	2C	#6C1008, #37B	#36B004 2 TERMINAL_IO (370.1s)
17	22	#43E008, #42B	#FBC04 2 TERMINAL_MANAGER (412.2s)
18	1A	#6B280B, #A7	#19CEE 2 INPUT (309.2s)
19	0	#6B280B, #38	#19CEE 1 INPUT (18.4s)

Task #1BCEE SCREEN TASK (289.)  
 DELAYING\_ON\_SELECT Parent=#12703004, #0 SA=#68E00B, #B

Fr Off	Pc	Outer	Lx
0	1F2	#68E00B, #4E71	#1BCEE 2 SCREEN_TASK (5077.2s)
1	0	#68E00B, #775	#1BCEE 1 SCREEN_TASK (16.5s)

Task #1F8EE COMMAND EDITOR (419.)  
 BLOCKING\_ON\_SELECT Parent=#12768004, #0 SA=#68A80B, #1713

Fr Off	Pc	Outer	Lx
0	5C	#68A80B, #23EB	#1F8EE 2 COMMAND_EDITOR (2737.2s)
1	0	#68A80B, #18B0	#1F8EE 1 COMMAND_EDITOR (30.4s)

Task #1E4EE TEXT OBJECT EDITOR (311.)  
 BLOCKING\_ON\_SELECT Parent=#22A804, #0 SA=#61000B, #2A43

Fr Off	Pc	Outer	Lx
0	6E	#61000B, #38DF	#1E4EE 2 TEXT_OBJECT_EDITOR (1205.2s)
1	0	#61000B, #2BBF	#1E4EE 1 TEXT_OBJECT_EDITOR (6.4s)

Task #1D4EE BUFFER TASK (121.)  
 BLOCKING\_ON\_SELECT Parent=#12717804, #0 SA=#6B180B, #B

Fr Off	Pc	Outer	Lx
0	36	#6B180B, #23F	#1D4EE 2 BUFFER_TASK (1102.2s)
1	0	#6B180B, #9E	#1D4EE 1 BUFFER_TASK (24.4s)

Task #230EE ADA EDITOR (174.)

Fr Off	Pc	Outer	Lx	
0	5C #5C080B, #1E63	#230EE	2	ADA_EDITOR (5041.2s)
1	0 #5C080B, #1729	#230EE	1	ADA_EDITOR (46.4s)

-----

!USERS.PHIL.GURU COURSE % SHOW\_TASKS STARTED 2:31:05 PM

-----

Processing file: !COMMANDS.SYSTEM\_MAINTENANCE.LATEST\_COMBINED\_MLOAD/V(1) ...  
1349 defined.

scanning stack of #588D3 tos = 147  
scanning stack of #5A0D3 tos = 13  
scanning stack of #5B0D3 tos = 16  
scanning stack of #5B8D3 tos = 42  
scanning stack of #5B4D3 tos = 46  
scanning stack of #5ACD3 tos = 13  
scanning stack of #59CD3 tos = 26  
scanning stack of #5BCD3 tos = 51

Task #588D3 JOB\_INITIATOR (180.)  
ACTIVATING\_MODULE Parent=#127B4804, #0 SA=#9940C, #343

Fr Off	Pc	Outer	Lx	
0	8C #BBB811, #384	#11004	2	KERNEL_ELABORATOR (278.8s)
1	81 #27009, #813	#1243F804	2	ENVIRONMENT_ELABORATOR_DATABASE (1083.1s)
2	7C #DD3C16, #1D4	#192004	2	EXECUTION (131.1s)
3	76 #DD3C16, #12B	#192004	2	EXECUTION (81.4s)
4	70 #56EC0B, #5A9	#1278B004	2	ADA_INTERPRETER (1628.2s)
5	68 #56EC0B, #102	#1278B004	2	ADA_INTERPRETER (18.10s)
6	59 #68A80B, #1038	#12768004	2	COMMAND_EDITOR (4110.34s)
7	48 #68A80B, #29A	#12768004	2	COMMAND_EDITOR (80.4s)
8	37 #5C080B, #956	#12766804	2	ADA_EDITOR (1138.5s)
9	32 #6B100B, #25B	#12798004	2	REGISTRATION (191.1s)
10	2E #6B140B, #339	#1FA004	2	OE_COUPLER (306.1s)
11	24 #56500B, #2577	#127B3C04	2	LIBRARY_OBJECT_EDITOR (820.9s)
12	1F #AB40C, #1107	#127B4004	2	LINE_EDITOR (664.3s)
13	0 #9940C, #3B2	#588D3	1	JOB_INITIATOR (20.6s)

Task #5B8D3 .TASKY.BUSY  
UNBLOCKED Parent=#59CD3, #14 SA=#21D03, #AB

Fr Off	Pc	Outer	Lx	
0	31 #2ED808, #167	#4B404	2	BOUNDED (145.2s)
1	2D #8DC08, #56F	#4BC04	3	STRING_UTILITIES (644.16s)
2	29 #8DC08, #553	#4BC04	3	STRING_UTILITIES (644.11s)
3	21 #8DC08, #232	#4BC04	2	STRING_UTILITIES (120.6s)
4	1B #8DC08, #208	#4BC04	2	STRING_UTILITIES (96.1s)
5	12 #21D03, #8D	#5B0D3	2	.TASKY.JUNK.BUSY2.2s
6	0 #21D03, #BA	#5B8D3	1	.TASKY.BUSY.7s

Task #5B4D3 .TASKY.BUSY  
UNBLOCKED Parent=#59CD3, #14 SA=#21D03, #AB

Fr Off	Pc	Outer	Lx	
0	2A #2ED808, #2AD	#4B404	2	BOUNDED (326.1d)
1	21 #8DC08, #223	#4BC04	2	STRING_UTILITIES (120.1s)
2	1B #8DC08, #208	#4BC04	2	STRING_UTILITIES (96.1s)
3	12 #21D03, #8D	#5B0D3	2	.TASKY.JUNK.BUSY2.2s
4	3D #8DC08, #553	#4BC04	3	STRING_UTILITIES (644.11s)
5	39 #8DC08, #553	#4BC04	3	STRING_UTILITIES (644.11s)
6	35 #8DC08, #553	#4BC04	3	STRING_UTILITIES (644.11s)
7	31 #8DC08, #553	#4BC04	3	STRING_UTILITIES (644.11s)
8	2D #8DC08, #553	#4BC04	3	STRING_UTILITIES (644.11s)

Package #59CD3 unknown - no table  
DELAYING Parent=#588D3, #8C SA=#43502, #13

Fr Off	Pc	Outer	Lx	
0	14 #21D03, #42	#5A0D3	2	.TASKY.3s
1	12 #43102, #1C	#5ACD3	2	unknown - no table
2	0 #43502, #39	#59CD3	1	unknown - no table

Task #5BCD3 .TASKY.BUSY  
UNBLOCKED Parent=#59CD3, #14 SA=#21D03, #AB

Fr Off	Pc	Outer	Lx	
0	29 #8DC08, #56A	#4BC04	3	STRING_UTILITIES (644.16s)
1	21 #8DC08, #232	#4BC04	2	STRING_UTILITIES (120.6s)
2	1B #8DC08, #208	#4BC04	2	STRING_UTILITIES (96.1s)
3	12 #21D03, #8D	#5B0D3	2	.TASKY.JUNK.BUSY2.2s
4	0 #21D03, #BA	#5BCD3	1	.TASKY.BUSY.7s

!USERS.PHIL.GURU COURSE % SHOW\_TASKS STARTED 2:32:03 PM

-----

Processing file: !COMMANDS.SYSTEM\_MAINTENANCE.LATEST\_COMBINED\_MLOAD/V(1) ...  
1349 defined.

scanning stack of #588D3 tos = 147  
scanning stack of #5A0D3 tos = 13  
scanning stack of #5B0D3 tos = 16  
scanning stack of #5B8D3 tos = 79  
scanning stack of #5B4D3 tos = 63  
scanning stack of #5ACD3 tos = 13  
scanning stack of #59CD3 tos = 26  
scanning stack of #5BCD3 tos = 48

Task #588D3 JOB\_INITIATOR (180.)  
ACTIVATING\_MODULE Parent=#127B4804, #0 SA=#9940C, #343

Package #5A0D3 .TASKY  
TERMINABLE\_AT\_END Parent=#59CD3, #0 SA=#21D03, #B

Package #5B0D3 .TASKY.JUNK/SPEC.1d  
TERMINABLE\_AT\_END Parent=#59CD3, #14 SA=#21D03, #63

```
Task #5B8D3 .TASKY.BUSY
UNBLOCKED Parent=#59CD3, #14 SA=#21D03, #AB

Task #5B4D3 .TASKY.BUSY
UNBLOCKED Parent=#59CD3, #14 SA=#21D03, #AB

Package #5ACD3 unknown - no table
TERMINABLE_AT_END Parent=#59CD3, #0 SA=#43102, #B

Package #59CD3 unknown - no table
DELAYING Parent=#588D3, #8C SA=#43502, #13

Task #5BCD3 .TASKY.BUSY
UNBLOCKED Parent=#59CD3, #14 SA=#21D03, #AB
```

# TASKING - REVIEW

- **Task lifecycle**

- **TCB**

- **Task state**

- **Commands:**

**Show\_Task\_States**

**Show\_Tasks**

**Roust**

- **Packages and tasks**

- **Wait service**



# VIRTUAL PROCESSORS

- **A Virtual Processor (VP) is a grouping of resources that the environment deals with as a unit.**
- **All objects on a particular VP exist on the same disk volume.**
- **All tasks running on a particular VP are scheduled as a unit.**
- **Part of the segment name of each segment identifies the VP. The bottom 10 bits of the segment name are the VPID number.**
- **Of the 1024 VPIDs, VP 1 through 255 are used for jobs. Others are used for data objects.**

# VIRTUAL PROCESSORS - SPECIAL VPs

<b>0 - 3</b>	<b>Reserved for kernel</b>
<b>4</b>	<b>System</b>
<b>5</b>	<b>Daemons</b>
<b>6 - 7</b>	<b>Reserved for System</b>
<b>8 - 25</b>	<b>Cross-compiled code segments</b>
<b>26 - 255</b>	<b>User jobs</b>
<b>256 - 1023</b>	<b>Storage of "permanent" data</b>

# VIRTUAL PROCESSORS - JOBS

- The VP number for a Job VP is the same as the job id.
- Information associated with a job
  - Priority
  - Name
  - Class - Attached, Detached, Server, CE, OE
  - State - Run, Wait, Idle, Disabled, Queued
  - Resource use
  - Access control identity
  - Session
  - Root task
  - Start time
  - Elapsed time
  - Termination message
  - Accounting information
- Task Name/VP conversion

<u>Task id (hex)</u>	<u>Job id</u>
4388004	4
1C8E7	E7 = 231
220DF	DF = 223

- Take right two hex digits. Next two higher bits always are zeros. (Why?)

*always on VP @ .1255*

# VIRTUAL PROCESSORS - JOB INFORMATION

- **Resource Use**

<b>CPU MS/S or Percentage</b>	<b>computed by scheduler</b>
<b>Module Count</b>	<b>Number of modules (tasks and packages). Computed by the microcode.</b>
<b>Cache</b>	<b>Number of pages in cache occupied by the jobs modules. Computed by Microcode.</b>
<b>Disk</b>	<b>Number of pages of disk occupied by the job's modules. Computed by the kernel.</b>
<b>Page Limit</b>	<b>Page limit - set by the job.</b>
<b>Disk Waits</b>	<b>Number of times job has waited for disk I/O. Roughly, the number of page faults (including I/O).</b>
<b>Disk Waits/Second</b>	<b>Rate of disk waits.</b>
<b>Job Segment Size</b>	<b>Number of pages in the job's garbage heap.</b>
<b>Working Set Size</b>	<b>Current job working set size in pages. Computed by scheduler.</b>
<b>Working Set Limit</b>	<b>Maximum size working set can be before scheduler discriminates against the job.</b>

# VIRTUAL PROCESSORS - JOB INFORMATION (CONT'D.)

- **Identity**

**Root Task**                      **The task id of the root module of the job. May be zero for system jobs and editor jobs.**

**Job Name**                      **A string name of the job. Usually the context and command.**

**Job Segment Name**              **Segment name of the job garbage heap. Environment code allocates space here on behalf of the job.**

**Group Identity**                  **The set of groups the job is a member of. Used for access control.**

- **Priority**                      **The priority of the job. User jobs start at 6. Bigger numbers are lower priorities.**

# VIRTUAL PROCESSORS - GETTING INFORMATION

- **From the Environment**

<b>What.Users</b>	<b>Shows user, session, elapsed time, state, and name</b>
<b>What.Jobs</b>	<b>Updated OE display of jobs. Can be set for different thresholds.</b>
<b>Show_Jobs</b>	<b>Shows, for individual, active, or all jobs, priority, CPU, status, module, cache, and disk page counts, etc. Very useful for watching resources go.</b>
<b>Show_Job_Names</b>	<b>Similar to Show_Jobs, but displays names, root task, and job segment name.</b>
<b>Show_Identity</b>	<b>Shows access control and user identity of a job.</b>

- **From the kernel (or kernel via EEDB)**

<b>Jobs</b>	<b>Same as Show_Jobs. Threshold =0 for interesting jobs only, = 1 for all.</b>
<b>Job</b>	<b>Same as Show_Jobs, but for a specific job.</b>
<b>Job_Names</b>	<b>Same as Show_Job_Names</b>
<b>Job_Name</b>	<b>Same as Show_Job_Names, but for a specific job.</b>

# VIRTUAL PROCESSORS - GETTING INFORMATION (CONT'D.)

<b>Jobs_MTS</b>	<b>Similar to Show_Jobs, includes scheduler information</b>
<b>Job_MTS</b>	<b>Jobs_MTS for a specific job</b>
<b>Show_VPs</b>	<b>Shows for each VP, the volume the VP is assigned to</b>
<b>MTSq</b>	<b>Show medium term scheduler queues</b>
<b>Hogs</b>	<b>Find big objects on specified VPs</b>

- **From programs**

<b>package</b>	<b>!.Commands.Scheduler</b>	<b>--</b>	<b>get information</b>
<b>package</b>	<b>!.Tools.System_Uilities</b>	<b>--</b>	<b>get information</b>
		<b>--</b>	<b>page limits</b>
<b>package</b>	<b>!.Commands.Program</b>	<b>--</b>	<b>create jobs</b>

What.Users

User Status on August 2, 1987 at 11:57:05 PM

User	Line	Job	S	Time	Job Name
*SYSTEM	-	4	RUN	1/05:41	System
		5	RUN	1/05:41	Daemons
		225	IDLE	1/05:32	Print Spooler
		227	IDLE	1/05:32	Console Command Interpreter
		228	IDLE	1/05:31	Problem Report Transfer Server
		231	IDLE	1/05:32	Ftp Server
		236	IDLE	1/05:32	Archive Server
NETWORK_PUBLIC	-	217	IDLE	6:30:35	Print Queue Server
PHIL	245	211	RUN	9:26:11	!USERS.PHIL.TEST_AREA & !USERS.PHIL.TEST_AREA.TASKY
		213	IDLE	13:12.090	[PHIL.S_1 Debugger]
		215	RUN	1:10:41	[PHIL.S_1 Editor]
		219	IDLE	1:10:36	[PHIL.S_1 Command]
		220	RUN	1.016	!USERS.PHIL.GURU COURSE & WHAT.USERS
		221	IDLE	13:16.808	!USERS.PHIL.TEST_AREA & !USERS.PHIL.TEST_AREA.TASKY

Show\_Jobs (Job => 0, Active\_Jobs\_Only => True);

Job	Pri	Stat	CPU%	ModCt	Cache	Disk	PgLim	DskWts	D/S	JSegSz	WsSiz	WsLim
4	0	R,AT	1	4052	9184	12101	65536	116128	0	777	11007	11000
211	0	R,DT	98	8	17	13	8000	12	0	78	26	50
213	0	I,SV	0	170	831	7	16000	499	0	18	58	75
215	6	R,CE	0	47	438	14	16000	400	0	67	164	150
228	0	I,SV	0	5	12	17	8000	200	0	178	42	75

Show\_Job\_Names;

Job	CPU%	Root	Job	Seg	Name
4	0	0	14F9903		System
211	95	588D3	1277902		!USERS.PHIL.TEST_AREA & !USERS.PHIL.TEST_AREA.T
213	0	BB4D5	1566503		[PHIL.S_1 Debugger]
215	0	0	1283D02		[PHIL.S_1 Editor]
228	0	2B4E4	1133901		Problem Report Transfer Server



!TOOLS.SYSTEM UTILITIES'V(1)

```

with Directory;
with Machine;
with Calendar;
with System;

package System_Uilities is

  pragma Subsystem (Tools, Closed);
  pragma Module_Name (4, 3973);

  subtype Job_Id is Machine.Job_Id range 4 .. 255;
  subtype Session_Id is Machine.Session_Id;

  subtype Version is Directory.Version;
  subtype Object is Directory.Object;

  subtype Port is Natural range 0 .. 4 * 16 * 16;
  subtype Tape is Natural range 0 .. 4;

  subtype Byte_String is System.Byte_String;

  -- Job (Process) characteristics
  function Get_Job return Job_Id;
  function Priority
    (For_Job : Job_Id := System_Uilities.Get_Job) return Natural;
  function Elapsed
    (For_Job : Job_Id := System_Uilities.Get_Job) return Duration;
  function Cpu (For_Job : Job_Id := System_Uilities.Get_Job) return Duration;
  function Job_Name
    (For_Job : Job_Id := System_Uilities.Get_Job) return String;

  -- Active Session characteristics
  function Get_Session return Session_Id;
  function Get_Session (For_Job : Job_Id) return Session_Id;

  function Session_Name
    (For_Session : Session_Id := System_Uilities.Get_Session)
    return String;
  function User_Name
    (For_Session : Session_Id := System_Uilities.Get_Session)
    return String;
  function Terminal (For_Session : Session_Id := System_Uilities.Get_Session)
    return Port;
  function Terminal (For_Session : Session_Id := System_Uilities.Get_Session)
    return Version;
  function Terminal (For_Session : Session_Id := System_Uilities.Get_Session)
    return Object;
  function Session (For_Session : Session_Id := System_Uilities.Get_Session)
    return Version;
  function Session (For_Session : Session_Id := System_Uilities.Get_Session)
    return Object;
  function User (For_Session : Session_Id := System_Uilities.Get_Session)
    return Version;
  function User (For_Session : Session_Id := System_Uilities.Get_Session)
    return Object;

  -- Inactive Session characteristics
  function Home_Library (User : String := User_Name) return String;

```

August 2, 1987 at 1:22:58 AM

!TOOLS.SYSTEM UTILITIES'V(1)

2

```

function Last_Login (User : String; Session : String := "")
  return Calendar.Time;
function Last_Logout (User : String; Session : String := "")
  return Calendar.Time;
function Logged_In (User : String; Session : String := "") return Boolean;

-- Names for Text_IO/Simple_Text_IO standard file names

function Output_Name
  (For_Session : Session_Id := System_Uilities.Get_Session)
  return String;
function Input_Name
  (For_Session : Session_Id := System_Uilities.Get_Session)
  return String;
function Error_Name
  (For_Session : Session_Id := System_Uilities.Get_Session)
  return String;
function Tape_Name (Drive : Tape := 0) return String;

-- Terminal characteristics

subtype Stop_Bits_Range is Integer range 1 .. 2;
subtype Character_Bits_Range is Integer range 5 .. 8;
type Parity_Kind is (None, Even, Odd);

function Terminal_Name
  (Line : Port := System_Uilities.Terminal) return String;
function Terminal_Type
  (Line : Port := System_Uilities.Terminal) return String;

function Input_Count
  (Line : Port := System_Uilities.Terminal) return Long_Integer;
function Output_Count
  (Line : Port := System_Uilities.Terminal) return Long_Integer;
-- The number of characters input/output from/to the specified terminal
-- since the machine was booted. Input from the terminal that has not
-- been read by a session or user program will not be counted as input.

function Input_Rate
  (Line : Port := System_Uilities.Terminal) return String;
function Output_Rate
  (Line : Port := System_Uilities.Terminal) return String;

function Parity (Line : Port := System_Uilities.Terminal)
  return Parity_Kind;
function Stop_Bits (Line : Port := System_Uilities.Terminal)
  return Stop_Bits_Range;
function Character_Size (Line : Port := System_Uilities.Terminal)
  return Character_Bits_Range;

function Xon_Xoff_Characters
  (Line : Port := System_Uilities.Terminal) return String;
function Xon_Xoff_Bytes
  (Line : Port := System_Uilities.Terminal) return Byte_String;

function Receive_Xon_Xoff_Characters
  (Line : Port := System_Uilities.Terminal) return String;
function Receive_Xon_Xoff_Bytes
  (Line : Port := System_Uilities.Terminal) return Byte_String;
-- returns a 2-element string consisting of Xon followed by Xoff

```

August 2, 1987 at 1:22:58 AM

```

function Flow_Control
  (Line : Port := System_Utilities.Terminal) return String;
function Receive_Flow_Control
  (Line : Port := System_Utilities.Terminal) return String;
-- return one of NONE, XON_XOFF, RTS, DTR

function Enabled (Line : Port := System_Utilities.Terminal) return Boolean;

function Disconnect_On_Disconnect
  (Line : Port := System_Utilities.Terminal) return Boolean;
function Logoff_On_Disconnect
  (Line : Port := System_Utilities.Terminal) return Boolean;
function Disconnect_On_Logoff
  (Line : Port := System_Utilities.Terminal) return Boolean;
function Disconnect_On_Failed_Login
  (Line : Port := System_Utilities.Terminal) return Boolean;
function Log_Failed_Logins
  (Line : Port := System_Utilities.Terminal) return Boolean;
function Login_Disabled
  (Line : Port := System_Utilities.Terminal) return Boolean;
function Detach_On_Disconnect
  (Line : Port := System_Utilities.Terminal) return Boolean;

-- Iterate over all active sessions

type Session_Iterator is private;
procedure Init (Iter : out Session_Iterator);
function Value (Iter : Session_Iterator) return Session_Id;
function Done (Iter : Session_Iterator) return Boolean;
procedure Next (Iter : in out Session_Iterator);

-- Iterate over all jobs for a session

type Job_Iterator is private;
procedure Init (Iter : out Job_Iterator;
  For_Session : Session_Id := Get_Session);
function Value (Iter : Job_Iterator) return Job_Id;
function Done (Iter : Job_Iterator) return Boolean;
procedure Next (Iter : in out Job_Iterator);

type Terminal_Iterator is private;
procedure Init (Iter : out Terminal_Iterator);
function Value (Iter : Terminal_Iterator) return Natural;
function Done (Iter : Terminal_Iterator) return Boolean;
procedure Next (Iter : in out Terminal_Iterator);

function System_Up_Time return Calendar.Time;
function System_Boot_Configuration return String;

function Image (Version : Directory.Version) return String;

procedure Set_Page_Limit (Max_Pages : Natural;
  For_Job : Job_Id := System_Utilities.Get_Job);

-- Set the upper limit for pages created by the specified job.
-- Attempts to create additional pages result in Storage_Error.
-- Requires operator capability if For_Job specifies a job belonging
-- to a user different from the caller. If For_Job parameter defaults,

```

```

-- then the limit applies to the calling job. In the worst case,
-- the job can allocate twice Max_Pages pages before getting a storage
-- error. Raises constraint_error if the job_id is illegal.

procedure Get_Page_Counts (Cache_Pages : out Natural;
  Disk_Pages : out Natural;
  Max_Pages : out Natural;
  For_Job : Job_Id := System_Utilities.Get_Job);

-- Return the counts for the specified job. Cache_Pages is the number of
-- pages presently in main memory; Disk_Pages is the number of pages that
-- have disk space allocated for them. Max_Pages is the current page
-- limit.

-- Operations for reading machine information:

type Bad_Block_Kinds is new Long_Integer range 0 .. 7;
Manufacturers_Bad_Blocks : constant Bad_Block_Kinds := 1;
Retargeted_Blocks : constant Bad_Block_Kinds := 2;
All_Bad_Blocks : constant Bad_Block_Kinds := 3;

type Block_List is array (Natural range <>) of Integer;

function Bad_Block_List
  (For_Volume : Natural;
  Kind : Bad_Block_Kinds := Retargeted_Blocks) return Block_List;
-- Return the list of bad blocks of the specified kind on the specified
-- disk. Return null array if Kind or volume are illegal.

function Get_Board_Info (Board : Natural) return String;
-- return information about the specified board in the machine. The
-- string identifies the information.
-- Board specifies the particular board:
-- 0 : IOA/IOC
-- 1 : SYS
-- 2 : SEQ
-- 3 : VAL
-- 4 : TYP
-- 5 : FIU
-- 100 : MEM0
-- 101 : MEM1
-- 102 : MEM2
-- 103 : MEM3
-- etc.
end System_Utilities;

```

# VIRTUAL PROCESSORS - COMMANDS TO CONTROL THEM

- **Environment**

- Job.Kill
- Scheduler.Set\_Job\_Attribute
- Scheduler.Set\_WSL\_Limits

*Scheduler.Disable  
; Enable*

- **Kernel**

- Disable\_Job                      **good to stop runaways and look around at them**
- Enable\_Job
- Abort\_Task                      *use in root task to kill job*

- **If a program has runaway and is using unexpected CPU or disk space (or job segment size large), you can disable it from the kernel and then use Show\_Tasks to dump the job state.**

# VIRTUAL PROCESSORS - REVIEW

- **Volumes**
- **VPIDs and Job Ids**
- **Job information**
- **Job resources**
- **Commands:**

**What.Users**

**What.Jobs**

**Show\_Jobs**

**Show\_Jobs\_Names**

**Show\_Identity**

**etc.**

```

with Machine;

package Scheduler is

  subtype Job_Id is Machine.Job_Id;
  subtype Cpu_Priority is Natural range 0 .. 6;
  subtype Milliseconds is Long_Integer;

  procedure Disable (Job : Job_Id);
  procedure Enable (Job : Job_Id);
  function Enabled (Job : Job_Id) return Boolean;

  function Get_Cpu_Priority (Job : Job_Id) return Cpu_Priority;

  type Job_Kind is (Ce, Oe, Attached, Detached, Server, Terminated);
  function Get_Job_Kind (Job : Job_Id) return Job_Kind;

  type Job_State is (Run, Wait, Idle, Disabled, Queued);
  function Get_Job_State (Job : Job_Id) return Job_State;

  -- returns the current state of job.
  -- RUN:      the job is currently runnable
  -- WAIT:     the job is runnable but being withheld by the scheduler.
  -- IDLE:     the job isn't using cpu time and has no unblocked tasks.
  -- DISABLED: an external agent has disabled the job from running.
  -- QUEUED:   the job is DETACHED and must wait for another to complete.

  function Get_Cpu_Time_Used (Job : Job_Id) return Milliseconds;

  -- returns the number of milliseconds of cpu time used by the job.
  -- belongs on the previous page, put here for compatability reasons

  function Disk_Waits (Job : Job_Id) return Long_Integer;

  -- returns the number of disk_waits the job has done since last initialized

  function Working_Set_Size (Job : Job_Id) return Natural;

  -- returns the number of pages in the job's working set.

  subtype Load_Factor is Natural;

  -- for run queues, number of tasks * 100

  procedure Get_Run_Queue_Load (Last_Sample : out Load_Factor;
                                Last_Minute : out Load_Factor;
                                Last_5_Minutes : out Load_Factor;
                                Last_15_Minutes : out Load_Factor);
  -- number of runnable tasks * 100

  procedure Get_Disk_Wait_Load (Last_Sample : out Load_Factor;
                                Last_Minute : out Load_Factor;
                                Last_5_Minutes : out Load_Factor;
                                Last_15_Minutes : out Load_Factor);
  -- number of tasks waiting for a page on the disk wait queue * 100

  procedure Get_Withheld_Task_Load (Last_Sample : out Load_Factor;
                                     Last_Minute : out Load_Factor;
                                     Last_5_Minutes : out Load_Factor;

```

```

                                Last_15_Minutes : out Load_Factor);
  -- returns the average number of tasks withheld from running by
  -- the scheduler * 100. In this call LAST_SAMPLE is the number of tasks
  -- held at the last 100ms scheduling cycle.

  procedure State;

  -- print scheduler state

  procedure Display (Show_Parameters : Boolean := True;
                    Show_Queues : Boolean := True);
  -- display current scheduler state and queues

  type Job_Descriptor is
    record
      The_Cpu_Priority : Cpu_Priority;
      The_State : Job_State;
      The_Disk_Waits : Long_Integer;
      The_Time_Consumed : Milliseconds;
      The_Working_Set_Size : Natural;
      The_Working_Set_Limit : Natural;
      The_Milliseconds_Per_Second : Natural;
      The_Disk_Waits_Per_Second : Natural;
      The_Maps_To : Job_Id;
      The_Kind : Job_Kind;
      The_Made_Runnable : Long_Integer;
      The_Total_Runnable : Long_Integer;
      The_Made_Idle : Long_Integer;
      The_Made_Wait : Long_Integer;
      The_Wait_Disk_Total : Long_Integer;
      The_Wait_Memory_Total : Long_Integer;
      The_Wait_Cpu_Total : Long_Integer;
      The_Min_Working_Set_Limit : Long_Integer;
      The_Max_Working_Set_Limit : Long_Integer;
    end record;

  function Get_Job_Descriptor (Job : Job_Id) return Job_Descriptor;

  -- use to get a consistent snapshot of a job's statistics.

  generic
    with procedure Put (Descriptor : Job_Descriptor);

  procedure Traverse_Job_Descriptors (First, Last : Job_Id);

  -- use to get a consistent, efficient snapshot of a range of
  -- job's statistics.

  procedure Set (Parameter : String := ""; Value : Integer);
  function Get (Parameter : String) return Integer;

  -- Programmatic versions of set and display
  --   initial parameters      Default      Units
  --   CPU_Scheduling           1             1 or 0 (true or false)
  --   Percent_For_Background   10          %
  --   Min_Foreground_Budget    -250        milliseconds (-5000..0)
  --   Max_Foreground_Budget    250         milliseconds (0..5000)
  --   Withhold_Run_Load        130         load * 100
  --   Withhold_Multiple_Jobs   0             1 or 0 (true or false)

```

```

-- Memory_Scheduling          1          1 or 0 (true or false)
-- Environment_Wsl           11000      pages
-- Min_Ce_Wsl                 400        pages
-- Max_Ce_Wsl                 1000      pages
-- Min_Oe_Wsl                 250        pages
-- Max_Oe_Wsl                 2000      pages
-- Min_Attached_Wsl          50         pages
-- Max_Attached_Wsl          2000      pages
-- Min_Detached_Wsl          50         pages
-- Max_Detached_Wsl          4000      pages
-- Min_Server_Wsl            400        pages
-- Max_Server_Wsl            1000      pages
-- Daemon_Wsl                 200        pages
-- Wsl_Decay_Factor           50         pages
-- Wsl_Growth_Factor          50         pages
-- Min_Available_Memory      2048      pages
-- Page_Withdrawal_Rate      1          n*640 pages/sec (n in 0..64)

-- Disk_Scheduling           1          1 or 0 (true or false)
-- Max_Disk_Load              250        Load_Factor
-- Min_Disk_Load              200        Load_Factor

-- Foreground_Time_Limit     60         seconds
-- Background_Streams         3
-- Stream_Time_N              2,5,20      minutes
-- Stream_Jobs_N              3,0,0       jobs
-- Strict_Stream_Policy       0          1 or 0 (true or false)

```

```

procedure Set_Job_Attribute (Job : Job_Id;
                             Attribute : String := "Kind";
                             Value : String := "Server");
function Get_Job_Attribute
  (Job : Job_Id; Attribute : String := "Kind") return String;

```

```

-- These interfaces exist to deal with ongoing changes to scheduler
-- characteristics without requiring new procedures.
--
-- The default parameters to Set_Job_Attributes make the indicated job
-- a server.
--
-- See the documentation for other attributes.

```

```

procedure Set_Wsl_Limits (Job : Job_Id; Min, Max : Natural);
procedure Get_Wsl_Limits (Job : Job_Id; Min, Max : out Natural);
procedure Use_Default_Wsl_Limits (Job : Job_Id);

```

```

-- Each class of job has a default for working set min and max.
-- Set_Parameter lets you change the default value. Set_Wsl_Limits lets
-- you override the default for a specific job. Use_Default_Wsl_Limits
-- restores the values to the defaults, cancelling any prior Set_Wsl_Limits
-- call.
-- Get_Wsl_Limits returns the current values for a specific job.
-- Min and Max specify the range (in number of pages) in which the
-- working set limit is set. The scheduler chooses the working set
-- limit based on prevailing conditions on the machine. If Min and
-- Max are the same, the a fixed limit is specified.
-- Min must be less than or equal to Max and Max less than the memory size.
-- Error messages are sent to an output window in the case of errors.

```

```

-- No message of any kind if success.

```

```

pragma Subsystem (Os_Commands);
pragma Module_Name (4, 3923);

end Scheduler;

```

# OBJECT MANAGEMENT SYSTEM

- **Data stack segments can be created that are not associated with a control stack. The data segments are called heaps or segmented heaps. They are used to store permanent objects in the environment.**
- **Segmented heaps are managed by a variety of mechanisms at several levels:**

**Kernel Level:** basic mechanisms for creating, copying, and deleting heaps; page differential copies; snapshot coordination.

**Object Management Level:** specific protocols and uses of heaps for storage of different types of objects; Action Semantics.

**Directory Level:** association of objects with names; versions; Ada semantics.

# OBJECT MANAGEMENT SYSTEM - HEAP MECHANISMS

- **Segmented heaps behave as Ada collections. Typed objects are placed in them using special allocator pragmas.**

```
type Heap_Ptr is access T;  
pragma Segmented_Heap (Heap_Ptr);
```

```
X: Heap_Ptr;
```

```
X:= New T;  
pragma Heap (Job_Segment.Get);
```

- **Not for customer use. Similar things can be done with Direct\_IO, Sequential\_IO, Polymorphic\_IO, etc.**
- **The first word in a heap indicates the next offset for allocation.**



# **OBJECT MANAGEMENT SYSTEM - TEMP HEAPS**

- **Internal structure similar to temp files.**
- **Each has a marker saying who owns it.**
- **Temp heaps are deleted during system boot.**
- **They usually are deleted explicitly by their creator, but this may not happen due to bugs or aborting of the creator.**
- **Recovery of this space is one reason to reboot on occasion.**

# **OBJECT MANAGEMENT SYSTEM - MANAGERS**

- **Each object manager provides some type-specific operations to support its object type. Managers generally coordinate access to their objects using an open/close protocol.**

- **Atomic transactions make it possible to read and update several objects atomically.**

- **The most significant object managers are:**

**Ada** provides storage and access for all Ada objects in the system

**File** provides storage and access for files

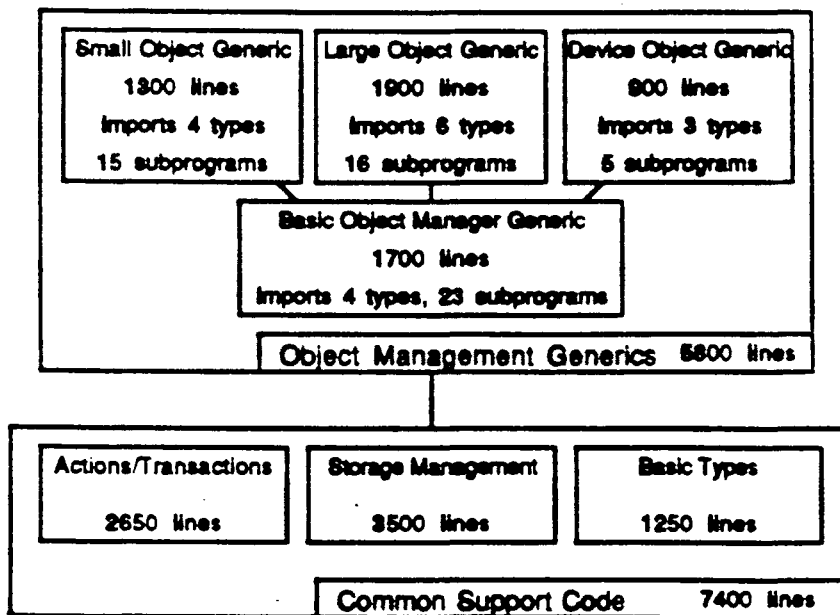
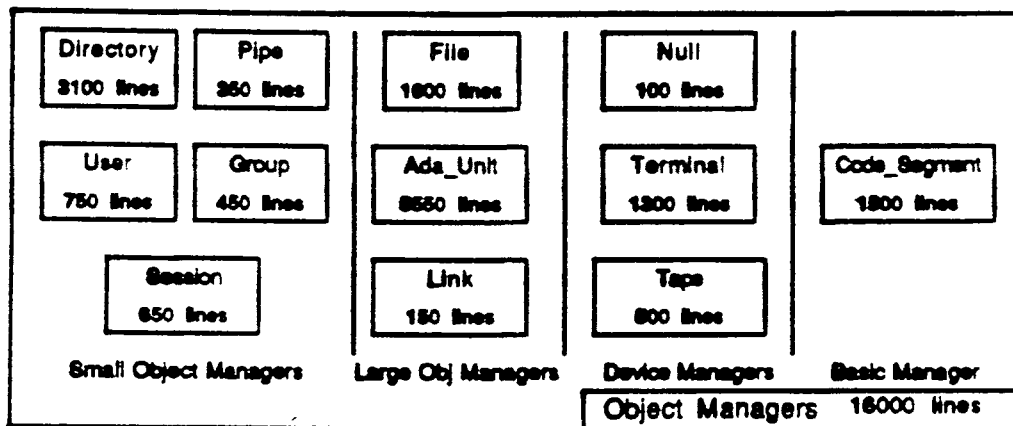
**Directory** stores mapping between string names and objects. Also stores information about objects.

- **Each manager stores its global information in a heap. During operation, new space is allocated in this heap and old space discarded.**
- **The daemon makes a copy of the heap, copying only the currently valid information. This is how compaction works.**

# **OBJECT MANAGEMENT SYSTEM - MANAGERS (CONT'D.)**

- **In addition, large objects that are managed by the object manager are each stored in their own heap or heaps.**
- **Small objects are stored in the manager heap and do not have their own heaps.**

**Figure 4-1: Object Management system structure and information.** Each of the managers in the Object Managers layer instantiates one of the Object Management Generics. The number of types and subprograms that appear as generic parameters are indicated. All code is based on the common support code.



# OBJECT MANAGEMENT - OBJECT IDs

- Each object has a unique object id.

- The object id is a triple:

`<manager id, object number, machine id>`

- The manager ids are:

0	Null	6	Session	12	Link
1	Ada	7	Tape	13	Null_Device
2	DDB	8	Terminal	14	Pipe
3	File	9	Directory	15	Archived_Code
4	User	10	Configuration		
5	Group	11	Code_Segment		

- Each manager assigns object numbers to their objects. This is usually done sequentially starting with one.
- The machine id is not really used and is always 1. (Except in a few cases.)
- Given an object id, you can find the name (if any) using:

`Action_Utilities.Display_Object (Class,Instance,Host)`

or

`Lib.Resolve ("<[Class,Instance,Host]>")`

- From the directory point of view, each version of an object has an object id and is, from the OM viewpoint, a separate object.

# OBJECT MANAGEMENT - ACTIONS

- **Actions serve two purposes:**
  - 1) **as atomic transactions that can be committed or abandoned**
  - 2) **as locks that are held on objects**
- **An action**
  - **has a unique action id**
  - **has an owner task**
  - **holds locks on zero or more objects**
  - **tracks changes in those objects**
- **When committed, the changes take effect and become "permanent" (at the next snapshot).**
- **When abandoned, the changes are all undone.**
- **Storing information to undo changes if actions are abandoned is one reason lots of disk garbage is produced doing routine things.**
- **If a task dies without explicitly committing or abandoning its actions, they are automatically abandoned the next time the action daemon runs or the next time someone attempts to update the affected object.**
- **Action "compaction" does not compact anything - it only abandons actions owned by dead tasks. The state size will only decrease when the system is shutdown and rebooted (not crashed and rebooted).**
- **To manually run the action manager daemon:**

```
Daemon.Run ("Action");
```

# OBJECT MANAGEMENT - SNAPSHOTS

- When a snapshot occurs, the state of each object manager and the action manager are saved. Each object modified since the last snapshot is flushed to disk and marked as committed.
- If the system crashes and is rebooted, during elaboration, each action that was in progress as of the last snapshot is abandoned. Thus, objects being modified when the snapshot occurred are restored to their pre-modified state.
- If there were bad objects lying around, they would be noticed at this point, probably when the action manager stopped in wait service due to a fatal error.
- During the snapshot, for part of the time, all object managers must save their state. During this time, not much else can happen. During the rest of the time, normal operations can proceed.

- To take a snapshot:

Daemon.Run ("Snap")

From EEDB

EEDB: Snap

- There is a kernel command called snapshot. Never use this. It will irrecoverably destroy the system (most likely).
- To be included, actions must be committed before the snapshot begins.

# OBJECT MANAGEMENT - BACKUP AND RECOVERY

- Backup saves address spaces in raw form. Individual objects would be difficult to recover as they are spread across segments and object manager state.
- When backup starts, a snapshot is taken. This is called a retained snapshot because it is saved until the backup completes. It is the contents of this snapshot that is saved.



# OBJECT MANAGEMENT - LOCKS

- **The Show\_Locks command can be used to display lock information and diagnose lock problems.**
- **There are several other commands to get information about locked objects:**

**What.Locks (M-F20) Gives simple information about a locked object, or objects.**

**Action\_Uilities.Lock\_Information**

**More detailed information about locked objects given either object id, string name, or Directory. Version (which is really an Object.Id)**

# OBJECT MANAGEMENT - ARCHIVE ON SHUTDOWN

- **Some software updates require changing the data structures stored by the object managers in their manager heap.**
- **"Archive on Shutdown" provides a mechanism to do this. When enabled, during shutdown, all managers convert their state to string form.**
- **On boot, the strings are converted back to internal form.**
- **Hopefully, the new software was loaded before re-booting.**
- **Operator.Archive\_On\_Shutdown (true);**

**Operator.Shutdown**

**---**

**after booting turn it off**

**Operator.Archive\_On\_Shutdown (false)**

**Daemon.Run ("Snap");**

**---**

**Op.Show\_Shutdown\_Settings displays state**

# OBJECT MANAGEMENT - LOCKS

- **There is a limit of 1024 actions. If a job starts actions and leaves them pending, the system could run out, or run low.**
- **Messages and errors "out of action ids" appear in this case.**
- **Show\_Locks can show who has what actions if it is possible to run commands. Parameters:**

**Show\_Job\_Action\_Summary**

- |                         |          |   |
|-------------------------|----------|---|
| <b>Show_All_Actions</b> | <b>-</b> | <b>show what each action has locked</b> |
| <b>Job_ID_Filter</b>    | <b>-</b> | <b>limit to this job</b>                |
| <b>Action_ID_Filter</b> | <b>-</b> | <b>limit to this action</b>             |

# OBJECT MANAGEMENT - REVIEW

- **Segmented heaps**
- **Object managers**
- **Object ids**
- **Actions**
- **Snapshots**
- **Backup and recovery**
- **Locks**
- **Commands:**

**Action\_Uilities.Display\_Object**

**Daemon.Run ("Action")**

**Daemon.Run ("Snap")**

**Show\_Locks**