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User: GPA

Object: !DELTA\_KK.REV1\_0\_0.UNITS.OM.OBJECT\_MANAGEMENT

Version: V(5) Request: 1266

Date: April 24, 1936 Queueo: 11:25:43 A/4 Printeu: 11:34:13 A/4



```
-- Export view --
-----
with Job;
with Disk:
with Error:
with Action:
with Naming;
with Object;
with Machine;
with Execution;
with Om_Services_1;
with Om_Definitions;
with Kernel_Services_1;
package Object_Management is
    -- This package defines the "standard export view" of the
    -- Object_Management subsystem (which for these purposes
    -- subsumes the Kernel functionality as well).
    -- Clients of this export view should always rename through this package.
    -- This will shield clients from the ravages of manual package integration,
    -- from a compilation point of view.
    -- Clients should ignore the following declarations:
    package P1 renames Job;
    package P2 renames Disk:
    package P3 renames Error;
    package P4 renames Action:
    package P5 renames Naming;
    package P6 renames Object:
    package Pl renames fachine:
    package 27 renames Execution:
    package F? renames to_Definitions:
    -- Clients should ignore the above declarations:
    -- Exported packages --
    ---------
    package Job renames P1;
    package Disk renames P2;
    package Link renames Om_Services_1.Link;
    package User renames Om_services_1.User;
    package view renames Om_Services_1.View;
    package Error renames P3;
    package world renames um_Services_1.World:
    package Action renames P4;
    package Namino renames 65;
    package ubject renames 25:
    package factine renames 27;
    package Trigger renames Kernet_Services_1.Trigger;
    package Execution renames P7:
    package Wait_Service renames Kernel_Services_1.Wait_Service;
    package Access_Control renames Om_Services_1.Access_Control:
    package Om_Definitions renames P8;
    package byte_String_Conversions renames
```

Kernel\_Services\_1.byte\_String\_Conversions;

end Object\_Management;

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User: GPA

Object: !DELTA\_KK.REV1\_0\_0.UNITS.OM.JOB

Version: V(33) Request: 1267

Date:

April 24, 1986

Gueued: 11:25:51 And Printed: 11:34:16 Add

(4) "Job\_yyy" is a directory;

page -- The job's garbage heaps live in this directory. Temp files created by -- Object.Create (with null name strings) go here. Rational created -- objects will follow the "\_Foo\_" convention to avoid conflicts with -- user applications, which are free to store objects in this directory. -- The system reserves the right to immediately expunge all objects in the -- directory (identified by the Garbage\_Directory function) at job -- termination, but does not guarantee to do so. It is highly recommenses -- that one never commit actions which create objects in the -- Garbage\_Directory, since this policy will guarantee that job termination -- does not leave objects in the Garbage\_Directory. -- The first function is equivalent to the second, supplying Job\_Stack and -- Top. -- Possible errors: Is\_Bad\_Version\_Handle (Status) Is\_Access\_Control\_Error (Status) function Fanaged\_Carbage\_Peab return Machine\_Segment\_Name; function tanaged\_Garbage\_Heap (Of\_Job : Om.Job\_Id) return Machine.Segment\_Name: -- This heap is created when the job starts, may be expunded when the job -- finishes, and is never "cut-back". Clients must manage their own -- garbage; that is, explicitly keep free lists. The first function -- is equivalent to the second, supplying Job. Current. This heap is created -- in the Garbade\_Directory, with simple name "\_Managed\_Garbage\_Heap\_". function Stacked\_Garbage\_Heap return Machine.Segment\_Name; function Stacked\_Garbage\_Heap (Of\_Job : Om.Job\_Id) return Machine.Segment\_Name: -- This neap is created when the job starts, may be expunged when the job -- finishes, and can be "cut-back" in a "stack-like" fashion. A program -- without tasking can use a technique where a supprogram can "mark" the -- stack on entry, and cut-back to the mark on exit. The first function -- is equivalent to the second, supplying Jop.Current. This heap is created -- in the Garbage\_Directory, with simple name "\_Stacked\_Garbage\_Heap\_". function Job\_World return Om.Object\_Handle; function Job\_World (uf\_Job : Om.Job\_Id) return Om.Object\_Handle; -- Returns the name of the directory object (a world) which contains the -- runtime representation of the identified job. This can be fed to the -- world package, which provides operations to get disk consumption -- statistics and operations to control disk space allocation. -- The first function is the same as the second if Current were supplied.

2

- -- Jon Attributes, Session invariant --
- -- These operations return the same value for every job in the -- same session.

function Session (Of\_Job : Om.Job\_Id := Current) return Om.Session\_Id; -- Same for every job in the same "session".

```
function User (Of_Job:: Om.Job_Id := Current) return Om.Jobject_Handle:
-- Object_handle of root directory of user's world.
-- Same for every job in the same "session".
function User return Om. User_Id:
function User (Of_Job : Om.Job_Id) return Om.User_Id;
-- Internal, shorter, identity of user.
function Account (Of Job : Om.Job Id := Current) return Om.Joject_Handle;
—— Object_handle of a directory with subclass Account.
-- Same for every job in the same "session".
       Stack conventions
-----
type Stack_Id is new Long_Integer:
-- type Stack_Id (For_Session : Boolean) is
       record
_ _
           case For_Session is
__
               when false =>
___
                   The_Job: Om.Job_Io:
                when true =>
                   The_Session : Om.Session_Id;
_ _
            end case;
       end record;
function Job_Stack (Of_Job : Om.Job_Id := Current) return Stack_Id;
function Session_Stack (vf_Session: Om.Session_Id := Session)
              return Stack_Id;
-- To operate on a stack associated with "your" job, simply supply the
-- result of the Job_Stack function, with default parameter. To operate
-- on a stack associated with "your" session, simply supply the result
-- of the Session_Stack function, with default parameter. Note that
___
          Job_Stack (Poot_Job (Session))
-- and Session_Stack (Current)
-- are NOT the same!
-- Subject to access control restrictions, one can operate on the stack
-- of any job or session.
subtype Stack_Location is Integer;
Top: constant Stack_Location := 0;
-- To look at the stack state that would be current following N
-- invocations of the Pop_State operation, use a Stack_Location of Top- .
-- The "depth" of the stack is the number of "pops" that can be berformed
-- before getting stack underflow.
          Job/Session state stack
function Stack_Depth (Stack : Stack_Id := Uop_Stack) return Natural;
```

function State\_Exists (Stack: Stack\_Id := Job\_Stack;

```
4
                       Location : Stack_Location := Top) return Boolean;
-- These can be used to determine the current depth of the stack, as well
-- as the existance of particular stack locations.
subtype Context is Om. Waming_Context;
function Default_Context (Stack : Stack_Id := Job_Stack;
                          Location : Stack_Location := Too)
              return Om.Pathname_String;
function Default_Context return Context;
function Default_Context (Stack : Stack_Id;
                          Location: Stack_Location) return Context;
-- The second function is equivalent to the third, supplying Job_Stack and
-- Top.
procedure set_Default_Context (The_Context : 3m.Unique_Wildcard;
                               Action : Om.Action_Id;
                               Status : out Error.Condition;
                               In_Context : Context := Default_Context;
                               Stack: Stack_Id := Job_Stack;
                               Location : Stack_Location := Top);
-- The_Context, Action and In_Context are used according to the usual rules
-- of Naming.Pesolve to compute a Version_Handle, which is then stored as
-- the default context for subsequent calls to Naming.Resolve.
-- Possible errors:
--
     Is_bau_Pathname (Status)
     Is_bad_Action (Status)
     Is_Access_Control_Error (Status)
procedure Set_Default_Context (The_Context : Context;
                               Status: out Error.Condition;
                               Stack : Stack_Id := Job_Stack;
                               Location : Stack_Location := Top);
-- Possible errors:
     Is_bad_Version_Handle (Status)
     Is_Access_Control_Error (Status)
function Name (Stack: Stack_Id := Job_Stack:
               Location : Stack_Location := Top) return String;
procedure Set_Name (New_Name : String;
                    Status: out Error.Condition:
                    Stack : Stack_Id := Job_Stack;
                    Location : Stack_Location := Top);
-- hame of the job (as displayed by What.Jobs). Always the null string
-- in the session stack. Job names longer than about 100 characters
-- are truncated at the left in order to fit within that limit.
-- Possible errors:
   ls_Access_Control_Error (Status)
```

5

```
function Termination_dessage
              (Stack : Stack_Id := Job_Stack;
              Location: Stack_Location:= Top) return String:
procedure Set_Termination_Message (New_Message : String;
                                   Status: out Error.Condition;
                                   Stack : Stack_Id := Job_Stack;
                                   Location : Stack_Location := Top);
-- When non-null, job completion will include this string in message that
-- goes to the initiator's message window. This can be used by a program
-- to display apprormal termination information. Always the null string
-- in the session stack. Termination messages longer than about 1000
-- characters are truncated in order to fit within that limit.
-- Possible errors:
-- Is_Access_Control_Error (Status)
function Default_Wait return Duration;
function Default_Wait (Stack : Stack_Id;
                      Location: Stack_Location) return Duration;
procedure Set_Default_Wait (Status : out Error.Condition;
                            Default_Max_wait : Duration := 5 * 60.0;
                            Stack : Stack_Id := Job_Stack;
                            Location : Stack_Location := Top);
-- Used as the default time to wait for object locks.
-- The first function is equivalent to the second, supplying Job_Stack and
-- Top.
-- Possible errors:
-- Is_Access_(ontrol_Error (Status)
function Max_Actions return Integer;
function Max_Actions (Stack: Stack_Id;
                      Location: Stack_Location) return Integer;
procedure Set_Max_Actions (Status: out Error.Condition;
                           Max_Actions : Integer := 25;
                           Stack: Stack_Id := Job_Stack;
                           Location : Stack_Location := Top);
function Max_Action_Locks return Integer;
function Max_Action_Locks (Stack : Stack_Id;
                           Location: Stack_Location) return Integer;
procedure Set_Max_Action_Locks (Status : out Error.Condition:
                                Max_Action_Locks : Integer := 250;
                                Stack : Stack_Id := Job_Stack;
                                Location : Stack_Location := Top);
-- hax_Actions serves to limit the number of actions which can be consumed
-- by a job. Currently, there is a limit of about 4000 concurrent actions.
-- Since commands are compiled and therefore consume actions, running the
-- system out of actions can be very painful! Each action lock consumes
```

-- system resources as well. Here the limit is the amount of disk space

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!DELTA_KK.REV1_0_0.JNITS.04.J08*V(83)
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page
-- avaitable.
-- Possible errors:
   Is_Access_Control_Error (Status)
package Profile is
    -- Currently a skeleton - to be filled in with "current spec"
   type Response_Profile is private;
   function Get (Stack : Stack_Id := Job_Stack;
                  Location : Stack_Location := Top)
                  return Response_Profile;
    procedure Set (New_Profile : Response_Profile;
                   Status: out Error.Condition;
                   Stack: Stack_Id := Job_Stack;
                   Location : Stack_Location := Top);
    -- Possible errors:
         Is_Access_Control_Error (Status)
    function Get_Default (Stack : Stack_Id := Job_Stack;
                          Location : Stack_Location := Top)
                  return kesponse_Profile;
    procedure Set_Default (New_Profile : Response_Profile;
                           Status : out Error.Condition;
                           Stack : Stack_Id := Job_Stack;
                           Location : Stack_Location := Top);
    -- Possible errors:
        Is_Access_Control_Error (Status)
private
    type Response_Profile is new Boolean: -- ha ha
end Profile;
procedure Push_State (Status : out Error.Condition;
                      Stack: Stack_Id := Job_Stack);
procedure Pop_State (Status : out Error.Condition;
                     Stack : Stack_Id := Job_Stack);
-- Pop_State throws away the top element of the identified state stack.
-- Push_State operation copies the top element into a new record, pushing
-- it on top of the state stack.
-- Possible errors:
    State_Stack_Overflow
    State_Stack_Underflow
    Is_Access_Control_Trror (Status)
         Job/Session View stack
-- This stack is operated independently of the state stack. It follows
```

-- the same conventions. An additional feature: The System view appears

-- indirect chain may be broken).

```
-- to be an unremovable entry at the bottom of the stack.
function. View_Stack_Depth
              (Stack : Stack_Id := Job_Stack) return Natural;
function View_State_Exists
              (Stack : Stack_Id := Job_Stack;
               Location : Stack_Location := Top) return Boolean;
-- These can be used to determine the current depth of the stack, as well
-- as the existance of particular stack locations.
function Default_View (Stack: Stack_Id := Job_Stack;
                       Location : Stack_Location := Top)
              return Om.Pathname_String;
function Default_View (Stack: Stack_Id := Job_Stack;
                       Location: Stack Location := Top) return Context;
procedure Set_Default_View (The_Context : Om.Unique_Wildcard;
                            Action : Om.Action_Id;
                            Status : out Error.Condition;
                            In_Context : Context := Default_Context;
                            Stack : Stack_Id := Job_Stack;
                            Location : Stack_Location := Top);
-- Possible errors:
     Is_bao_Pathname (Status)
     Is_Bad_Action (Status)
     Is_Access_Control_Error (Status) *
procedure Set_Default_View (The_Context : Context;
                            Status : out Error.Condition:
                            Stack : Stack_Id := Job_Stack;
                            Location : Stack_Location := Too);
-- These operations simply read/write the identified location in the
-- identified view stack. Note that if the stack is empty, the functions
-- will return the System view.
-- Possible errors:
     Is_Bad_Version_Handle (Status)
     Is_Access_Control_Error (Status)
function Default_korld_View (Of_World : Machine.Object_Worlds)
              return Om. Version_Id;
function Default_world_View (Of_World : Machine.Object_Worlds;
                              Stack: Stack_Id) return Om. Version_Id;
-- Scans down the identified universe view stack looking for a universe
-- view which has a non-nil slot for the given world. For this purpose,
-- treats the stack as if the System view is at the bottom. Returns the
-- Version_Io of the selected universe view. If the stack does not contain
-- such a universe view, returns Nil. This operation does NOT follow the
-- indirect chain which may lead from the slot of the returned universe
-- view. Note that a non-nil result does not imply that the referenced
-- universe view selects a world view which exists (in particular, an
```

```
-- This operation is used by the Object package to bind Version_Handle's.
-- This first function is equivalent to the second, when given Jot_Stack.
procedure Push_View_State (Status : out Error.Condition;
                           Stack : Stack_Id := Job_Stack);
procedure Pop_View_State (Status : out Error.Condition;
                          Stack : Stack_Id := Job_Stack);
-- Pop_Universe_State throws away the top element of the identified view
-- stack. Push_Jniverse_State operation copies the top element into a new
-- record, pushing it on top of the identified view stack.
-- Possible errors:
    View_Stack_Overflow
     View_Stack_Underflow
     ls_Access_Control_Error (Status)
-- Job's files --
-- Note that current input/output/error have their own independent job
-- state stacks. Job start and finish semantics for job files is
-- implemented using a coupling mechanism which is not generally available.
-- See the Job_Files package in DIO.
-- Execution context --
-- Each job is session has an execution state stack which follows the same
-- basic paradion as the job view stack. See the Execution backage for
-- details.
-- Job scheduling --
package Scheduling is
    function Kind (Cf_Jou: Om.Job_Id)
                  return Om.Job_Kind;
    -- Returns Om. Terminated for jobs which do not "exist".
    function Priority (3f_Job : Om.Job_Id := Job.Current)
                  return Om.Jot_Priority;
    function State (of_dob : Om.Job_Id := Job.Current)
                  return Om. Job_State;
    function Cpu_Consumption (Of_Job): Om.Job_ld := Job.Current)
                  return Long_Integer: -- msec
    function Disk_Consumption (Of_Job : Om.Job_Id := Job.Current)
                  return Long_Integer: -- blocks transferred
    function Current_Memory_Consumption
                  (Of_Job : Om.Job_Id := Job.Current)
```

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```
return Long_Integer; -- pages in cache
   -- Returns Nil for non existant jobs.
   procedure Enable (The_Job : Om.Job_Id:
                     Status: out Error.Condition);
   procedure Disable (The_Job : Om.Job_Id;
                      Status: out Error.Condition);
   -- Possible errors:
        Is_Access_Control_Error (Status)
   function Working_Set_Limit (Of_Job : Om.Job_Id := Job.Current)
                 return Long_Integer: -- cache pages
   procedure Set_Working_Set_Limit
                  (The_Job : Om.Job_Id := Job.Current;
                   New_Limit : Long_Integer:
                                              -- cache pages
                   Status : out Error.Condition);
   -- Fossible errors:
        Is_Access_Control_Error (Status)
   procedure Attach (Job : Um. Job_Id;
                     status : out Error.Condition);
   procedure Detach (The_Job : Om.Job_Id;
                     Status: out Error.Condition);
   -- Possible errors:
        Is_Access_Control_Error (Status)
   function Run_Load (Interval : Duration := 60.0) -- seconds
                  return Long_Integer: -- load * 100
   function Disk_Load (Interval : Duration := 60.0) -- seconds
                  return Long_Integer: -- load * 100
    function Withheld_Load (Interval : Duration := 60.0) -- seconds
                  return Long_Integer; -- load * 100
    -- more stuff to follow
end Scheduling;
-- Session control --
_____
function Exists (The_Session : Om.Session_Id) return Boolean;
-- Returns true iff there exists at least one job for which the Session
-- function matches given Session_Id parameter.
function Default_Unit (Of_Session : Om.Session_Id := Job.Session)
              return Disk.Unit_Number;
procedure Set_Lefault_Unit (Of_Session : Om.Session_Id := Job.Session;
                            Unit : Disk.Unit_Number:
                            Status: out Error.Condition);
```

generic

```
-- / value of Nil means that the system will choose the "best" unit
-- on which to run new jobs for the given session. A non-nil value
-- causes the system to always run new jobs on the specified unit (unless
-- extenuating circumstances require otherwise).
-- Possible errors :
     Is_Access_Control_Error (Status)
function Root_Job ()f_Session: Om.Session_Id := Job.Session)
              return Om. Job_Id:
-- Returns the id of the first job created under the given session. If the
-- session parameter is garbage, or the job does not exist, returns Nil.
function Subsystem_State (Of_Session: Om.Session_Id := Job.Session)
              return Om. Job_Id;
-- A session may have an associated job which is used to keep persistently
-- elaborated copies of subsystems. This function returns Nil when the
-- session does not exist or does not have such an associated job.
-- the Execution package for further details.
procedure Kitl_Session (The_Session : Om.Session_Id;
                        Status : out Error.Condition);
-- Does a Kill on every job with the given Session_Id.
-- Possible errors:
   includes those for Job.Kill on a Job_Id
procedure wait_For_Termination (The_Session : Om.Session_Id);
-- Returns immediately when The_session is garbage or otherwise does not
-- exist.
function Terminal (Of_Session : Om.Session_Id := Job.Session)
              return Om. Version_Handle;
-- Identifies the terminal to which the session is "attached",
-- possible Nil.
function Message_Window (Of_Session : Om.Session_Id := Job.Session)
              return Om. Version_Handle;
function Io_window (Of_Bession : Om.Session_Id := Job.Session)
              return Um.version_Handle;
-- Return files with subclass Message_window and Io_window, respectively.
-- Used with DID to read/write to screen via editor.
-- Job/Session creation/deletion --
-- There is no relationship between jobs other than the fact some jobs have
-- the same Session_Id. In particular, there is no notion of parent/child
-- etc.
```

-- are as follows:

page type Parameters is private; with procedure Execute (P: in out Parameters); package Initiate\_Generic is procedure Initiate (Name : String: P: in out Parameters; New\_Job : out Om.Job\_Id; Status : out Error.Condition; -- Initial job files: Current\_Output : Om.Unique\_Wildcard := ""; Current\_Error : Om.Unique\_Wildcard := ""; Current\_Input : Om.Unique\_wildcard := ""; -- New session parameters: Start\_New\_Session : Boolean := False: New\_User : Om.Object\_Handle := Om.Vil\_Object\_Handle; Users\_Password : String := ""; users\_Account : String := ""; Terminal : Om.Object\_Handle := Om.vil\_Object\_Handle; -- Miscellaneous parameters: Unit : Disk.Unit\_Number := Job.Default\_Unit; After : Duration := 0.0; Job\_Kind: Om.Job\_Kind:= Om.Attached\_Job; wait\_For\_Completion': boolean := False); -- The job's world will live on the specified unit (or "best" unit -- when Nil is specified). -- The first architectural module created in a new job is known as the -- "root thread"; there is a function to fetch this value. The -- generic formal supprogram (Execute) is run on the root thread of -- the new job. The job terminates when this root thread terminates. -- The first job created in a new session is known as the "root job"; -- there is a function to fetch this value. For interactive -- sessions, this job runs the editor. -- Setting Start\_New\_Session to true is used to run a job under a -- different Session Id. If the value of New\_User is Nil, the new -- session runs with the id of Job-User, and no password check is -- performed. Otherwise, checks password. Checks other login profile -- parameters. Updates time of last login; when the last job of the -- session terminates, will update time of last logout; also appends -- appropriate entries to the accounting log. -- Caller can implement almost arbitrary rules for the state of the -- new job by passing information via the Parameters to the Execute -- procedure (and feeding these values to the various job state Set\_ -- operations, above).

(1) when starting a new session, both the new session and its root job have state and view stacks which are copies of the

-- The built in rules for inheritance of the state and view stack

--

-- remaining cases.

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```
initiator's job stacks. This allows the session initiator
               to use its job stacks as parameters to this operation.
    _ _
           (2) when starting a new job (within an already existing
               session), the state stack, view stack and execution state
               stack of the new job are initially copies of the
    ___
               corresponding session stacks.
    -- Note that "disconnect" simply detaches an already running job,
    -- leaving it with whatever job stacks it had. The new job (created to
    -- run commands) will start with session state (as defined in rule ?
    -- above). Thus, jobs share changing state only when they explicitly
    -- manipulate session state.
    -- In any case, a Set_Name (and Push_State if necessary) is
    -- performed to set the job name as specified.
    -- The io files for the new job are always as specified by the
    -- job files parameters. Recall that setting an io file to "" will
    -- cause first reference to the file to attempt to open the
    -- session's terminal.
    -- The persistent subsystem state of sessions & jobs is specified
    -- in the Execution package.
    -- Job termination will send a non-Nil Termination_Message to the job's
    -- current error file. Recall that "termination" includes normal
    -- job completion as well as operations such as Job. Terminate_Self
    -- and Job.Kill.
    -- Possible errors:
    ___
         Is_resource_Limit_Error (Status)
             o System_Is_Out_Of_Job_World_Numbers
         Is_Bad_Start_Job_Parameter (Status)
    --
             o Bad_Job_Current_Output_Parameter
    ___
             o Bad_Job_Current_Error_Parameter
             o Sad_Job_Current_Input_Parameter
    --
         Is_Login_Error (Status)
    ___
             o _ogin_User_Id_Does_Not_Exist
             o _ogin_User_Password_Invalid
end Initiate_Generic;
procedure Kill (The_Job : Om.Job_Id;
                Status: out Error.Condition;
                Message : String := "");
-- Works on self and editors. If the caller supplies a non-null Message,
-- its value is assigned to the killed job's Termination_Message.
-- Otherwise, the job's Termination_Message will be automatically set to
       (1) "<job> has been killed"
       (2) "(job) has been killed by (user)"
-- where the value of (job) comes from the killed job's Name, and the
-- value of (user) comes from the User function of the executing job.
-- The first string is produced when the session of the killer is the same
-- as the session of the killed job. The second string is produced in the
```

13

New\_Task : out Pointer:

```
-- Possible errors:
    Is_Execution_Error (Status):
_-
         o Subsystem_ls_Still_Shared
             The specified job has elaborated copies of subsystem which
___
__
             are currently shared by other jobs.
     Is_Access_Control_Error (Status);
procedure Terminate_Self;
-- Same as Job.Kill on self, but produces no job termination message.
procedure wait_For_Termination (The_Job : Om.Job_Id);
-- Returns immediately when The_Job is garbage or otherwise does not exist.
-- Tasks within job --
_____
function My_Task_Id return Machine.Task_Id;
-- "task id" of the executing thread. Note that packages also have
-- task ids.
function Job_Ot (The_Task : Machine_Task_Id) return Om_Job_Id;
-- The job to which the task belongs. When this job terminates, the task
-- is certainly terminated.
function Root_Thread (Of_Job : Om.Job_Id := Current) return Machine.Task_Id;
-- The "task id" of the first thread created in the job. If this task
-- dies, the job will certainly be dealt a severe blow (its actions would
-- be abandoned, for example). Other than some interesting issues related
-- to allocation of tasks from persistently elaborated collections, the
-- job would probably eventually die.
function Is_Callable (The_Task : Machine.Task_Id) return Boolean;
-- As per the 'Callable attribute in the LRM. Used by the action manager
-- to determine if an action's guardian is "alive" (not alive causing the
-- action to be auto abandoned).
procedure Set_Associated_Heap (To_Heap: Machine.Segment_Name);
function Get_Associated_Heap return Machine.Segment_Name;
function Get_Associated_Heap (For_Module : Machine.Module_Name)
              return Machine.Segment_Name:
-- The above operations are used by Diana to set/get the task's current
-- Diana unit. The second Get_ operation is subject to access control.
generic
    type Task_Type is limited private;
    type Pointer is access Task_Type;
procedure Allocated_New_Noot_Thread
     (In_Job : Om.Job_Id := Job.Current;
```

## Status : out Error.Condition);

- -- This operation allows one to allocate a task in another job. The new
- -- task is considered to be an additional root thread in the target job.
- -- Termination of the target job will cause the termination of all these
- -- additional root threads.
- -- In the Gamma system the rules for task allocation (in a collection) are
- -- that the allocated task runs in the same job as the allocating thread
- -- (unless pragma Virtual\_Processor is used). In Epsilon, the pragma
- -- Virtual\_Processor is not available, and tasks cannot be allocated in
- -- other jobs without using this operation.
- -- Possible errors:
- -- Is\_Access\_Control\_Error (Status);

end Jou:

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User: GPA

Object: !DELTA\_KK.REV1\_0\_0.UNITS.KK.DISK

Version: V(30) Request: 1268

Date: April 24, 1986 Queued: 11:25:58 AM Printed: 11:35:20 AM

```
!DELTA_KK.REV1_(_0.JVITS.KK.DISK'V(30)
                                                                      page 1
with Sysi
with Machine;
with Um_Definitions:
package bisk is
    package System renames Sys;
    Nil : constant := 0;
    -- Null value for all scalar types in this package,
    -- unless stated otherwise.
    type Unit_Number is new Long_Integer range 0..15;
    -- Denotes the physical drive unit number. Note that the unit number of
    -- a drive can change when its jumpers are changed.
    type Spindle_Size is
        record
            Cylinders: Long_Integer range 0..2 ** 16 - 1;
                    : Long_Integer range 0..2 ** 8 - 1;
            Sectors
                     : _ong_Integer range 0..2 ** 8 - 1;
        ena record;
    pragma Assert (Spindle_Size'Size = 32);
    type Sector_Number is new Long_Integer range 0..2 ** 8 - 1;
    type Track_Number is new Long_Integer range 0..2 ** 8 - 1;
    type Cylinder_Number is new Long_Integer range 0..2 ** 12 - 1;
    type Address is new Long_Integer range 0..2 ** 28 - 1;
    --
          record
    --
              Cylinder: Cylinder_Number:
    __
              Track : Track_Number;
              Sector
                      : Sector Number:
    __
          end record:
    -- Corresponds to the disk address at the IOP interface
    Can_Pad_Address_To_Get_Iop_Address : constant Boolean := True;
    -- Can pad Address on the left with zeroes, and treat it as a subtype
    -- of the iop's disk address format.
    function To_Address
                  (Cylinder: Cylinder_Number:
                   Track: Track_Number:
                   Sector: Sector_Number) return Address;
    function Cylinder (Address) return Cylinder_Number;
    function Track (Ador : Address) return Track_Number;
    function Sector (Addr : Address) return Sector_Number;
    function Increment (Addr : Address; Size : Spindle_Size) return Address;
    function Decrement (Addr : Address; Size : Spindle_Size) return Address;
    -- Return the addresses of the next and prior blocks, respectively, of
    -- the given address.
    type Address_Array is array (Natural range <>) of Address;
```

2

```
type Address_Range is new Long_Integer range 0..2 ** 55 - 1;
    record
___
        First: Address:
--
        Last : Address;
--
    end record;
function To_Address_Range
              (First : Address; Last : Address) return Address_Range;
function First (The_Pange : Address_Range) return Address;
function Last (The_Pange : Address_Range) return Address;
Null_Address_Range : constant Address_Range := To_Address_Range (1, 0);
function Is_Null_Range (The_Range: Address_Range) return Boolean;
function In_Range (Addr : Address; The_Range : Address_Range)
              return Boolean;
function Ranges_Overlap (Range_1, Range_2: Address_Range) return Ecolean;
type Address_Pange_Array is array (Natural range <>) of Address_Range;
Sector_Size_In_Rytes : constant := 512;
subtype Region_Number is Cylinder_Number:
-- A region is composed of one or more adjacent disk cylinders, and is
-- denoted by the Cylinder_Number of the first cylinder in the region.
subtype Region_Count is Long_Integer range 0..2 ** 16 - 1;
pragma Assert (Region_Count'Size >= Cylinder_Number'Size);
subtype Block_Address is Address;
-- A block is composed of one or more adjacent sectors.
subtype Block_Count is Long_Integer range 0..2 ** 32 - 1;
pragma Assert (5lock_Count'Size >= Address'Size);
Block_Size_Equals_Page_Size : constant Boolean := True;
Blocks_Span_Tracks
                            : constant Boolean := False;
subtype Cluster_Address is Address;
-- A cluster is composed of one or more adjacent blocks. The cluster
-- size may vary from world to world. Note that the Volume_Manager
-- allocated clusters, not blocks.
Cluster_Size_Is_Multiple_Of_Block_Size : constant Boolean := True;
Clusters_Span_Tracks: constant Boolean := False:
Cluster_Size_Equals_Page_Size
                                       : constant Boolean := True;
-- True in the first implementation, but false in future revs.
subtype Cluster_Count is Long_Integer range 0..2 ** 32 - 1;
pragma Assert (Cluster_Count'Size >= Address'Size);
type Full_Address is new Long_Integer range 0..2 ** 32 - 1;
      record
          Unit : Unit_Number:
```

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!DELTA_KK.REv1_0_0.JNITS.KK.DISK*V(30)
                                                                      page
              Cylinder: Cylinder_Number:
    __
              Track : Track_Number:
    ___
              Sector
                      : Sector_Number;
          end record:
    function To_Full_Address
                  (Jnit: Unit_Number;
                   Cylinder: Cylinder_Number:
                   Track : Track_Number;
                   Sector : Sector_Number) return Full_Address;
    function To_Full_Address
                  (Unit : Unit_Number; Addr : Address) return Füll_Address;
    function To_Address (Addr : Full_Address) return Address;
    function Unit (Addr : Full_Address) return Unit_Number;
    function Cylinder (Addr : Full_Address) return Cylinder_Number;
    function Track (Addr : Full_Address) return Track_Number;
    function Sector (Addr : Full_Address) return Sector_Number;
    function Increment (Adar : Full_Address;
                        Size : Spindle_Size) return Full_Address;
    function Decrement (Adar : Full_Address;
                        Size : Spindle_Size) return Full_Andress;
    -- Return the addresses of the next and prior blocks, respectively, of
    -- the given address.
    type Full_Address_Array is array (Natural range <>) of Full_Address;
    subtype Full_flock_Address is Full_Address;
    subtype Full_Cluster_Address is Full_Address;
    All_Dirty_World_Info_Mapped_Under_Its_Vp : constant Boolean := True;
    -- This invariant is fundamental to the window of vulnerabilty mechanism,
    -- actions, and crash recovery. Therefore, one should take great care
    -- with the world parameter in the following function.
    -- Be extra careful when using "special worlds" (like Machine.Om_World and
    -- Machine.Miscellaneous_World).
    function To_Disk_Mapping_Address
                  (Addr : Full_Address;
                   world: Machine.World_Number) return Machine.Page_Address;
    -- All unwired kk_om data structures are in cache pages whose vp
    -- corresponds to the world in which the data structure lives. Data
    -- structures which are not implemented in heaps use this function to map
    -- their disk adoresses to cache pages.
    -- Converts the disk address to a page address as follows:
    ___
         ΰ
                         o wits
          unit
                         4 pits
                                   --> segment
    ___
          cylinaer
                         12 pits
          world param
                        10 pits
                                 ---> vp
          track
                        & bits \
          sector
```

3 pits

3 bits

Ű

(

--> page number

10 bits ---> bit offset

--

. . .

```
3 bits ---> segment kind (data)
                     64 bits
___
-- Note that this mapping implies that ucode must not allocate module
-- names with segment numbers smaller than 2 ** 16 (ie, 1/64 of the
-- name space is preallocated by CM).
-- Note that it is possible for the same Full_Address to be mapped by many
-- different Fage_Addresses, each with a different World_Number component.
-- By convention, DON'T DO THAT!
function Is_Disk_Mapping_Page (Addr : Machine.Page_Address) return Boolean;
-- Returns true iff the most significant 6 bits of the address are 0 and
-- the least significatnt bits of the address are 4 (kind).
function To_Disk_Address (Addr : Machine.Page_Address) return Full_Address;
-- Inverts the above mapping. Returns a bogus disk address if the given
-- page address is not a disk mapping page (the bogus result is guaranteed
-- to cause an error when used to issue a physical disk ID).
type Region_Id is new Long_Integer range 0..2 ** 10 - 1;
--
     record
               : Unit_Number:
         Region : Region_Number:
     end record:
function To_Region_la (Unit : Unit_Number;
                       Region: Region_Number) return Region_Id;
function Unit (Id : Region_Id) return Unit_Number;
function Region (Id: Region_Id) return Region_Number;
function Base_Address (Of_Region : Region_Id) return Address;
function Full_Base_Address (Of_Region : Region_Id) return Full_Address;
-- Returns an address with same cylinder #, and zero for track and sector.
type Block_Kinds is new Long_Integer range 0...2 ** 8 - 1;
-- "Standard header" for disk blocks:
     record
         Kind
                  : block_Kinds;
__
                   : Filler.Filler_8;
         Instance : Um_Definitions.Unique_Id;
___
                   : Full_Address:
         Home
         -- the above fields consume the first 96 bits.
_-
-- "Semi-standaro header" for disk blocks:
__
     record
_ _
                  : Filler.Filler_64;
         . . .
                   : Filler.Filler_64;
___
         • • •
                   : Block_Kinds;
         Kind
                   : filler.Filler_8;
_ _
         Instance : Um_Definitions.Unique_Id;
___
                   : Full_Address;
         -- the above fields consume the first 128 + 95 bits.
```

```
-- This convention is fundamental to the ability to detect garbage in
-- kernel data structures. It also greatly simplifies one's ability to
-- poke around at disk data structures, and make sense out of them.
-- The Pack_Label has a non-standard header, for historical reasons.
-- The List_Section_block and Object_Header_Block have semi-standard
-- headers for compatability with ucode interpretation of the first
-- word in a heap. Unless stated otherwise, all other block kinds
-- have standard headers. Block_Kinds takes on the following values:
Size_Of_The_Standard_Block_Header : constant :=
    Block_Kinds'Size + 3 +
    Um_Definitions.Unique_Id'Size +
    Full_Address'Size;
Uninitialized_Block
                                    : constant Block_Kinds := 15#0#;
-- Block_Kind has not yet been assigned to the block.
Retarget_Database_Block : constant Plock_Kinds := 16#1#;
-- Individual block in the retarget database.
World_Allocation_Elock
                                     : constant Block_Kinds := 16#2#;
-- Map from world_Number to location of world_deader_Block.
World_Header_Block
                                     : constant Block_Kinds := 15#3#;
-- At the beginning of the root region of an existing world.
List_Section_Block
                                     : constant Block_Kinds := 16#4#;
-- Used to construct lists of disk addrésses.
-- Uses "semi-stan and header".
Object_Header_Slock
                                    : constant Block_Kinds := 15#8#;
-- Contains ucope control, KK_OH control, and initial user data of heap.
-- Uses "semi-standard header".
Free_Object_Header_Stock
                                    : constant Block_Kinds := 16#9#;
-- An unused object header block.
-- Uses "semi-standard header".
Machine_Segment_Index_Block : constant Block_Kinds := 16#10#;
-- Index_Block for a segment created by the architecture.
Heap_Segment_Index_Block
                                     : constant Block_Kinds := 15#11#;
-- Index_Block for a segment underlying a permanent object.
Object_Catalog_Branch_:lock : constant Block_Kinds := 16#20#;
Object_Catalog_Leaf_block : constant Block_Kinds := 15#21#;
Job_Catalog_Branch_Block : constant Block_Kinds := 16#22#;
Job_Catalog_Leaf_Block : constant Block_Kinds := 16#23#;
-- Biree nodes for the segment catalogs.
Structural_Kap_Branch_Glock : constant Block_Kinds := 15#30#;
Structural_Map_Leaf_Glock : constant Block_Kinds := 16#30#;
-- Etree nodes for directory's structural map.
                             : constant Block_Kinds := 16#32#;
Version_Map_branch_Stock
Version_Map_Leaf_block
                                     : constant Block_Kinas := 15#33#;
```

-- biree nodes for directory's version map.

```
: constant Block_Kinds := 16#34#;
Version_Map_Data_Flock
-- Contains information about particular versions of objects.
Link_Pack_Branch_Block_1 : constant Block_Kinds := 16#36#:
Link_Pack_Leaf_block_1 : constant Block_Kinds := 16#37#;
-- Btree nodes for the reconstructed portion of the Link_Pack state.
Link_Pack_Branch_Elock_2 : constant Block_Kinds := 16#38#;
Link_Pack_Leaf_Block_2 : constant Block_Kinds := 16#39#;
-- Btree nodes for the temporary portion of the Link_Pack state.
                                      : constant Block_Kinds := 16#40#;
Action_Log_Head_Block
Action_Log_Block : constant Block_Kinds := 16#41#;
-- Used to build action logs.
```

end Disk:

```
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      qqqa
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      p 5
  Ĝ
             A A
      Р
          c
G
             A
                 Ą
G GG
      9999
             AAAAA
6 6
      Ł,
             Α
                 A
G
   G
      F
             A
                 Α
GGGG
             Α
                 A
```

User: GPA

Object: !DELTA\_KK.REv1\_0\_0.UNITS.KK.ERROR

Version: V(51) Request: 1269

Date: April 24, 1986

Queued: 11:26:05 AM Printed: 11:35:46 Am

1

```
package Error is
    type Status is new Long_Integer range û..2 ** 16 = 1;
    -- Vanilla status code
    type Condition_Information is new Long_Integer range 0..2 ** 48 - 1;
    type Condition is new Long_Integer:
    -- Status plus info which is a function of the status
    -- record
          Info
                : Condition_Information;
          Reason: Status;
    -- end record;
    function Info (The_Condition: Condition) return Condition_Information;
    function Reason (The_Condition : Condition) return Status;
    function To_Condition (Info : Condition_Information;
                           Reason: Status) return Condition;
    ______
    -- Success --
    Successful_Status : constant Status := 0;
    Successful
                     : constant Condition := 0; '
    function Is_Successful (The_Status : Status) return Boolean;
    function Is_Successful (The_Condition: 'Condition) return Boolean;
    -- Returns true for the following values of Status:
       Successful
    -- It is considered 3% for clients to use "=" to compare either a
    -- Status or Condition to the value Successful, instead of calling the
    -- Is_Successful function, which cannot be inlined due to compiler/debugger
    -- restrictions.
    -- Bad Object_ld/Version_ld (010x) --
    function Is_bad_Version_ld (The_Status : Status) return Boolean;
    function Is_Bad_Version_Id (The_Condition : Condition) return Boolean;
    function Is_Bad_Object_Id (The_Status : Status) return Boolean;
    function Is_Bad_Object_Id (The_Condition : Condition) return Boolean;
    function Is_Bad_Version_Handle (The_Condition: Condition) return Boolean;
    function Is_bad_Object_Handle (The_Condition: Condition) return Boolean;
                                   : constant Status := 16#0101#;
    world_Does_Not_Exist
                                  : constant Status := 16#0192#;
    Object_Class_Does_Not_Exist
    Object_Index_Does_Not_Exist : constant Status := 16#0103#;
    Object_Instance_Does_Not_Match : constant Status := 16#0104#;
    Object_Version_Does_Not_exist : constant Status := 16#0105#;
    No_world_View : constant Status := 16#0106#;
                                   : constant Status := 16#0107#;
    Object_Mot_In_view
                                   : constant Status := 16#0103#;
    Universe_Does_Not_Exist
                                  : constant Status := 16#0109#;
    Max_Indirection_Exceeded =
```

```
-- Bad pathname (011x) --
function Is_bad_Pathname (The_Status : Status) return Boolean;
function Is_Bad_Pathname (The_Condition : Condition) return Boolean;
Filename_Syntax_Error
                               : constant Status := 16#0111#;
      Condition identifies the character in the name which caused the
--
      error, as well as the token class which was expected.
Child_Object_Does_Not_Exist : constant Status := 16#0112#;
     Condition identifies the segment (of the name) which caused the
--
      error. The prefix to the left of this segment resolved without
     error. The identified segment does not identify any child object.
Link_Pack_Entry_Does_Not_Exist : constant Status := 16#0113#;
--
     Condition identifies a segment (of the name) which was resolved by
__
      looking through the link pac, which failed.
Child_Name_Is_Ambiguous : constant Status := 16#0114#;
      Similar to the previous status, but the identified name segment
      identifies more than one child object.
Set_Causes_Ambiguity : constant Status := 16#0115#;
      Condition identifies a segment (of the name) which uses set
      notation to identify more than one object/version.
Indirect_File_Causes_Ambiguity : constant Status := 16#0116#;
      Condition identifies a segment (of the name) which identifies an
      indirect file which identifies more than one object/version.
-- Ead action (012x) --
function Is_Bad_Action (The_Status : Status) return Boolean;
function Is_Bad_Action (The_Condition: Condition) return Boolean;
Action_Does_Not_Exist
                                : constant Status := 15#0121#;
Not_The_Action_Gwner
                                : constant Status := 16#0122#;
Commit_Has_Been_Frevented
                                : constant Status := 15#0123#;
Action_Already_Has_Update_Locks : constant Status := 16#0124#;
Update_Has_Been_Prevented
                            : constant Status := 16#0125#;
Action_Owner_Not_Willing
                                : constant Status := 16#0126#;
Cant_Xfer_Ownership_Cross_Jobs : constant Status := 16#0127#;
Action_Owner_Not_Dead
                                : constant Status := 16#0128#;
Master_Action_Does_Not_Exist
                               : constant Status := 15#0129#;
                              : constant Status := 15#012A#;
Action_Is_A_Dependent_Action
Action_Already_Has_Dependent
                               : constant Status := 16#0123#;
Action_Guardian_Not_Dead
                                : constant Status := 16#012C#;
-- Lock errors (013x) --
------
function Is_Lock_Error (The_Status : Status) return Boolean;
function Is_Lock_Error (The_Condition: Condition) return Boolean;
Cant_Get_Object_Lock
Cant_Get_Record_Lock
                                : constant Status := 16#0131#;
Cant_Get_Record_Lock : constant Status := 16#0132#;
Cant_Get_View_Stot_Lock : constant Status := 16#0133#;
Cant_Get_Lock_On_Predecessor : constant Status := 16#0134#;
```

```
Ittegat_Lock_Mode
                                : constant Status := 16#0135#;
Invalid_Object_Lock_mandle
                               : constant Status := 16#3136#;
                                : constant Status := 16#0137#;
Object_Is_Not_Locked
 ______
-- Load_Image errors (014x) --
: constant Status := 16#0141#;
Load_Image_Is_Too_Big -
Not_A_Secondary_Load_Image : constant Status := 16#0142#;
-----
-- Code segment errors (015x) --
-------
                               : constant Status := 16#0151#;
Not_A_Cross_Cg_Segment
Not_In_Cross_Cg_Object_World : constant Status := 16#0152#;
                                : constant Status := 16#0153#;
Not_A_New_Code_Segment
_____
-- job/session errors (016x) --
------
Bad_Stack_Id : constant Status := 16#0160#;
Bad_Job_Id : constant Status := 16#0161#;
Bad_Session_Id : constant Status := 16#0162#;
                       : constant Status := 16#0163#;
: constant Status := 16#0164#;
State_Stack_Overflow
State_Stack_Underflow
                              : constant Status := 16#0165#;
View_Stack_Overflow
                               : constant Status := 16#0166#;
View_Stack_Underflow
Invalid_Stack_Location
                                : constant Status := 16#0167#;
-- Other identification errors (02xx) --
                                : constant Status := 16#0221#;
Unit_Does_Not_Exist
Parent_Object_Does_Rot_Exist : constant Status := 16#0222#;
Predecessor_Version_ls_Not_Known : constant Status := 16#0223#;
                                : constant Status := 16#0224#;
Segment_Does_Not_Exist
                                : constant Status := 16#02a0#;
world_Already_Exists
                                : constant Status := 16#02A1#;
Link_Pack_Entry_Already_Exists
-- Link Pack Errors (02Bx) --
function Is_Link_Error (The_Status : Status) return Boolean;
function Is_link_Error (The_Condition: Condition) return Boolean;
Link_Name_Joes_Not_Exist : constant Status := 16#0280#;
                            : constant Status := 16#0201#;
Cant_Delete_Internal_Link
Duplicate Link Name
Duplicate_Link_Name : constant Status := 16#02B2#;
External_Link_To_Local_Object : constant Status := 16#02B3#;
Cant_Create_Internal_Link : constant Status := 16#02B4#;
Cant_Create_External_Link : constant Status := 16#0285#;
Cant_Delete_External_Link : constant Status := 16#0286#;
```

```
Cant_Identify_Link
                             : constant Status := 16#0267#;
-- View Errors (020x) --
function Is_View_Error (The_Status : Status) return boolean;
function Is_View_Error (The_Condition : Condition) return Boolean;
Dad_Object_Index_Constraint
                                     : constant Status := 15#02D0#;
-- Miscellaneous object operation errors (03xx) --
Illegal_Class : constant Status := 16#0301#;
Object_Mismatch : constant Status := 16#0302#;
Object_Class_Mismatch : constant Status := 16#0303#;
Too_Many_Versions : constant Status := 16#0304#;
Version_Is_Still_Referenced
                                      : constant Status := 16#0305#;
Multiple_Versions_Not_Supported
                                     : constant Status := 16#0336#;
world_View_Is_Frozen : constant Status := 16#0307#;
Version_Is_Frozen : constant Status := 16#030°#;
world_Still_Contains_Objects : constant Status := 16#0309#;
World_Still_Contains_Subworlds

**Trandy In Use
                                     : constant Status := 16#030A#;
Object_Index_Already_In_Use
                                      : constant Status := 16#030B#;
Explicit_Version_Creation_Not_Allowed : constant Status := 16#030C#;
Object_Already_Exists : constant Status':= 16#030D#:
Object_Is_Not_Reserved
                                      : constant Status := 10#0305#;
Must_Update_via_View : constant Status := 16#030F#:
-- Resource limit errors (04xx) --
function Is_Resource_Limit_Error (The_Status: Status) return Boolean;
function Is_Resource_Limit_Error (The_Condition : Condition) return boolean;
Object_Is_Out_Of_Version_Numbers
                                          : constant Status := 16#0401#;
World_Is_Out_Of_Segment_Numbers
World_Is_Out_Of_Code_Segment_Numbers
                                          : constant Status := 16#0410#;
world_Is_Past_Low_Space_warning_Threshold : constant Status := 16#0413#;
Insufficient_Object_Indices_Available : constant Status := 16#0414#;
Unit_Is_Past_Low_Space_Warning_Threshold : constant Status := 16#0420#;
Unit_Is_Out_Of_Disk_Space : constant Status := 16#0421#;
System_Is_Out_Of_Action_Numbers : constant Status := 16#0430+;
System_Is_Out_Of_Job_World_Numbers : constant Status := 16#0431#;
System_Is_Out_Of_Object_world_*umbers : constant Status := 16#0432#;
System_Is_out_Of_Session_Aumbers : constant Status := 16#0433#;
```

Jou\_Is\_At\_Action\_Limit : constant Status := 15#0440#;

```
Job_Is_At_Action_Node_Limit
                                          : constant Status := 15#6441#;
-- Access control errors (05xx) --
function Is_Access_Control_Error (The_Status: Status) return Boolean;
function Is_Access_Control_Error (The_Condition: Condition) return Boolean;
-- Physical I/O errors (101x) --
                              : constant Status := 16#1011#;
Page_Does_Not_Fit_In_Cache
                          : constant Status := 16#1012#;
Page_Is_Not_In_The_Cache
                                 : constant Status := 15#1013#;
Page_Is_Already_wired
Page_Is_Already_Loading
                                 : constant Status := 15#1014#;
Unrecoverable_Disk_Error
                                  : constant Status := 15#1015#;
-- KK/OM errors (Dxxx) --
-- Status conditions for internal KK/OM interfaces (Exxx) --
                                  : constant Status := 16#2001#;
Page_Is_Already_Defined
Page_Is_Atready_Defined : constant Status := 15#E002#;
Page_Is_Atready_Shadowed : constant Status := 15#e003#;
Header_Already_Stores_In_Free_List : constant Status := 15#E004#;
-- Status conditions that are not even visible from --
-- internal KK/Om interfaces (Fxxx)
```

end Error;

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              Ant
         )
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G
              A A
       P
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                   A
               Ą
G
       PPPP
  66
              AAAAA
G
       Р
   ú
               A
                   A
       P
G
   Ĝ
               Α
                   A
 6666
       Р
               Α
                   A
```

User: GPA

Object: !DELTA\_KK.REV1\_0\_0.UNITS.OM.ACTION

Version: V(93) Request: 1270

Date: April 24, 1986 Queued: 11:20:12 A4

Printed: 11:30:07 AM

```
page
with Job:
with Sysi
with Error;
with Machine;
with Version_Map: -- private eyes
with Om_Definitions;
pragma Private_Eyes_Only;
with Action_Layout;
package Action is
    -- pragma Subsystem (Object_Management);
    -- pragma Module_Name (4, ?);
    package System renames Sysi
    package Om renames Om_Definitions;
    subtype Version_Id is Jm. Version_Id;
    subtype Action_Id is Om.Action_Id;
    subtype Lock_Mode is Om.Lock_Mode;
    subtype Record_Key is Om.Record_Key;
    function Max_Length_Of_Log_Record return Natural;
    subtype Log_kecord_Value is System.Bit_String;
    -- Basic action operations --
    procedure Start_Action (Id : out Action_Id;
                            Status: out Error.Condition;
                            Upoate_Prevented : Boolean := False;
                            Commit_Prevented : Boolean := False;
                            Guardian : Machine.Task_Id := Job.Root_Thread;
                            master : Action_Id := Om.Nil);
    -- Used to start an action. Update_Prevented = true is equivalent to
    -- calling Prevent_Update. Similarly, Commit_Prevented = true is
    -- equivalent to calling Prevent_Commit. Note that Update_Prevented
    -- makes starting and finishing the action considerably faster (since
    -- no 10 is required to write commit records).
    -- The calling task is said to be the "owner" of the action. One will
    -- observe that only the owner can operate on the action (but there is
    -- a mechanism for transferring ownership).
    -- If both the owner and the guardian die (become not *callable), the
    -- action will eventually get auto-abandoned. The default value for
    -- the Guardian parameter causes actions to almost never get auto abandoned
    -- (because death of the root thread implies termination of the job, and
    -- job termination implicitly invokes Abandon_All_Actions_In_Jou).
    -- Specifying a Nil Master causes the new action to be a master.
    -- Specifying a non-kil Haster causes the new action to be a dependent of
    -- the named Master, and to become commit prevented. When the master is
```

-- finished (by explicit Commit, Abandon, or auto-abandon), its associated -- dependent is Abancon's. The guardian and owner of a dependent action

```
-- are those of its master.
-- Possible errors:
--
     Is_Bad_Action (Status):
--
         o Master_Action_Does_Not_Exist
               The identified Master does not exist.
___
         o Not_The_Action_Owner
_ _
               Not the owner of the Master action.
--
         o Action_ls_A_Dependent_Action
               The Master parameter refers to a dependent action.
--
         o Action_Already_Has_Dependent
--
               Master's can have at most one dependent.
--
    Is_Resource_Limit_Error (Status):
--
         o System_Is_Out_Of_Action_Numbers
--
         o Job_Is_At_Action_Limit
__
         o Job_Is_At_Action_Node_Limit
-- Internal operations --
procedure Acquire_Object_Lock
               (Id : Action_Id;
               On_Version : Version_Id:
               Map_Location : Version_Map.Data_Item_Pointer;
               In_Mode : Lock_Mode:
               For_New_Version : Boolean:
               Status: out Error.Condition;
               Max_kait : Duration := Duration'Last);
type Object_Lock_Handle is private;
Nil_Object_Lock_Handle : constant Object_Lock_Handle;
procedure Acquire_Record_Lock_Intention
               (Id : Action_Id:
                Un_Version : Version_Id;
                Map_Location : Version_Map.Data_Item_Pointer;
                In_Mode : Lock_Mode;
                Handle: out Object_Lock_Handle;
                Status: out Error.Condition;
               Max_Wait : Duration := Duration*Last);
procedure Acquire_Record_Read_Lock
               (Id : Action_Id;
                Un_Version : Object_Lock_Handle;
                On_Item : Record_Key;
                Status: out Error.Condition;
                Max_Wait : Duration := Duration'Last);
generic
    with function Provide_Log_Entry return Log_Record_Value;
procedure Acquire_Necond_Write_Lock
     (Id : Action_ld;
      On_Version : Ouject_Lock_Handle;
      Un_Item : Recond_Key:
      Status : but Error.Condition;
```

dax\_wait : Duration := Duration\*Last);

- -- These operations are used by an action to acquire a lock.
- -- The Acquire\_Object\_Lock and Acquire\_Record\_Lock\_Intention operations
- -- are used to acquire object level locks. The second operation must be
- -- used prior to acquiring record locks.
- -- The Acquire\_Record\_Kead\_Lock and Acquire\_Record\_Write\_Lock operations
- -- are used to acquire record locks. The client passes in the
- -- Object\_Lock\_Handle which was returned by a previous call to
- -- Acquire\_Record\_Lock\_Intention (which locked some object "X"). In the
- -- case of Acquire\_Record\_Read\_Lock, the call gets a Record\_Read lock on
- -- the record (specified by On\_Item) in object "X".
- -- In the case of Acquire\_Record\_Write\_Lock, the call gets a Record\_Write
- -- lock on the record (specified by On\_Item) on object "X". In addition,
- -- once the Record\_write lock has been acquired the client must supply
- -- the "log record", via Provide\_Log\_Entry, which will be stored in the -- action los, and processed in order to commit/abandon the action.
- -- The Provide\_Log\_Entry operation should not acquire any action locks
- -- or mutex's, else you risk getting the action manager deadlocked.
- -- The appropriate calls to the Heap\_Segment\_Manager are made. In
- -- particular, For Acquire\_Object\_Lock: if For\_New\_Version is true, then
- -- calls Create\_New\_Version; else calls Open\_For\_Read or
- -- Spawn\_New\_Generation, as appropriate. Similarly, for
- -- Acquire\_Record\_Lock\_Intention, calls Open\_For\_Read or
- -- Spawn\_New\_Generation, as appropriate: Assumes caller does NOT
- -- currently hold the mutex in the supplied version map entry. Will
- -- acquire and release the mutex once the object has been locked.
- -- Modifies fields of the segment descriptor and version map as specified
- -- by the Heap\_Segment\_Manager.

## -- Lock compatibility matrix:

 		Current Lock (by different action) RO WO RRO RWO RR RW						SR0	U	
	RO :	Х		х					X	
	WO							1	X	
Desired Lock	RRO	X		X	<b>X</b>			X	X	
	E #3			X	X			1	X.	
	RF :					X		<b>X</b> ,	, <b>X</b>	
	15.#×							1	Y	
	- <b>S</b> RO	, X	1	X	1	X	1	<b>X</b>	)	
	; ;	X	· 大	X	×	Х	, X	X	χ	

-- Absence of an "X" indicates that the desired access will not be granted

-- if any UTHER action (not including requesting action) has the indicated

-- current access. The 1 indicates that while this package does not

-- allow compatibility, the Object package will "do the right thing" to -- allow users the view that they are compatible. Via Max\_Wait, queueing -- is available when access is denied.

## -- Lock upgrade matrix:

 	Current Lock (by same action)									
		₹0	¥0	R30	RWO		RW	SRO	U	
	RO	: x	*	X				X	X	
	WO	<b>* *</b> 2	*					X	X	
Desired Lock 	RHO	. x		X	X			x	X	
	<b>RNO</b>	•		X	X			x	X	
	КÐ					X	*	Х	X	
	ř.	•				<b>*</b> 2	*	У	x	
	SRO	i L	χ	Χ	λ	<b>X</b>	X	x	<b>y</b>	
	U	: X	х	X	X	X	X	X	X	

-- Absence of an "X", "\*", or "\*2" indicates that the desired access
-- will not be granted. Presence of one of these indicators indicates
-- that the desired access will be granted, even though there are other
-- actions waiting for incompatable access; but note that an action can
-- upgrade from RO to wO iff it is the only action with RO. The three
-- positions with "\*"s identify the actual upgrade cases. Note that there
-- are "X"s in the same positions as in the compatibility matrix.

The semantics of Supersedeable\_Read\_Object are as follows: Obviously, if the version in question does not undergo updates while the s-reader has the version, the s-reader sees the same value as a vanilla reader. The remaining cases: (1) The version is open for update when the s-reader arrives. In this case, reader sees the writer's before image, and continues to see it even after the writer commits. The implementation will cause the reader to make a full copy of the before image at open time. (2) The version is already open by an s-reader when the write arrives. The reader continues to see the same image. The writer simply makes a new version, in a fashion similar to the implicit versions caused by various version control policies. Implementation note: most of the implementation of supersedeable read actually happens in the Object package.

The semantics of Unsychronized are as follows: If the version is not already open, the u-opener will have no access. If the version is already open, the u-opener will have access to the same generation as the other opener's. Take note that this means that the u-opener has access to the same generation as a writer, and can in fact make changes to the writer's image. When the version is closed by all of the other opener's, the u-opener will have no access. The u-opener should be capable of dealing with Non\_Existent\_Page\_Error at any time, since access may go away at any time. The above semantics may change from software release to software release, so don't write lots of code that

```
-- depends on it. In particular, this mode should NOT be directly
-- available to customers (exception: ok when shielded by the IO
-- packages).
-- In some cases, the action manager considers the action to be in
-- deadlock, or near deadlock, and returns Lock_Error immediately, even
-- thought the acquire operation gave a non-zero Max_Wait. Example:
-- Actions A1 and A2 hold R0 locks on object X; action A2 attempts to
-- upgrade to w0, but can't because of the RO held by A1; action A2
-- will get Lock_Frror immediately. Another example: Another example:
-- Action A holds a wo lock on object X and attempts to get a Rwo lock
-- on X with non-zero Max_Wait: it will get Lock_Error immediately.
-- ***WARNING*** The presence of the upgrades marked by "*" allows for
-- strange behaviour. 2 scenarios are described below:
-- Consider: task T1 uses action A to get a record write lock on some
-- record λ; task T1 passes ownership to task T2 which proceeds to also
-- get a record write lock on the same record X: tasks T1 and T2 each
-- proceed to operate on the record concurrently; unless these tasks are
-- prepared for this, one may get strange results. Tasks can protect
-- themselves via the following technique: acquire the lock (which may be
-- via upgrade) and do not relinquish ownership of the action until after
-- leaving the "critical region".
-- Consider a example which involves just a single task: Procedure X has a
-- handle for action A, gets a "file handle" for write on object O, does
-- work (like building up context in the file), and calls some procedure
-- Y, passing it A. The exact effect of procedure Y is not precisely
-- known by X. Y gets yet another "file handle" for write on object 0
-- (this is allowed by the upgrade semantics), and completely changes the
-- contents of object 0. On return, X gets very confused. The only known
-- solution is for applications to understand whats going on.
-- These operations are only available to the OM subsystem, and to other
-- "registered" environment clients (like the IO packages). In particular,
-- these operations are not available to "users".
-- Possible errors:
     Is_Bad_Action (Status):
         o Action_Does_Not_Exist
--
         o Not_The_Action_Owner
         o Update_Has_Been_Prevented
     Is_Lock_Error (Status):
         o Cant_Get_Object_Lock (1st & 2nd ops only)
__
--
         o Cant_Get_Record_Lock (3rd and 4th ops only)
--
         o Illegal_Lock_Mode
              The Acquire_Object_Lock operation only takes modes
              Read_Object and Write_Object. The
--
              Acquire_Pecond_Lock_Intention operation only takes modes
              Record_Read_Object and Record_Write_Operation.
     Is_Resource_Limit_Error (Status):
         o Job_Is_At_Action_Node_Limit
procedure Acquire_Record_Read_Lock_And_Object_Lock
               (Id : Action_Id;
```

Record\_Read\_Lock\_Version : Object\_Lock\_Handle;

On\_Item : Record\_Key;

```
Object_Lock_Version : Version_Id;
                Map_Location : Version_Map.Data_Item_Pointer;
                Object_Mode: Lock_Mode:
                Status: out Error.Condition;
                Max_wait : Duration := Duration*Last);
procedure Acquire_Record_Read_Lock_And_Intention_Lock
               (Id : Action_Id;
                Record_Read_Lock_Version : Object_Lock_Handle:
                On_Item : Record_Key:
                Object_Lock_Version : Version_Id;
                Map_Location : Version_Map.Data_Item_Pointer;
                Object_Mode: Lock_Mode;
                Handle: out Object_Lock_Handle;
                Status: out Error.Condition;
                Max_Wait : Duration := Duration*Last);
-- These operations are a composite of Acquire_Record_Read_Lock and
-- Acquire_Object_Lock or Acquire_Record_Read_Lock and
-- Acquire_Record_Lock_Intention. Either both locks are acquired or
-- neither is.
procedure Append_Log_Entry (Id : Action_Id;
                            On_Version : Version_Id;
                            On_Item : Record_Key;
                            Log_Record : Log_Record_Value:
                            Status: out Error.Condition);
-- In some special cases, additional information is needed to abandon an
-- action. If On_Item is nil, the action must have a Write_Object lock
-- on the version. If On_Item is not nil, the action must have a
-- Record_Write_Object lock on the version and a Record_Write lock on
-- On_Item. Currently, this operation is only used by load image's.
procedure Re_Open_Segment (On_Version : Version_Id;
                           Map_Location : Version_Map.Jata_Item_Pointer;
                           Status : out Error.Condition);
-- Due to the internal workings of commit in the face of multiple writers
-- on the same object, tasks will occasionally handle
-- Write_To_Read_Only_Page conditions by calling this routine.
-- This operation is available only to CM subsystems.
-- Possible errors:
     Version_Is_Frozen
procedure Abandon_All_Actions_In_Job (The_Job : Machine.Job_Worlds;
                                      Status: out Error.Condition);
-- Used by Job termination to abandon every action created and still open
-- by the specified job. Assumes the job's modules have already been
-- terminated.
-- This operation is available only to (of subsystems.
-- Possible errors:
     Is_bad_Action (Status):
```

-- Possible errors:

```
o Action_Owner_Not_Dead
                o will get this if there exists a task (in The_Job, other
___
                  than the thread running this procedure) which still
                  claims to be owner of an action owned by The_Job.
         o Action_Guaroian_Not_Dead
                o will get this is the guardian task is not dead.
-- Basic action operations --
procedure Commit (Id : Action_Id;
                  Status : out Error.Condition);
-- None of the updates of this action actually take place until Commit
-- is called.
-- Possible errors:
--
     Is_bad_Action (Status):
         o Action_Does_Not_Exist
         o Not_The_Action_Owner
         o Commit_Has_Been_Prevented
         o Action_Is_A_Dependent_Action
--
procedure Abandon (Id : Action_Id;
                   Status: out Error.Condition);
-- Used to undo the updates of this action. Implicitly called when the
-- owning task/job is abnormally terminated, and when the system crashes.
-- Possible errors:
__
     Is_Had_Action (Status):
         o Action_Does_Not_Exist
__
         o Not_The_Action_Owner
         o Action_is_A_Dependent_Action
procedure Prevent_Commit (Id : Action_Id;
                           Status : out Error.Condition);
-- Used to prevent an action from being committed. Useful when an
-- error is detected at the "bottom" of an algorithm, and it wants to
-- guarantee that erroneous changes to open objects are not committed.
-- Ok if the action is already commit-prevented.
-- Possible errors:
     Is_5ad_Action (Status):
--
         o Action_Does_Not_Exist
--
         o Not_The_Action_Gwner
procedure Prevent_Update (Id : Action_ld;
                           Status : out Error.Condition);
-- Used to prevent an action from acquiring an Object_write or Record_write
-- lock on any object. This is implicitly called by Object. Upen when the
-- Supersedable parameter is true. Ok if the action is already
-- update-prevented.
```

```
Is_Bad_Action (Status):
__
         o Action_boes_vot_Exist
         o Not_The_Action_Owner
         o Action_Already_Has_Update_Locks
procedure Give_Up_Ownership (Id : Action_Id:
                             Status: out Error.Condition);
procedure Take_Ownership (Id : Action_Id;
                          Status : out Error.Condition);
-- These operations are used to cause ownership to be transfered from one
-- task to another. The current owner (task I1) executes Give_. Any
-- future operations by T1 will result in Not_The_Action_Owner. The
-- proposed new owner (task T2) then executes Take_ and becomes the new
-- owner.
-- In other words, once vive has been called, the only exported
-- operation that can be successfully applied to the action is Take_.
-- the window between the call to Give_ and the call to Take_, the
-- action is protected (from auto-abandonment) by the continued existance
-- of task T1. The new owner must be in the same job as the previous owner.
-- If a task is currently the owner, and dies, it does an implicit Give_.
-- Possible errors:
___
     Is_dad_Action (Status):
--
         o Action_Does_Mot_Exist
___
         o Not_The_Action_Owner (1st op only)
         o Action_Owner_wat_Willing (2nd op only)
         o Cant_Afer_Ownership_Cross_Jobs (2nd op only)
         o Action_Is_A_Dependent_Action
-- Information requests --
function Is_In_Progress (Id : Action_Id) return boolean;
-- When false, the reason could be one of: (a) The given Id is
-- garbage, and has never been started, (b) the action has committed,
-- and (c) the action has abandoned. There are no exported ways to
-- distinguish between these cases.
function Is_Commit_Prevented (Id : Action_Id) return Boolean;
function Is_Update_Prevented (Id : Action_Id) return Boolean;
-- Returns the corresponding attributes of the action.
                                                        Returns false for
-- garbage input.
function Is_Locked (The_Version : Version_Id) return Roolean;
function Is_Locked (The_version : Version_Id:
                    The_ltex : Second_Key) return Boolean;
-- True iff the object/item is currently locked (in some mode, other
-- than Unsynchronized). Returns false for garbage version_Id.
```

function Owner (Ic : Action\_Id) return Machine.Task\_Id;

page function Owner (Id : Action\_la) return OmiJob\_Id; -- keturns Nil when the given action is not in progress, or the Action\_Id -- is garbage. function Guardian (Ia : Action\_Id) return Machine.Task\_Id; -- Returns Nil when the given action is not in progress, or the Action\_Io -- is garbage. function Master (Id : Action\_Id) return Action\_Id; -- Returns Nil when the given action is a master, is not in progress, or -- the Action\_Id is garbage. function Dependent (Id : Action\_Id) return Action\_Id; -- Returns Nil when the given action is not a master, is not in progress, -- or the Action\_Id is garbage. function Consumed\_Log\_block (Id : Action\_Id) return Integer; function Action\_Nodes (Id : Action\_Id) return Integer; -- Returns 0 when the given action is not in progress, or the Action\_Id -- is garbage. These values reflect how much of the Action\_Manager's -- resources are being consumed by the given action. generic with procedure Visit (Locker\_Id : Action\_Id; Locker\_Job : Om.Job\_Id; In\_Mode : Lock\_Mode); procedure Traverse\_Lockers (The\_Version : Version\_Id; The\_Item : Record\_Key); -- Used to find all the actions which have some lock on the given -- version/item. Calls Visit iff the version/item is locked in some -- mode other than Unsynchronized. May call Visit more than once for the -- same action. Noop for garbage input. generic with procedure Visit (The\_Version : Version\_Id; The\_Item : Record\_Key; In\_inode : Lock\_Mode); procedure Traverse\_Locks (Id : Action\_Id); -- Used to find all the locks currently held by the given action. Does not -- call Visit if the given action is not in progress. In\_Mode will never -- be Unsynchronized. May call Visit more than once for the same -- version/item. Noop for garbage input. procedure Set\_Daemon\_Wait\_Interval (Interval: Duration);

(·

private type Object\_Lock\_Handle is new Action\_Layout.Lock\_List\_Pointer; Nil\_Object\_Lock\_Handle : constant Object\_Lock\_Handle := null;

1 ^

```
GGG
      pppp
              AAA
G
      Ρ
         >
  G
             A
      P
G
G 65
      babb.
             AAAAA
GG
      F
                 A
             Α
6
   G
      P
             Α
                 Α
GGGG
      P
                 Á
```

User: GDA Object: !DE

!DELTA\_KK.REV1\_0\_0.UNITS.OM.NAMING

Version: V(41) Request: 1271

Date: April 24, 1986 wueued: 11:26:18 AM Printed: 11:36:50 AM

```
!DELTA_KK.RE.1_0_0.UNITS.OM.NAMING.V(41)
```

page

1

```
with Job:
with Error:
with Om_Definitions;
package Naming is
    -- pragma Subsystem (Object_management);
    -- Note that there is a parallel package, Naming_Messages, which will
    -- interpret bad status and produce error message strings.
    package Om renames Om_Definitions;
    subtype Context is Om. Naming_Context;
    -- Name resolution --
    -- The rules from TRW's Gamma Max document need to be reproduced here.
    -- Additional feature: Assume that we have a name of the form
    -- "...Foo'V(...). Let "!...Kernel" denote the root of the world
    -- containing the object identified by "...Foo". The following are
    -- examples of uses of the "'V" attribute:
    --
         (a) ...Foo'v(7)
                                        refers to version 7 of Foo
         (b) ...Foo'V(Gamma.Green)
                                        "!.i.Kernel.Gamma" is a directory
                                         with subclass Path:
    --
                                         "!...Kernel.Gamma.Green" is a
                                        universe view; assuming it refers
                                        to a world view which refers to version
    --
                                        14 of Foo, then the given string refers
    --
                                        to version 14 of Foo.
    __
         (c) ...Foo'V(Gamma)
                                        same as "...Foo'v(Gamma.Current)"
    ___
    --
         (d) ...Foo'V(_View'V(3))
                                         "!...Kernel._View'V(3)" is a world
                                         view; assuming it refers to version 21
    --
                                         of Foo, then the given string refers to
    --
                                         version 21 of Foo.
    -- Note that "...Foo'V(Gamma)" is not just a shorthand for "...Foo'V(7)".
    -- In particular, as the destination of an operation (say "Copy", for
    -- example), "...Foo'V(Gamma)" can cause the new version to be created in
    -- the Gamma view.
    -- We do not currently support full pathnames or complex wildcards in the
    -- "'V" attribute. In particular, the only special characters supported
    -- in the "'\" attribute are "\" and "#".
    -- Additional feature: Rational reserves all simple names which both start
    -- and end with an underscore. The empty string is not considered well
    -- formed.
    procedure Resolve (Name : Um.Unique_wildcard;
                       Action : Om.Action_Id;
```

Result : out Om. Version\_Handle;

In\_Context : Context := Job.Default\_Context; Max\_Wait : Duration := Job.Default\_wait); -- To be successfully resolved, the given wildcard must match exactly -- one object. Exception: search lists cause success on first match. -- The object referenced by the given wildcard is identified by -- Result. Version. If the given wildcard specifies a specific version -- number (by using "'V(7)", for example), then Result. Version. Number -- identifies that particular version; otherwise, Result. Version. Number is -- Nil. If the given wildcard specifies a universe/world view (by using -- "'V(Gamma)", for example), then Result. Universe identifies that -- particular universe/world view; otherwise, Result.Universe is Nil. -- Note that the syntax is such that Result. Version. Number and -- Result.Universe cannot both be non-Nil. These rules are such that -- operations like Object.Open "do the right thing". -- Action and Max\_Wait are used iff Name refers to indirect files. -- Note that this form of resolve does NOT produce Diana. Tree. -- Possible errors: Is\_Bad\_Object\_Id (Status) \_\_ o These are caused by a context which does not identify an existing object. These are possible iff Name does not -start with "!". ----Is\_Bad\_Pathname (Status) \_\_ Is\_Bad\_Action (Status) generic with procedure Visit\_Roof (Roof: Om.Object\_Handle); with procedure Visit (Result : Om. Version\_Handle); with procedure Handle\_Error (Status : Error.Condition); procedure Wild\_Resolve (Name : Om.Full\_Wildcard; Action : Om.Action\_Id; In\_Context := Job.Default\_Context; Suppress\_Duplicates : Boolean := False; Max\_Versions : Integer := 2 \*\* 16: Calculate\_Roof : Soolean := False: Max\_wait : Duration := Job.Default\_Wait); -- Similar to the Resolve procedure, with these differences: Is happy -- with 0 or more matches, calling Visit for each match. After an error, -- this procedure will simply stop searching further in the active context, -- going back to any remaining contexts in its stack. The caller is free -- to raise arbitrary exceptions in the generic formal supprograms; these -- terminate the resolution. -- when Suppress\_Duplicates is false, it is possible for sets, indirect -- files, and wildcarded universes to cause the same object to be Visit'd -- more than once. When true, this procedure has to build up an internal -- map to filter out duplicates; this has considerable performance impact.

Status: out Error. Condition;

```
-- The Visit_Roof procedure is called iff Calculate_Roof is true. In this
-- case, Visit_Roof is called prior to any calls to Visit. A "roof" is
-- an object which is a parent of every object identified by visit.
-- routine attempts to find the "best roof", namely the parent which
-- if furthest from the root. Use of this feature causes the procedure
-- to build up an internal map and has considerable performance impact.
-- Max_Versions is used if either Suppress_Duplicates or Visit_Roof is
-- true. It limits the size of the internal map and therefore the number
-- of versions which can result from the resolution. If the error
-- Too_Many_Versions is returned, resolution stops, and Visit is called
-- for all of the versions currently in hand.
-- Possible errors:
--
     Is_Bad_Version_Handle (Status)
         o Something wrong with In_Context.
--
     Is_Bao_Pathname (Status)
--
         o Something wrong with Name.
--
     Too_Many_Versions
--
         O 'ax_Versions is smaller than the number of objects which
--
           match the wildcard. Cannot be returned unless
___
           Suppress_Duplicates or Calculate_Roof is true.
--
    Is_Bad_Action (Status)
type Create_kind is new Long_Integer range 0..2 ** 3 - 1;
Creating_New_Version : constant Create_Kind := 1;
Creating_New_Object : constant Create_Kind := 2;
Creating_Temb_Object : constant Create_Kind := 3;
procedure Resolve_For_Create (Name : Om.Unique_Wildcard;
                              Action: Om.Action_Id;
                              Kind: out Create_Kind;
                              Result : out Om.Object_Handle;
                              Child_Name : out Om.Simple_Name;
                              Class_Of_New_Object : out Om.Class_Number;
                              Status : out Error.Condition;
                              In_Context := Job.Default_Context;
                              hax_Wait : Duration := Job.Default_Wait);
-- This operation is used by the string form of Object. Create. Assuming
-- Status is is Successful: If the given Name resolves to an already
-- existing object, then Kind will be Creating_New_Version, and Result
-- will identify the object (for which a new version is to be created),
-- just as if the vanilla Resolve had been invoked. If the string is empty
-- (with the possible exception of a "'C" attribute) then Kind will be
-- Creating_Temp_Object, Child_Name will be a simple name string generated
-- from a unique id (returned by Om.Get), and Class_Of_New_Object will be
-- Nil (or the value of the optional ""C" attribute). Otherwise, Kind will
-- be Creating_Vew_Object, Result will identify the parent of the new
-- object, Child_Name will contain the simple name of the new object, and
-- Class_Of_New_Object will be Nil (or the value of the obtional "'("
-- attribute).
-- Possible errors:
     includes all those that can be returned by vanilla Resolve
```

```
procedure Substitute (The_Object : Om.Version_Handle;
                      Source : Om.Full_Wildcard;
                      Destination: Om.Full_Wildcard;
                      Result : out Om.Pathname;
                      Status: out Error.Condition);
-- Given a source name (with wild cards) that matches the name of the
-- given object, return a Result string in which substitution characters
-- have been replaced by the matching portions of the object's name as
-- indicated by the destination pattern.
-- Substitution semantics need to be fully specified.
-- Used by Object.Wild_Copy.
-- RESTRICTION: not supported in 1st qtr impl.
-- Possible errors:
     Is_Ead_Version_Handle (Status)
--
         o Something wrong with The_Object.
___
     Is_Bad_Pathname (Status)
__
         o Something wrong with Source or Destination.
           Is_well_Formed function to see which one is bad.
-- Detailed structure of names --
function Is_Well_Formed (Name : Om.Full_Wildcard) return boolean;
procedure Check_Well_Formed (Name : Om.Full_Wildcard;
                             Status: out Error.Condition);
-- Possible errors:
     Filename_Syntax_Error
         o Status identifies the character in Name, which caused the error,
             as well as the token class which was expected.
function Prefix (Name: Om.Fult_Wildcard) return Om.Fult_Wildcard;
function Head (Name: Jm.Full_Wildcard) return Om.Full_Wildcard;
function Tail (Name: Om.Full_Wildcard) return Om.Full_Wildcard;
-- Given ill formed names, these functions produce undefined results.
-- Prefix removes the last segment. Head returns only the first segment.
-- And Tail removes the first segment.
      Prefix ("A.B.C") => "A.B"
_ _
      Prefix ("A") => ""
--
__
      Head ("A.B.C") => "A"
      Head ("A") => "A"
      Tail ("A.3.C") => "3.C"
--
      Tait ("A") => ""
function Attributes (Name: Om.Full_wildcard) return String;
function Class_Attribute (Name : Om.Full_Wildcard) return String;
function Version_Attribute (Name : Om.Full_kildcard) return String;
function Part_Attribute (Name : Om.Full_wildcard) return String;
-- Given ill formed names, these functions produce undefined results.
```

-- Attributes returns the attributes at the end of the Name, including the

```
!DELTA_KK.REV1_0_0.U4115.04.WAMING*V(41)
                                                                       page 5
    -- leading "'", or a null string if Name has no a attributes.
    -- remaining functions return the indicated attribute.
    procedure Get_Full_Name (The_Object : Om.Object_Handle;
                             Full_Name : out Om.Pathname;
                             Status : out Error.Condition);
    procedure Get_Simple_Name (The_Object : Om.Object_Handle:
                               Name : out Om.Simple_Name:
                               Status: out Error.Condition);
    -- Returns the full/simple name for The_Object, excluding qualifying
    -- attributes (like the "'V" attribute). The Version_Number component of
    -- The_Object is ignored.
    -- Possible errors:
          Is_Bad_Version_Handle (Status)
    procedure Get_Unique_Full_Name (The_Object : Om.Version_Handle;
                                    Full_Name : out Om.Pathname;
                                    Status : out Error.Condition);
    procedure Get_Unique_Simple_Name (The_Object : Om.Version_Hangle;
                                      Name : out Om.Simple_Name;
                                      Status : out Error.Condition);
    -- Returns the fully/simple name for The_Object, including qualifying
    -- attributes (like the "'V" attribute):
    -- Possible errors:
          Is_Bad_Version_Handle (Status)
    -- Detailed directory structure --
    function hoot return Om. Object_Mandle;
    -- Returns the Object_Id for "!".
    procedure Get_Parent (Child: Om.Object_Handle;
                          Parent : out Om.Object_Handle:
                          Status: out Error.Condition);
    -- Returns the handle of the Child's parent. If Child refers to "!", then
    -- returns Nil_Object_Handle.
    -- Possible errors:
         Is_bad_Object_Id (Status)
    procedure Get_Containing_Directory
                   (Of_Object : Om.Object_Handle;
                    Containing_Directory : out Om.Object_Handle;
                    Status : out Error.Condition);
    procedure Get_Containing_world
                   (Of_Object : Om.Object_Handle;
                    Containing_world : out Om.Object_Handle;
```

Status: out Error.Condition);

worlds : Boolean := True;

```
procedure Get_Associated_Directory
               (For_Object : Om.Object_Handle:
                Associated_Directory: out Om.Object_Handle;
                Status: out Error.Condition);
procedure Get_Associated_World
               (For_Object : Om.Object_Handle;
                Associated_World: out Om.Object_Handle;
                Status: out Error.Condition);
-- The Containing_ functions return the nearest enclosing world/directory.
-- The Associated_ functions are the same as the Containing_ functions,
-- except if the argument is a world/directory, the argument is returned.
-- Possible errors:
     Is_Bad_Version_Handle (Status)
procedure Get_Child (Parent : Om.Object_Handle;
                     Child_Name : Om.Simple_Wildcard:
                     Child_Id : out Om.Object_Handle;
                     Status: out Error.Condition);
-- Returns the Object_Handle of the one child (of parent) whose simple
-- name matches Child_Name (which may include wildcards).
-- Possible errors:
     Is_Bad_Version_handle (Status).
        o Something wrong with Parent
     Is_bad_Pathname (Status)
--
        o Something wrong with Child_Name
generic
    with procedure Visit (Child_Id : Om.Object_Handle;
                          Child_Name : Om.Simple_Name);
procedure Traverse_Children (Parent : Om.Object_Handle;
                             Status: out Error.Condition);
-- Calls Visit for every child object of Parent. Visit is called in
-- sort order, by Child_Name.
-- Possible errors:
     Is_Bad_Version_Handle (Status)
        o Something wrong with Parent
generic
    with procedure Visit (Depth: Positive;
                          The_Object : Om. Version_Handle;
                           Its_Name : Om.Simple_Name_String;
                           Go_Deeper : out Boolean);
    with procedure Handle_Error (Status : Error.Condition);
procedure wepth_first_Traversal
     (Root_Of_Traversal : Om_Object_Handle;
      Pattern : UmiSimple_Wildcard := "0";
      Class: Om.Class_Number := Om.Nil;
      Subunits : Boolean := True;
      Directories : Boolean := True;
```

Versions : Boolean := False);

- -- Calls Visit for Noot\_uf\_Traversal, with Depth = 1. Its children will -- have a depth of 2, etc. Child objects are visited in sort order by -- Its\_Name. Versions are visited in sort order by Version\_Number. Assume -- that objects "C1" and "C2" are children of object "P". "depth first" -- implies that all of C1's children (including all their versions) are -- Visit'd prior to visiting C2 or its children. It also implies that all -- versions of C1 are visited prior to any version of C2.
- -- The Go\_Deeper out parameter of Visit can be used to prune the traversal.
- -- The Pattern, Class, Subunits, Directories, Worlds and Versions -- parameters are filters which prune the traversal.
- -- The caller is free to raise arbitrary exceptions in the generic formal -- subprograms; these terminate the resolution.
- -- Possible errors:
- -- Is\_bad\_Version\_handle (Status)
- o Something wrong with Root\_Of\_Traversal
- -- Is\_Eau\_Pathname (Status)
- o Something wrong with Pattern

end Naming:

```
PPPP
 GGG
               AAA
G
       p 5
  G
              Α
                  Α
6 GG
G G
       Р
           P
              A
                  Α
       PPPP
              AAAAA
      Р
              A
                  Α
G
   Ğ
       Ρ
              Α
                  A
 GGGG
       P
              A
                  Α
```

User: GPA

Object: !DELTA\_KK.REV1\_0\_0.UNITS.OM.OBJECT

Version: V(173) Request: 1272

Date: April 24, 1986 Queued: 11:20:30 AM Printed: 11:37:20 AM

```
!DELTA_KK.REV1_9_0.UNITS.O4.UBJECT V(173)
                                                                      page 1
with Job;
with Error;
with Action;
with Machine;
with Calendar;
with Om_Definitions;
package Object is
    -- pragma Subsystem (Object_Management);
    package Om renames Om_Definitions;
    subtype Unique_Id is Om_Definitions.Unique_Id;
    subtype Context is Om. Naming_Context;
    -- Basic Operations --
    -------
    -- Normally, objects have 1 or more versions; that is, expunging the last
   -- version of an object will normally cause the object itself to be
    -- expunged. However, the user model includes the ability for objects to
    -- exist with no versions. In particular, when the last version of an
   -- object is expunyed, but the object still has children, then the parent
    -- object will continue to exist, with no versions. When the child objects
    -- are expunsed, the parent object (with no versions) will also be
    -- expunged).
    -- A view can select a version of a child object without selecting a
    -- version of the parent object.
    -- Note that the above semantics apply to directories as well as other
    -- objects (since all object classes may have children).
    -- The phrase "usual rules for Naming.Resolve" means that the operation
    -- takes 3 parameters:
                Name : Um.Unique_wildcard:
    --
                Action: Om.Action_Id:
                In_Context := job.default_context;
    -- These parameters are passed to Naming.Resolve to compute the
    -- Version_Handle. Typically a second form of the operation is also
    -- provided which takes a Version_Handle, to avoid calling Naming.Pesolve.
    -- The term "retention list" refers to the set of versions selected by
    -- a world view. The "current" version (selected by a world view) is
    -- the "top" item of the retention list. "Pushing" causes a new version
    -- to be added to the top of the retention list, becoming current;
    -- note however that if current is Nil, pushing simply causes the new
    -- version to replace current. If the retention list was already full (as
    -- specified by the retention count in force for the world view), then
    -- pushing may also cause the "bottom" item to be removed from the list,
    -- the removed version is expunded if possible (as defined by the Expunde
    -- procedure later in this spec). The retention list (together with the
    -- Reserved bit) is also referred to as the "slot" for the object.
```

<sup>--</sup> The phrase "usual rules for interpretation of Version\_Handle" means:
-- (a) As the source of an operation (or the target of an operation

```
2
         which does not create a new version): If
         Version_Handle.Version.Number is not Nil, then the specified
__
_ _
         version is used, and Version_Handle.Universe is ignored.
         Otherwise, we use the Version_Number which is current in the
         selected world view. The "selected view" is specified by the
__
         .Universe component or Job.Default_world_View when the .Universe
--
         component is Nil. A Record_Read lock is acquired on the slot.
___
     (b) As the target of an operation which creates a new version,
         Version_Handle.Version.Number is ignored, and the new version is
--
_--
         created in the selected universe/world view; the definition of
__
         "selected" is identical to that above. A Record_write lock is
         acquired on the slot. (If the Record_Write lock fails because
         the action already has Object_Write, updates to world views
         will use the Object_Write lock instead.)
-- The phrase "usual rules for creation of new version" means:
-- "Explicit version creation" (calling Create) causes the new
-- version to be pushed onto the retention list (for the object in the
-- selected world view), thereby becoming its "current version". This may
-- cause the bottom predecessor version to become no longer selected by
-- the world view (as required by the retention count); this predecessor
-- version will be expunded at action commit time. A Record_Write lock is
-- acquired on the slot.
-- "Implicit version creation" (calling Open with an update mode and
-- appropriate Version_Control_Policy) is identical to the above case when
-- the Version_Control_Policy does NOT have the
-- Implicit_Versions_Replace_Current bit set. Otherwise, the new version
-- simply replaces the current version in the retention list; the replaced
-- version is expunded at commit time; and a Record_Write lock is acquired.
-- Included in the usual rules for new versions are the following Ada
-- specific properties: If the object has Class_Ada and a parent which is
-- not Class_Ada, and the "current version" (in the selected world view)
-- was previously Nil, then an appropriate Link_Pack entry is made. If the
-- selected view's Link_Pack is currently shared with other views, then a
-- new version is made (this copy is committed immediately, and will
-- therefore not be undone by action abandon). This latter rule is
-- referred to as the "usual rule for undifferentiating the Link_Pack".
-- Unless specified otherwise, all locks are acquired on behalf of the
-- given Action, and the operation will suspend for up to Max_Wait for
-- each lock. The order in which operations acquire locks is not specified.
-- Unless specified otherwise, one can assume that an operation which
-- fails leaves "the world" the way it was prior to the operation.
```

procedure hesolve (The\_Object : Om. Version\_Handle; The Version : out Om. Version\_Id; Status: out Error.Condition);

-- made, the action will become Commit\_Prevented.

-- Takes the given Version\_Handle and computes the actual Version\_Id

-- An exception: when an operation fails, it does NOT release locks that -- it acquired prior to the failure; if no modifications have been made, -- the action simply holds extraneous locks; if modifications have been

```
-- according to the usual rules for interpretation of source
-- Version_Handle's.
-- Possible_trrors:
    Is_bad_Version_Handle (Status);
type Open_Handle is private:
Nit_Open_Handle : constant Open_Handle;
procedure Create (Name : Om.Unique_Wildcard;
                  Action : Om.Action_Id;
                  Handle: out Open_Handle;
                  Status : out Error.Condition;
                  Class: Om.Class_Number := Om.Nil;
                  In_Context : Context := Job.Default_Context;
                  Initialize : Boolean := False;
                  Success: Om. Subclass_Number := Om. Nil;
                  Version_Policy : Om.Version_Control_Policy := Om.Nil;
                  Index : Om.Object_Index := Om.Wil;
                  Creation_Id : Om.Net_Unique_Id := Om.Nit_Net_Unique_Id;
                  Max_wait : Duration := Job.Default_Wait);
procedure Create (The_Unject : Om.Object_Handle;
                  Action : Om.Action_Id;
                  Handle: out Open_Handle;
                  Status : out Error.Congition;
                  Initialize : Boolean != False:
                  Subclass : Om.Subclass_Number := Om.Nil;
                  Max_Wait : Duration := Job.Default_Wait);
-- In the first form, Naming. Resolve is used to compute the Object_Handle.
-- If this succeeds, then a new version of the object is created.
-- Otherwise, if the Name is null (except for an optional "'C" attribute),
-- then a "temporary object" is created. In this case, In_Context is
-- ignored. The object is created in the Job.Sarbage_Directory with a
-- unique name constructed from the time and date. The given action
-- becomes commit prevented.
-- Otherwise, Prefix(Name) and In_Context are used to compute the
-- Object_Id of the parent (of the object to be created); the last segment
-- is used as the simple name of the child object (to be created).
-- For a newly created object, the "*C" attribute may be used to specify
-- the class of the object, as well as the Class parameter.
-- For a newly created object, the Version_Policy parameter specifies
-- how future operations on the object will behave. All versions of the
-- object will follow the same policy. If a Mil Version_Policy is supplied
-- the Default_Version_Policy is used.
-- For a newly created object, the Reserved bit is set.
-- For a newly created object, it is possible to specify the desired
-- Object_Index and Creation_Id. This facility is used to restore an
-- an object from "archive" or move it from machine to machine, in the
-- context of a particular view or world.
```

- -- In the second form, The\_Object refers to an existing object, for -- which a new version is created.
- -- The following rules apply to both forms:
- -- Usual rules for interpretation of the Object\_Handle. Usual rules for -- creation of a new version.
- -- The Version function can be used (with the the returned Open\_Handle) to -- determine the id of the newly created version. The Version\_Number field
- -- (of the id) will be non-Nil. The new version is held with an
- -- Object\_Write lock, and is open.
- -- If Initialize is false, the new version is empty, and we say that "an
- -- empty version was created". If the operation creates the first version
- -- of a new object, the Initialize parameter is treated as if it were
- -- false. Otherwise, an Object\_Read lock is acquired on the previous
- -- version (selected by the selected world view) and its data bits are
- -- copied into the newly created version, and we say that "a non-empty
- -- version was created".
- -- The Predecessor\_Update\_Id function can be used (on the returned
- -- Open\_Handle) to determine the stamp of the copied version.
- -- If a non-Nil value is supplied for Subclass, then the subclass of the
- -- new version is set as specified (note that, unlike class, different
- -- versions may have different subclasses). If a Nil value is supplied,
- -- and Initialize is treated as True, then the subclass of the new version
- -- is copied from the predecessor version. In the remaining cases, the
- -- subclass of the new version is Nil.
- -- If Initialize is treated as True, the Predecessor\_Update\_Id attribute
- -- is set to the Last\_Update\_Id of the source (of the new version's
- -- data bits), else Predecessor\_Update\_Id is Nil.
- -- Class specific semantics (these are in addition to or countermand the
- -- general rules given above); there are no subclass specific semantics:
- -- Class\_Directory: The directory will NOT be the root of a new world;
- -- must use the World package to do that. The returned Open\_Handle will
- -- be Nil. The Initialize parameter is ignored.
- -- Class\_File: An empty file has organization Binary\_File, and contains
- -- no data bits. A non-empty file has the same organization, and data
- -- bits. Only the ID packages should use this operation: "users" should
- -- use the Create in the IU packages.
- -- Class\_Pipe: An empty pipe has some default buffer size, whereas a
- -- non-empty pipe has the same buffer size as the predecessor. The pipe
- -- is locked as specified by the Mode, and is open from the point of
- -- view of this package. Note that a mode of Record\_Write is typically
- -- more useful than the default of Object\_Write. Only the IO packages
- -- should use this operation; "users" should use the Create in the ID
- -- packages.
- -- Class\_Ada: Additional semantics supplied by the Ada coupler and are as
- -- yet unspecified.

\_\_\_

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```
-- Class_Ada_Attributes: Additional semantics supplied by the Ada coupler
-- and are as yet unspecified.
-- Class_Code: _ike other classes of objects (such as files), a code
-- version is modified using the segment operations at the end of this
-- package. Just as for other classes of objects, the contents of the
-- version can only be addressed (as a heap) while Open. In addition to
-- this "heap name", each version of a code object also has a distinct
-- code segment name assigned to it; the contents of the version can
-- always be addressed via the code segment name, even when the version is
-- not Open; but the version cannot be modified via the code segment name.
-- A page fault on page "I" of the code segment is resolved by creating
-- the code page I as a copy of the current value of page I+1 of the head
-- page (of the corresponding version).
-- Class_Load_Image: Initialize is always treated as false.
-- Possible errors:
--
     Is_Bad_Version_Handle (Status):
___
     Is_Bad_Pathname (Status):
--
         o Might be because an object must be created, but the last segment
--
           (name of new child) contains wildcards.
__
     Is_Bad_Action (Status)
--
     Is_Lock_Error (Status):
--
     Predecessor_Version_Is_Not_Known
___
         o Attempting to create a new version of an existing object, and
--
           the Initialize parameter is true, but the selected world view
___
           does not currently select a version of the object, therefore,
___
           we don't know the version number of the "predecessor".
__
     Link_Pack_Entry_Already_Exists
--
         O Attempting to create a new Ada library unit with a simple
___
           name which conflicts with the name of some unit which is
--
           already in the Link_Pack of this world.
--
     Illegal_Class
___
         o This operation is not supported for Class_World_View,
--
           Class_Universe_View, Class_Link_Pack, Class_Tabe, or
--
           Class_Terminal.
--
     Object_Class_Mismatch
--
         o Can occur in 1st form only. If the object already existed,
--
           means that the existing object does not have the class specified
--
           by the non-Nil Class parameter. If the object did not already
___
           exist, and the class param and "'C" attribute are both specified
--
           and neither is Nil, then this means that they are not equal.
__
     Object_Index_Already_In_Use
___
         o The Index parameter is non-Nil and refers to an already
--
           existing object.
__
     world_View_Is_Frozen
--
         o The specified world view is frozen; therefore it is not
--
           possible to create a version of the object.
--
     Is_desource_Limit_Error (Status):
___
     Is_Version_Policy_Error (Status):
--
         o Undefined_Version_Policy
--
         o Multiple_Versions_Not_Supported
               Means that either the class of the existing object does not
--
               support multiple versions, or the
--
               Allow_Explicit_Create_Version bit is not set in the given
```

policy options. Note that directories can have at most 1

```
version.
--
         o Object_Is_Not_Reserved
               Attempting to create a new version of the object, but
               the object's slot (in the target view) does not have the
               Reserved bit set.
procedure Open (Name : Om.Unique_Wildcard;
                Action : Om.Action_Id;
                Mode : Om.Lock_Mode:
                Handle: out Open_Handle;
                Status: out Error.Condition;
                Class : Om.Class_Number := Om.Nil;
                In_Context : Context := Job.Default_Context;
                Ignore_Reserved_Bit : Boolean := False;
                Max_Wait : Duration := Job.Default_Wait);
procedure Jpen (The_Version: Om.Version_Handle;
                Action : Om.Action_Id;
                Mode : Om.Lock_Mode;
                Handle: out Open_Handle:
                Status: out Error.Condition:
                Ignore_Reserved_Bit : Boolean := False;
                Max_Wait : Duration := Job.Default_Wait);
-- Usual rules for Naming.Resolve, and interpretation of the Version_Handle.
-- The following special rules take precedence over the usual rules:
       (a) The_Version.Version.Number specifies a Nil version and
-- If
       (b) Mode is one of the update modes (eg, Write_Object) and
___
       (c) The selected world view is not frozen and
___
       (d) Either
              (1) The Private_Version_On_Every_Update policy bit
<u> -</u> -
__
                  is set, or
              (2) The Private_Version_On_First_Update policy bit
                  is set, and the version that would be selected as a
                  target under the normal rules is either frozen or
--
                  referenced more than once
--
-- then an Open_Handle is generated by an implicit call to Create, with
-- Initialize = True; this is known as "differentiation". If the
-- Implicit_versions_Replace_Current policy bit is set, the management of
-- the retention list is slightly different (than the rules specified for
-- the Create procedure). Specifically, the new version simply replaces
-- the current version.
       (a) The_Version.Version.Number is not Nil and
-- If
       (b) Mode is one of the update modes (eq. write_Object) and
       (c) The Prevent_Update_Without_View policy bit is not set and
       (d) The specified version is not frozen
-- then the identified version is open'd for update. Note that this rule
-- will allow the following: One can update an unfrozen version which is
-- shared by more than one view. One can update a deleted (but not yet
-- expunged) version.
-- If The_version. Version. Number is Nil and the Mode is not in the set
-- {Unsynchronized, Supercedeable_Read}, then the action acquires a
```

-- Record\_Read lock on the slot; in all other cases, the given action

```
-- acquires no slot lock. In any case, the action acquires the requested
-- lock on the selected version. Note that if differentiation results, a
-- seperate action is created which acquires a slot Record_Write lock, and
-- Object_Read lock on the copied (predecessor) version, and is committed.
-- Due to internal implementation requirements, when a new version is
-- implicitly created (as defined by the above paragraph), the
-- undifferentiation is actually done on a separate action which is
-- committed immediately: thus, the undifferentiation will not be undone
-- by action abandon. This should be "transparent".
-- When a new version is implicitly created, the new version inherits
-- the subclass attribute from the source (of the new version's data bits).
-- When a new version is implicitly created, the Predecessor_Update_Id
-- attribute is set to the Last_Update_Id of the source (of the new
-- version's data bits), else Predecessor_Update_Id is Nil.
-- When an object is first Open'd (for Supersedable read), and later
-- Open'd for update, the Prior_Update_Id function should be used to
-- detect a situation in which the object has been modified between the
-- initial Open (for read) and the subsequent Open for update.
-- To find out what version was Open'd, feed the Open_Handle to the Version
-- function.
-- Class specific semantics:
-- Class_Code: Observe that if Node is one of the update modes, but
-- differentiation was not performed, this operation gives one the
-- ability to change the contents of an existing code segment.
-- Remember that code segment names are assigned to new versions, not
-- new generations. See the Refresh_Code_Segment operation.
-- Class_Load_Image: Update modes are not allowed in conjunction with
-- objects of this class. One must use the Create operation and get a
-- new version. This restriction might be relaxed at some future date.
-- Possible errors:
--
     Is_Bad_Version_Handle (Status)
--
     Is_Bad_Pathname (Status)
--
     Is_Bad_Action (Status)
--
     Is_Lock_Error (Status):
--
     Illegal_Class
--
         o This operation is not supported for Class_Directory,
--
           Class_Universe_View, Class_World_View, and Class_Link_Pack.
--
     Version_Is_Frozen
__
     Is_Resource_Limit_Error (Status):
--
     Is_Version_Policy_Error (Status):
--
         o Object_Is_Not_Reserved
__
               Attempted to Open with one of the update modes,
___
               The_version.Version.Aumber is Nil, the
_ _
               Ignore_deserved_bit parameter is False, and the
--
               the object's slot (in the target view) does not have the
--
               Reserved bit set.
__
         o ust_Update_Via_View
___
               The_Version.Version.Number is not Nil, but the
___
               Prevent_Update_Without_View policy bit is set.
```

```
procedure Delete (Name : Om.Unique_Wildcard;
                  Action : Om.Action_Id;
                  Status : out Error Condition;
                  In_Context : Context := Job.Default_Context;
                  Max_Wait : Duration := Job.Default_Wait);
procedure Delete (The_Object : Om.Object_Handle;
                  Action : Om.Action_Id;
                  Status : out Error.Condition;
                  max_Wait : Duration := Job.Default_Wait);
-- The usual rules for Naming. Resolve. Usual rules for selecting a
-- world view (from The_Object).
-- The following additional rules apply: Object.Version.Number is ignored.
-- Nil is pushed onto the retention list (of the selected view),
-- following the usual rules for creation of a new version (the new
-- version simply happens to be Nil).
-- Note that the "deleted" version(s) can still be named (using string
-- names) via other referencing world views, or by use of the "'V"
-- attribute; and they can, of course, be named by their Version_Id*s.
-- If the deleted version has Class_Ada and a parent which is not
-- Class_Ada, then the corresponding Link_Pack entry is deleted, with the
-- usual rules for undifferentiating the Link_Pack. There are additional
-- (unspecified) obsolecense rules for Class_Ada, Class_Ada_Attributes, and
-- Class_Code.
-- The action will acquire a Record_Write lock on the corresponding slot.
-- Possible errors:
--
     Is_sad_Version_Handle (Status)
         o Object_Not_In_View
               o The object is already deleted in this view.
__
--
     Is_Bad_Pathname (Status)
--
     Is_Bad_Action (Status)
--`
     Is_Lock_Error (Status):
--
     Illegal_Class
__
         o This operation is not supported for Class_Universe_View,
--
           Class_World_View, and Class_Link_Pack.
     World_View_Is_Frozen
--
__
     Is_Resource_Limit_Error (Status)
     Is_Version_Policy_Error (Status):
--
--
         o Object_Is_Not_Reserved
__
               The object's slot (in the target view) does not have the
               Reserved bit set.
generic
    with procedure Visit (Version_To_Be_Deleted : Om.Version_Handle);
    with procedure Handle_Resolve_Error (Status : Error.Condition);
    with procedure Handle_Delete_Error (Status : Error.Condition;
                                         bad_Version : Om.Object_Handle);
```

procedure wild\_Delete (Name : Om.Full\_Wildcard;

```
In_Context : Context := Job.Default_Context;
                       Yax_Wait : Duration := Job.Default_Wait);
-- Basically, employs Naming. Wild_Resolve (using Name, Action, and
-- In_Context) to produce a sequence of Version_Handle's; errors produced
-- by Wild_Resolve are fed to Handle_Resolve_Error. For each
-- Version_Handle produced by Wild_Resolve, calls Visit (with a fully
-- resolved Version_Handle), and then the corresponding form of Delete;
-- its errors are fed to Handle_Delete_Error.
-- As usual, the caller is free to raise arbitrary exceptions in the
-- generic formal subprograms; these terminate the deletion process,
-- however, it is up to the client to abandon the action, if that is what
-- is desired.
-- Note that one can use "...Foo'V(@)" to cause the object to be deleted
-- from all views.
-- Note that one can do things like: set Max_Wait to 0, and ignore
-- calls to Handle_Delete_Error which produce Status for which
-- Is_Lock_Error is true.
-- Clearly, this operation can do lots of damage, especially when followed
-- by World.Expunge.
-- RESTRICTION: not supported in 1st qtr impl.
-- Possible errors:
        o Union of those produced by Naming. Wild_Resolve and the
          Version_Handle form of Delete.
procedure Undelete (Name : Om.Unique_Wildcard;
                    Action : Om.Action_Id;
                    Status: out Error Condition;
                    In_Context := Job.Default_Context;
                    Max_wait : Duration := Job_Default_Wait);
procedure Undelete (The_Version : Om.Version_4andle;
                    Action: Om.Action Id;
                    Status : out Error.Condition;
                    Max_Wait : Duration := Job.Default_Wait);
-- Usual rules for Naming.Resolve.
-- The version to be undeleted (and the view from which it is undeleted) is
-- chosen by the normal rules for the interpretation of a source
-- Object_Handle.
-- The following additional rules apply:
-- If The_version.version.Number = Nil then
__
       If the current version is Nil then
           pop the Mil off the top of the retention list;
           causing the acquisition of a Record_Write lock on the slot;
___
           the operation fails:
       end it:
```

Action : Om.Action\_Id;

```
-- else
___
       If it refers to a version in the retention list then
           the identified version is removed from the retention list, and
           pushed on at the top of the list, becoming current, and
           causing the acquisition of a Record_Write lock on the slot;
       else
           the operation fails:
       end if:
__
-- end if;
-- When Nil is popped off, a Link_Pack entry may be created, according to
-- the usual rules for creation of new versions.
-- There are additional (unspecified) obsolecense rules for Class_Ada,
-- Class_Ada_Attributes, and Class_Code.
-- The action will acquire a Record_Write lock on the corresponding slot.
-- Possible errors:
     Is_bau_Version_Handle (Status)
         o Object_Not_In_View
--
               o Either the specified version is not in the retention list,
--
                 or no version was specified and the retention list
                 contains all Nils.
__
     Is_Bad_Pathname (Status)
--
     Is_Bad_Action (Status)
     Is_Lock_Error (Status)
--
__
     Itlegal_Class
--
         o This operation is not supported for Class_Universe_View,
___
           Class_World_View, and Class_Link_Pack.
--
     World_View_Is_Frozen
     Is_Resource_Limit_Error (Status)
--
--
     Is_Version_Policy_Error (Status):
__
         o Object_Is_Not_Reserved
_-
               The object's slot (in the target view) does not have the
               Reserved bit set.
--
procedure Expunse (Name : Om.Full_Wildcard;
                   Number_Expunged : out Natural:
                   Disk_Blocks_Freed : out Natural;
                   Status : out Error.Condition;
                    In_Context : Context := Job.Default_Context;
                   Violate_Safeguards : Boolean := False);
procedure Expunge (The_Version : Om.Version_Handle;
                   Number_Expunged : out Natural;
                   Disk_Blocks_Freed: out Natural;
                   Status : out Error.Condition;
                    Violate_Safeguards : Boolean := False);
-- The meaning of the word "expunge" is defined by the semantics of this
-- operation. Note that these semantics apply to all classes of object,
-- including interesting cases like universe views, world views, and
-- directories.
-- Note that this is the only operation for destroying universe views.
```

-- and therefore the only operation for removal of unwanted releases, etc.

!DELTA\_KK.REV1\_0\_0.UNITS.OM.OBJECT\*V(173) page -- Note that one can protect frozen views from expunge via access control. -- Usual rules for Maming.Resolve and interpretation of Version\_Handle. -- Plus the following: -- If The\_Version.Version.Number is Nil then \_\_ remove all deleted versions in the retention list \_\_ (by "deleted", we mean one which is not current), -causing the acquisition of a Pecord\_Write lock on the slot; -- else -remove the specific version from the retention list, if its there, causing the acquisition of a Record\_Write lock on the slot; -- end if; -- Note that the above rule is simply a noop in the case where the -- selected universe view does not reference a world view, or the slot -- cannot be locked. -- If any versions removed (from the retention list) in the above step are -- expungeable, then reclaim their disk space. -- If (at this point) the object itself is expungeable, then it ceases to -- exist. Note that if the object itself is not expungeable, this -- operation may leave the object in existence, without any versions. -- If the object itself is expunged, and the object's parent is -- expungeable, then the parent is also expunded. This rule is applied -- recursively. -- A version is physically expungeable if it is not open (in a mode -- other than Unsynchronized). -- A version's expunse safeguards are satisfied if the version is not --(1) Referenced by a world view; and not \_ \_ (2) A world view referenced by a universe view in the -same world; and not \_\_ (3) A world view protected by a still existing universe view \_\_\_ in a potentially different world; and not --(4) A secondary load image referenced by a still existing -primary; and not (5) An Ada unit referenced by a load image; and not (6) A code segment referenced by a load image. \_\_ -- A version is expungeable if (1) It is physically expungeable; and \_\_ (2) Either the version's expunge safeguards are satisfied, or the Violate\_Safeguards parameter is true. -- An object is expungeable if (1) There are 0 versions of the object; and (2) The object has no children. -- In the event that an object/version is expunsed even though the -- safeguard conditions are not met, the operation does NOT attempt to -- first remove the references, rather it leaves the references -- "dangling"; this behaviour is desirable in cases where the references

-- come from damaged or mysteriously vanished objects. Generally speaking,

-- this option is used only as a "last resort". It might be used to

```
-- recover from objects which have been damaged by software bugs or disk
-- errors.
-- Note that this operation does not follow action semantics.
-- On return, Number_Expunged indicates how many versions this operation
-- actually expunded, and Disk_Blocks_Freed indicates how many 1K byte
-- blocks were freed.
-- Class specific semantics:
-- Class_Directory (subclass World): The single version of the world
-- is not expunded by this operation. Use World.Destroy instead.
-- Class_Load_Image: If any of the referenced Ada units or code segments
-- are now no longer referenced, then the referenced version (and possibly
-- object) is expunded. For a primary load image, expundes all of the
-- referenced secondary load image versions (and probably objects).
-- Possible errors:
___
     Is_Sad_Version_Handle (Status)
     Is_Resource_Limit_Error (Status)
     world_View_Is_Frozen
___
     Illegal_Class
--
--
         o This operation is not supported for Class_Universe_View,
__
           Class_World_View, and Class_Link_Pack.
__
     Is_Version_Policy_Error (Status):
___
         o Object_Is_Not_Reserved
               The object's slot (in the target view) does not have the Reserved bit set. This will only result when one attempts to
___
--
                effect the "current" version.
procedure Copy (Source : Om.Unique_Wildcard;
                 Destination: Om. Unique_Wildcard;
                 Action : Om.Action_Id;
                 handle: out Open_Handle;
                 Status: out Error.Condition;
                 Source_Context : Context := Job.Default_Context;
                 Destination_Context : Context := Job.Default_Context;
                 Max_Wait : Duration := Job.Default_Wait);
procedure Copy (Source : Om.Object_Handle;
                 Destination: Om.Object_Handle:
                 Action : Om.Action_Id;
                 Handle: out Open_Handle:
                 Status: out Error.Condition;
                 Max_wait : Duration := Job.Default_Wait);
-- Usual rules for Naming. Resolve and interpretation of Object_Handle's.
-- Open the source basically as follows:
      Object.Open (Source, Action, Mode => Read_Object);
-- This acquires the usual slot and version lacks, as specified by Open.
        (a) Source and Target refer to the same object; and
```

(b) Either the Private\_Version\_On\_First\_Update or

Private\_Version\_Un\_Every\_Update policy bits are set

```
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-- then the target view is modified to reference the source version, as
-- follows. If the object is deleted in the target view, or Open for
-- update would create and push a new version, then the version is pushed
-- onto the target retention list; otherwise, the version simply replaces
-- current in the target retention list.
-- Note that above capability can be used to construct arbitrary world views
-- and can therefore be used to merge paths, etc, by building the
-- appropriate differential view.
-- In the remaining cases, we actually copy data. If the object is deleted
-- in the target view, then we open the target as follows:
     Create (Destination, Action, Initialize => False);
-- else:
      Open (Destination, Action, Mode => Write_Object);
-- The data bits of the source are then copied into the target.
-- virtual length of the target becoming the same as the virtual length of
-- the source.
-- The target version is left open, by the returned Open_Handle.
-- korks for Class_Load_Image; in the case of a physical copy, does the
-- "right thing" with respect to reference counts on the referenced
-- Ada and code versions.
-- Possible errors:
        o Union of those produced by Open and Create.
__
        o Illegal_Class
___
              o This operation is not supported for Class_Directory,
___
                Class_Universe_View, Class_World_View, Class_Link_Pack,
                Class_Tape, and Class_Terminal.
--
        O Object_Class_Mismatch
--
              o If the destination object already exists, it must have the
___
                same class as the source.
--
     Is_Version_Policy_Error (Status):
___
         o Object_Is_Not_Reserved
--
              o Can result from attempting to copy by reference in a slot
--
                (in the destination view) which does not have the
--
                Reserved bit set. Can result from the implicit calls to
                Create or Open.
generic
    with procedure Visit (Source: Om. Version_Handle;
                          Destination: Om. Version_Handle);
    with procedure Handle_Resolve_Error (Status : Error.Condition);
    with procedure Handle_Substitution_Error
                        (Status: Error.Condition;
                         Bad_Version : Om. Version_Handle);
    with procedure Handle_Copy_Error
                        (Status : Error Condition;
                         Attempted_Source : Om.Pathname_String;
                         Attempted_Target : Om.Pathname_String);
procedure wild_Copy (Source : Om.Full_wildcard;
                     Destination: Om.Full_Wildcard;
```

Action: Om.Action\_Id;

Source\_Context : Context := Job\_Default\_Context; Destination\_Context : Context := Job.Default\_Context; Max\_Wait : Duration := Job.Default\_Wait);

- -- Employs Naming.wild\_Resolve (using Source, Action, and Source\_Context)
- -- followed by Object.Resolve to produce a sequence of Version\_Handle's:
- -- errors are fed to Handle\_Resolve\_Error. For each Version\_Handle produced
- -- by resolve, employs Naming. Substitute (using the generated
- -- Version\_Handle, its pathname, and the Destination substitution pattern)
- -- to produce the target name; errors are fed to
- -- Handle\_Substitution\_Error. Calls the string form of copy for each
- -- generated pair of pathnames (using the full pathnames, Action, and
- -- Option); errors are fed to Handle\_Copy\_Error; calls Visit prior to each
- -- call to Copy.
- -- Properly handles creation of directories in the target context.
- -- What does it do with source universe and world views?
- -- RESTRICTION: not supported in 1st qtr impl.
- -- Possible errors:
- o Union of those produced by Naming.Wild\_Resolve,
- Naming. Substitute, Open, Create and Copy.

procedure Refresh\_Code\_Segment (Handle: in out Open\_Handle; Status : out Error Condition);

- -- Recall that Open will allow one to modify an existing code segment.
- -- Recall that code pages will lay around in the cache until reclaimed
- -- by the page replacement policy. Thus, in order to see the changes
- -- (that have been made in the heap segment) via the code segment, one
- -- must first employ this operation.
- -- Clearly this operation is a waste of time when applied to a new
- -- version of the code object.
- -- Possible errors:
- Is\_Bad\_Version\_Handle (Status)
  - o Version (Handle) does not exist probably a bad handle.
- -- Contents of an Open\_Handle --

-- The functions in this section will return the closest thing to Nil when -- given bad parameters or other errors occur.

function Version (Handle: Open\_Handle) return Om.Version\_Handle;

- -- Identifies the object, version and universe in which it was
- -- Create'd/Open'd. If Create/Open created a new version, identifies the
- -- newly created version.

function Predecessor\_Version (Handle: Open\_Handle) return Om. Version\_handle:

- -- If the Create/Open created a new version, then this function returns -- the identity of the version from which the new version was created,
- -- else Nil. Can be used to detect when the view is differentiated.

function Segment (Handle: Open\_Handle) return Machine.Segment\_Name;

- -- Identifies the memory addresses which will be used to manipulate the
- -- contents of the open version. Can use the Pointer\_To\_First\_User\_Bit
- -- function (in the segment operations, below) to get a typed pointer
- -- to the data in the segment.

function Prior\_Jpdate\_Id (Handle : Open\_Handle) return Om.Unique\_Id;

- -- Returns the value of Last\_Update\_Id as it existed just prior to the
- -- Open, or Nil when the Open created a new version. This information can
- -- be used to detect that the contents of an object have changed between
- -- an Open for read (perhaps supersedeable) and later Open/Create. A
- -- common application is to validate an editor image at first edit.

- -- These operations are restricted to internal OM and other "registered" -- clients; specifically these are not available to user programs.
- -- Read only object/version attributes '--
- -- The functions in this section will return the closest thing to Mil when
- -- given bad parameters or other errors occur. The procedures will be
- -- noops when given bad parameters or other errors occur.
- -- The procedures all take Open\_Handle's, which must have Write\_Object
- -- access to the desired version, else the operation is a noop.
- -- Note that if one applies a Get\_xxx operation to a version which happens
- -- to be open for update, one will receive the values which are current
- -- for the open generation.

function Exists (The\_Version: Om\_Version\_Id) return Boolean;

- -- If The\_Version.Number is Nil, returns true iff the specified object
- -- exists, else returns true iff the specified version exists. This is
- -- one of the few operations which will take Version\_Id\*s for job
- -- worlds and segments.

function Exists (The\_Version: Om.Version\_Handle) return Boolean;

- -- Returns true iff the specified version is the "current" version in the
- -- specified universe. If the version number component of The\_Version is
- -- Nil, this is equivalent to asking if the specified object is deleted
- -- (from the point of view of the given universe).

function Is\_Frozen (The\_Version : Om.Version\_Handle) return Boolean;

-- True iff the specified version exists and is either a frozen universe -- view, a frozen world view, or is referenced by one or more frozen world -- views.

function Version\_Count (The\_Object : Om.Object\_Handle) return Positive;

-- Returns the number of existing versions of the object.

-- Will return the Last\_Update\_Id of the previous version whose data bits -- were copied into this version when it was created. Very subtle note: -- when a write creates a new version to carry out its portion of the -- Supersedeable\_Read\_Object protocol, the new version keeps the same -- Predecessor\_Update\_Id value.

-- This information can be used by the configuration management tools to -- detect a situation in which a sublibrary is being reintegrated into its -- parent library, and the parent's unit has been recompiled, which might -- force the sublibraries's unit to be reintegrated in the source state.

-- Returns the network wide creation timestams for the identified version.
-- This is one of the few operations which will take Version\_Id's for job
-- worlds and segments.

function Disk\_Blocks\_Consumed (Version : Om.Version\_Handle) return Integer;

-- Number of 1K byte disk blocks actually consumed by the version itself.
-- This is one of the few operations which will take Version\_Id\*s for job
-- worlds and segments.

function Virtual\_Fage\_Length (Version : Om.Version\_Handle) return Integer:

-- Assuming no holes, indicates the number of pages which can can be
-- addressed (without getting Nonexistent\_Page\_Error). This is one of the
-- few operations which will take Version\_Id's for job worlds and segments.

function User\_Bit\_Length (Version : Om.Version\_Handle) return Integer;

```
-- Number of user bits stored in the version. Does not include ucode/04
-- overhead.
function Disk_Blocks_Reclaimable_By_Expunge
              (Version : Om. Version_Handle) return Integer;
-- Number of 1K byte disk blocks which would be reclaimed if the
-- version was expunded. For normal objects, same result as the
-- Disk_Blocks_Consumed function. But for universe views, world views,
-- and load images, includes the space that would be reclaimable as
-- a result of fewer references. This is one of the few operations which
-- will take Version_Id's for job worlds and segments.
-- Read/Write version attributes --
-- The functions in this section will return the closest thing to Nil when
-- given bad parameters or other errors occur. The procedures will be
-- noops when given bad parameters or other errors occur.
-- Most procedures take Open_Handle's, which must have Write_Object
-- access to the desired version, else the operation is a noop.
-- Note that if one applies a Get_xxx operation to a version which happens
-- to be open for update, one will receive the values which are current
-- for the open generation.
function Get_Subclass (Version : Om.Version_Handle)
              return Om.Subclass_Number;
procedure Set_Subclass (Handle : in out Open_Handle;
                        New_Subclass : Om.Subclass_Number);
function Acl (version : Om. Version_Handle) return Om. Access_Control_List;
procedure Set_Acl (Handle: Om.Version_Handle;
                   New_Acl : Om.Access_Control_List;
                   Status: out Error.Condition);
-- RESTRICTION: not supported in 1st qtr impl.
function Reserved (Version: Om. Version_Handle) return Boolean;
procedure Set_Reserved (Handle: in out Open_Handle;
                        New_Version : Boolean);
-- Used by configuration mgmt tools to control access to objects.
function Code_Name (Version : Om. Version_Handle) return Machine.Module_Name;
procedure Set_Code_Name (Handle: in out Open_Handle;
                         New_Name : Machine.Module_Name:
                         Status : out Error.Condition);
-- The procedure can only be used for MV generated code segs and requires
-- update access to every load image which references the open version.
-- Possible errors:
     Illegal_Class
         o Can only be applied to objects with Class_Code.
     Not_A_Cross_Cg_Segment
```

```
o World (New_Name) must be in Machine.Cross_Cg_Worlds.
--
     Not_In_Cross_Cg_Object_world
         o The associated world (of the given object) must be that
__
___
           associated with "!Machine".
--
     Not_A_New_Code_Segment
         o This must be a new version of the object not referenced by any
__
           load image.
function Last_Update_Time (Version : Om.Version_Handle)
              return Calendar.Time;
procedure Set_Last_Update_Time (Handle : in out Open_Handle;
                                Update_Time : Calendar.Time);
function Want_Backup (Version : Om.Version_Handle) return Boolean;
procedure Set_Want_Backup (Handle : in out Open_Handle;
                           New_Value : Boolean);
-- Versions are packed up iff the Want_Backup flag is set. The default
-- is Want_Backup => true.
-- Ada specific read/write version attributes --
-- The functions in this section will return the closest thing to Nil when
-- given bad parameters are other errors occur. The functions do NOT
-- get action locks, thus garbage results are theoretically possible.
-- The procedures will get a Record_Write lock on the attribute.
function Unit_State (Version: Om.Object_Handle) return Om.Ada_Jnit_State;
generic
    with procedure Make_Change (Unit_State : in out Om.Ada_Unit_State);
procedure Set_Unit_State (Version : Om.Object_Handle;
                          Status : out Error Condition);
-- The function returns the current state of the identified unit.
-- The procedure can make arbitrary changes to the unit state, via the
-- supplied Make_Change operation, acquiring a Record_write lock on
-- the unit state of the identified unit.
function Depends_On (Object : Om.Object_Handle)
              return Um.Dependency_Vector:
generic
    with procedure Make_Change (Depends_On : in out Om.Dependency_Vector);
procedure Set_Depends_On (Object : Om.Object_Handle;
                           Status: out Error.Condition);
-- The function returns the current value of the identified row of the
-- Dependency_Matrix. The procedure can make arbitrary changes to the row,
```

-- via the supplied Make\_Change operation, acquiring a Record\_Arite lock on

function Dependents (Object: Om.Object\_handle) return Om.Dependency\_Vector:

-- the row for the identified unit.

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-- The function returns the current value of the identified column of the
-- Dependency_Matrix. There is no operation to change a column of the
-- matrix, without resorting to the change row operation (above).
-- Relations between objects --
generic
    with procedure Visit (The_Version : Om.Version_Handle);
procedure Traverse_Versions (Of_Object : Om.Object_Handle;
                             Status: out Error.Condition);
-- If Of_Object.Universe is non-Nil and it selects a world view, then this
-- operation calls Visit for every version Of_Object in the specified
-- universe. If the slot (in the world view) has non-Nil entries, Visit
-- will be called for trailing Nil's. For example, if the retention list
-- is (nil, nil, v1, v4, v7, nil) (the versions are listed in the order
-- that they were pushed), Visit will be called 4 times, for v1, v4, v7
-- and nil, in that order.
-- Otherwise (.Universe is Nil, or doesn't select a world view), calls
-- Visit for every version Of_Object, regardless of what views they
-- participate in. Visit is called in sort order, by Version_Number.
-- Possible errors:
     Is_Bad_Version_Handle (Status)
        o Something wrong with Of_Object
generic
    with procedure Visit (Referencing_Version : Om.Version_Handle);
    with procedure Handle_Error (Status : Error.Condition);
procedure Locate_Expunge_Preventing_References
     (The_Version : Om.Version Handle:
      Status: out Error.Condition);
-- The operation finds the expunge-preventing references to the specified
-- version of the object.
-- Calls Visit for every expunge-preventing reference. Depending upon the
-- class of The_Object, Referencing_Version may refer to a universe view,
-- world view, or load image. It is possible that some of the information
-- may not be available due to locks held by other actions; calls
-- Handle_Error in these cases. As usual, the caller is free to raise
-- arbitrary exceptions in the generic formal subprograms; these terminate
-- the search.
-- This operation is useful for determining why a particular version fails
-- to be expunyed.
-- Possible errors:
--
     Is_Bad_Version_Handle (Status)
___
     Is_Lock_Error (Status);
     Is_kesource_Limit_Error (Status)
```

```
generic
    with procedure Visit (The_Object : Om_Object_Handle);
procedure Locate_Objects_With_Given_Class
     (In_World : Om.version_Handle;
      Class: Om.Class_Number;
      Status: out Error.Condition);
-- The world number component of In_World is used to identify the
-- world to be searched. Calls Visit for every object (in the specified
-- world) which has the specified Class.
-- Possible errors:
     Is_Bad_Version_Handle (Status)
     Illegal_Class
generic
    with procedure Visit (The_Universe_View : Om.Version_Handle);
procedure Locate_Referencing_Universes
     (The_world_View : Om.Version_Handle;
      Status: out Error.Condition);
-- Looks through all worlds (on this machine), calling Visit for every
-- universe view which references the specified world view version.
-- Possible errors:
     Is_Bad_Version_Handle (Status)
     Illegal_Class
-- Class and Subclass names, and registration --
function To_String (Class: Om.Class_Number) return String;
-- Serves as the 'Image function for Om. Class_Number. Returns "(undefined)"
-- for bogus input.
function To_String (Class: Om.Class_Number;
                    Subclass: Om. Subclass_Number) return String;
-- Serves as the 'Image function for Om. Subclass_Number. Note that subclass
-- numbers are defined in the context of a particular class. Returns
-- "(undefined)" for bogus input.
package Directory_Subclass is
    function World return Om.Subclass_Number;
    function rath return Om.Subclass_Number;
    function Session return Om.Subclass_Number;
end Directory_Subclass;
--pragma Integrate (Directory_Subclass);
package File_Subclass is
    -- The following subclasses are predefined:
```

function Finary return Om.Subclass\_Number;

```
page
                                                                           21
    function Text return Om.Subclass_Number;
    function Message_Window return Om.Subclass_Number;
    function Text_Io_window return Om.Subclass_Number;
    function Switches return Om. Subclass_Number:
    function Editor_Macros return Om.Subclass_Number;
end File_Subclass;
--pragma Integrate (File_Subclass);
package Ada_Attribute_Subclass is
    -- The following subclasses are predefined:
    function Compatibility_Database return Om.Subclass_Number;
    function Ada_Image return Om.Subclass_Number:
    function Cg_Attributes return Om.Subclass_Number;
end Ada_Attribute_Subclass;
--pragma Integrate (Ada_Attribute_Subclass);
procedure Define_Subclass
               (Class: Om.Class_Number;
                Subclass : Om.Subclass_Number:
                Name : String;
                Policy : Om. Version_Control_Policy := Om. Vil;
                Status: out Error.Condition);
-- Used to define a new subclass. If the specified Default_Version_Policy
-- is Nil, will implicitly use the value returned by the
-- Default_Version_Policy function, below. The new subclass definition
-- is now permanent. Note that the assignment of subclasses is typically
-- done via a manual registration process, at Rational.
-- Possible errors:
--
     Object_Class_Does_Not_Exist
__
     Object_SucClass_Already_Exists
--
         o The given Subclass and Name may not previously exist.
     Is_Version_Policy_Error (Status):
__
         o Undefined_Version_Policy
procedure Undefine_Subclass
               (Class: Om.Class_Number;
                Subclass : Om.Subclass_Number;
                Name : String;
                Status: out Error.Condition);
-- Used to delete a previously defined subclass. Does not verify that
-- there are no objects which claim to have the now deleted subclass.
-- Possible errors:
     Object_Class_voes_Not_Exist
     Object_SubClass_Does_Not_Fxist
     Object_SubClass_Is_Predefined
function Default_Version_Policy
              (Class: Om. Class_Number;
               Subclass: Om. Subclass_Number := Om. Nil)
              return Om. version_Control_Policy;
-- Returns Standard_Differential_Version_Policy for all predefined
-- subclasses, including vit. Exceptions:
```

```
-- (1) The Compatibility_Database subclass of Ada_Attribute's, uses the
      Shared_Version_Policy.
-- (2) Directory class uses the Shared_Version_Policy.
-- (3) Universe view class uses the Shared_Version_Policy.
-- (4) Device classes use the Shared_Version_Policy.
-- Segment operations:
    Restricted to Rational programs only --
-- The following operations assume that one obtained the Segment_Name
-- argument by using the Segment function on a legitimate Open_Handle.
-- Segment operations fed bad Segment_Name parameters will raise the
-- following exceptions:
--
       Illegal_Reference
--
           o The Module_Name portion of the given Segment_Name is 0.
____
       Nonexistent_Page_Error
--
           o The Module_Name portion of the given Segment_Name does not
__
              exist.
       Itlegal_Heap_Access
           o Attempt to address outside the range of "user data bits".
function Corresponding_Version (The_Segment : Machine.Segment_Name)
              return Om.Version_Handle;
-- Returns a Version_Handle with a fully resolve Version component, and
-- nil Universe component. May raise the standard bad Segment_Name
-- exceptions, as defined above.
function End_Of_Segment (The_Segment : Machine.Segment_Name)
              return Machine.Bit_Offset;
function End_Of_Segment (The_Segment: Machine.Segment_Name)
              return Om.Segment_Position:
-- Returns the identity of the first unused "offset" in the segment.
-- The first function returns the value in terms appropriate for comparison
-- with the Bit_Offset component of Machine. Address's. The second function
-- returns the value in terms of "user data bits". May raise the standard
-- bad Segment_Name exceptions, as defined above.
function Virtual_Length (Of_Segment : Machine.Segment_Name)
              return Long_Integer;
function Logical_Length (Of_Segment : Machine.Segment_Name)
              return Long_Integer;
-- The Virtual_ength of a segment is the number of bits between the
-- first and last legal Bit_Offset's, inclusive. The Logical_Length of
-- a segment is the number of bits between the first and last legal
-- Segment_Position's, inclusive. By "legal" we mean that if one had
-- the segment open for update, one could "touch" the addressed bit.
```

-- hay raise the standard and Segment\_Name exceptions, as defined above.

```
23
generic
   type Element is limited private;
    type Pointer is access Element:
    pragma Segmented_Heap (Pointer);
function Get_Segment (P : Pointer) return Machine.Segment_Name;
-- May raise the standard bad Segment_Name exceptions, as defined above.
generic
   type Element is limited private;
    type Pointer is access Element:
    pragma Segmented_Heap (Pointer);
function Get_Position (P : Pointer) return Om.Segment_Position;
-- May raise the standard bad Segment_Name exceptions, as defined above.
-- Segment operations:
    Extremely dangerous operations --
     Restricted to "trusted code" --
generic
   type Element is private;
    type Pointer is access Element;
   pragma Segmented_Heap (Pointer);
function Pointer_To_First_User_Bit
     (Into: Machine.Segment_Name) return Pointer;
-- Used to return a typed pointer to the first "user bit" in the segment.
generic
    type Element is limited private;
    type Pointer is access Element;
    pragma Segmented_Heap (Pointer);
function Make_Pointer (Into : Machine.Segment_Name;
                       At_Position : Om.Segment_Position) return Pointer;
-- Can be used to construct pointers from their components. Since one
-- can make pointers at arbitrary positions, the resulting equivalencing
-- can be the source of severe damage to the system. One can make a null
-- pointer by supplying Nil for At_Position. May raise the standard bad
-- Segment_Name exceptions, as defined above.
generic
    type Element is limited private;
    type Pointer is access Element:
    pragma Segmented_Heap (Pointer);
function Convert_To_Pointer (Into : Machine.Address) return Pointer;
-- Can be used to construct pointers from its address. Since one can make
-- pointers at arbitrary positions, the resulting equivalencing can be
-- the source of severe damage to the system. Will raise
-- Illegal_Heap_Access if the given address refers to an object with
-- other than Class_File, Class_Ada, Class_Ada_Attributes, or Class_Code.
generic
```

```
type Element is limited private;
       type Pointer is access Element:
       pragma Segmented_Heap (Pointer);
   procedure Assign_Value (Value : Element;
                            Into : Pointer);
   -- For scalars and records without discriminants, this operation is the
   -- same as "Into.all := Value". In particular, if the Element type is
   -- scalar, will raise Contraint_Error and Numeric_Error, as would be
   -- expected from "Into.all := Value".
   -- In the cases where the Element type is an array or record with
   -- discrimants, the semantics are completely different than "Into.all :=
   -- Value"; in this case, the operation does not generate Constraint_Error;
   -- in particular, it simply copies the bits of the source (by looking
   -- inside the source object, to determine 'size) to the memory addressed
   -- by the given pointer; but does not look inside the target (to check
    -- array bounds, discriminants, constraints, etc); thus it simply assumes
   -- that *size bits have already been allocated at the target address.
   -- May raise the standard bad Segment_Name exceptions defined above.
   procedure Adjust_End_Of_Segment (Of_Heap : Machine.Segment_Name;)
                                     New_End_Of_Segment : Om.Segment_Position):
   -- Causes the End_Of_Segment operation to return New_End_Of_Segment. When
   -- New_End_Of_Segment is larger than the current End_Of_Segment, this
   -- operation "extends the heap". When New_End_Of_Segment is smaller than
    -- current End_Of_Segment, this operation "cuts back the heap". May raise
   -- the standard bad Segment_Name exceptions defined above.
private
    type Open_handle is
       -- For details, see the section containing functions on Open_Handle.
        record
            Version
                                : Om. Version_Handle;
            Predecessor_version: Om.Version_Id:
                                : Machine.Segment_Name:
            Segment
                                : Action.Object_Lock_Handle;
            Lock_Handle
            Version_Map_Pointer : Machine.Address:
        end record;
    Nil_Open_Handle : constant Open_Handle :=
                    (Version => Om.Nil_Object_Handle,
                     Predecessor_Version => Om.Nil_Version_Id,
                     Segment => Machine.Null_Segment_Name,
                     Lock_Handle => Action.Nil_Object_Lock_Handle,
                     Version_Map_Pointer => Machine.Null_Address);
```

end Object:

```
GGG
       cded
               AAA
G
      Ρ
         P
  G
              A A
G
       Р
           Ρ
G
      PPPP
  GG
              AAAAA
G
   G
      P
                  Α
              A
G
   G
       P
              A
                  Α
GGGG
       P
                  A
```

User: GPA
Object: !DE

: !DELTA\_KK.REV1\_0\_0.UNITS.KK.MACHINE

Version: V(14) Request: 1273

Date: April 24, 1986 Queued: 11:26:36 AM

Printed: 11:39:09 A4

```
!DELTA_KK.REV1_0_0.JNITS.KK.MACHINE *V(14)
```

```
page
with Sys;
package Machine is
   -- Gamma "Machine" package to die. This package is part of the
    -- standard OM export view and contains machine dependent definitions,
   -- mostly addressing conventions.
    -- pragma Subsystem (Kernel);
    -- pragma Module_Name (4, ?);
    package System renames Sys;
    -- Basic storage units --
    ------
   Bits_Per_Byte
                            : constant := 8;
   Bits_Per_Word
                             : constant := 128;
    Bits_Per_Page
                              : constant := 2 ** 13;
    Bits_Per_Page_Is_Constant : Boolean := True;
    -- In the future, the first value may be a function and change from
    -- machine to machine, perhaps from boot to boot. Assumptions to the
    -- contrary should assert the _Is_Constant invariant.
    --=========
    -- Segments --
    -------------
    -- Taken together, World_Number, Segment_Number and Segment_Kind uniquely
    -- identify a segment in the machine. A segment is the fundamental unit
    -- of storage in the virtual memory system. Kinds Illegal and System are
    -- not used.
    -- Note that in Job_worlds Control, Typ, Queue, and Data segments have the
    -- same lifetime, and are collectively referred to as a module. Import and
    -- Code segments each have distinct lifetimes and are typically referred
    -- to as import spaces and code segments, respectively.
    -- World_Number --
    type world_Number is new Long_Integer range 0..2 ** 10 - 1;
    subtype Kernel_Worlds is World_Number range 0...4;
    -- All Kernel_Layer modules are elaborated in segments in worlds
    -- in this range.
    Ipl_world : constant World_Number := 4;
    subtype Environment_worlds is World_Number range 4..5;
```

-- The module and import spaces of the environment subsystems are -- elaborated in these job worlds. When the virtual memory system is -- initialized, the above job worlds are automatically created.

subtype Env\_backup\_worlds is World\_Number range 6..7;

- -- After a system crash, the stuff in world 4 is "copied" to world 6,
- -- and world 5 "copied" to world 7. This allows virtual memory to be
- -- started and the system to elaborate prior to dumping the state of -- virtual memory associated with the last crash.

subtype Job\_Worlds is world\_Number range 8..255;

- -- Jobs run in worlds in this range. Module amd import names may be
- -- reclaimed when the job terminates and all traces of it have been wiped
- -- from the face of the earth. Until we stop building environment systems
- -- on the MV's, we will not actually give out job numbers in the range
- -- Cross\_Cg\_Worlds.

subtype Cross\_Cg\_Worlds is world\_Number range 8..25;

- -- It is still the case that mv generated code segments live on vps 8..25.
- -- These pages are always read\_only and clean. There is 40 corresponding -- object world.

subtype Iop\_Cross\_Cg\_Worlds is Cross\_Cg\_Worlds range 8..10;

- -- These are the Cross\_Cg\_Worlds that are loaded from the IOP.
- -- \*\*\* These numbers are not the real ones. They are yet to be decided on.

subtype Env\_Cross\_Cg\_Worlds is Cross\_Cg\_Worlds range 11..25;

- -- These are the Cross\_C3\_Worlds that are loaded into EED3.
- -- \*\*\* These numbers are not the real ones. They are yet to be decided on.

subtype Task\_Worlds is World\_Number range 0..255;

-- Worlds in which the architecture/mts are willing to run tasks.

subtype Object\_Worlds is World\_Number range 256..1022;

-- Segmented heaps (file spaces) live in worlds in this range. These may -- be either permanent and temporary.

Om\_World : constant Object\_Worlds := 256; Miscellaneous\_World: constant World\_Number := 1023;

- -- Om\_World contains object management state which is not accounted for by
- -- legitimate directory objects. There is a corresponding disk world.
- -- Current uses: (a) Action log blocks.
- -- There is no corresponding disk world for Miscellaneous\_World. It is
- -- used for addresses of special gook in the memory cache. The following
- -- things are placed in the cache under these addresses: (1) The ucode
- -- assist routines are trapped by doing field executes into this world.
- -- (2) The kernel's page pool pages go in this world. (3) Some of the
- -- Volume\_hanager state is mapped via this world.
- -- Beware of the fact that the window of vulnerability mechanism, actions,
- -- and crash recovery all assert the following invariant:
- Disk.All\_Dirty\_world\_Info\_Mapped\_Under\_Its\_Vp
- -- Moral of the story: be VERY careful with permanent/reconstructed state

```
-- which lives in the cache under these world numbers.
------
-- Segment_Number --
_____
type Segment_Number is new Long_Integer range 0..2 ** 22 - 1;
subtype Code_Segment_Number is Segment_Number
     range 0...Segment_Number (2 ** 16 - 1);
First_Allocatable_Segment : constant := 2 ** 6;
-- For all worlds (both job and file), the first 2**6 segment numbers
-- are reserved for use as disk mapping pages by Object_Management.
type Segment_Kind is new Long_Integer range 0...7;
Illegal_Segment : constant Segment_Kind := 0;
Control_Segment : constant Segment_Kind := 1;
Typ_Segment : constant Segment_Kind := 2;
@ueue_Segment : constant Segment_Kind := 3;
Data_Segment : constant Segment_Kind := 4;
Import_Segment : constant Segment_Kind := 5;
Code_Segment : constant Segment_Kind := 6;
System_Segment : constant Segment_Kind := 7;
type Module_Name is new Long_Integer range 0..2 ** 32 - 1;
-- The architecture often knows the Segment_Kind implicitly (from
-- context) and therefore uses Module_Name to refer to segments. The
-- following operations define the representation of a Module_Name.
function Extract_Segment (Name : Module_Name) return Segment_Number;
function Extract_World (Name : Module_Name) return World_Number;
function Construct_Module_Name
              (Segment : Segment_Number:
               World: World_Number) return Module_Name;
-----------------
-- Addresses --
--=========
type Bit_Offset is new Long_Integer range 0..2 ** 32 - 1;
-- Physical bit offset within a particular segment; O identifies first
-- physical bit in the segment.
subtype Address is System. Address:
Null_Address: constant Address:= System.Null_Address;
-- Addresses carried around by the machine are of this form:
-- type Address is
__
       record
           Segment : Segment_Number:
           World : World_Number;
```

Page

: Page\_Number:

```
!DELTA_KK.REV1_0_0.JNITS.KK.MACHINE *V(14)
```

۷,

```
Byte
                  : 0 -- 1023;
___
           Bit
                   : 0 .. 7:
--
       end record;
-- or, equivalently:
--
      record
--
           Name : Module_Name:
--
           Offset : Bit_Offset;
       end record;
-- Warning: an Address does NOT contain the segment kind; it must be
-- known implicitly. The following operations define the representation
-- of addresses.
function "<" (L, R : Address) return Boolean;
function "=" (L, R : Address) return Boolean renames System."=";
function Extract_Segment (Addr : Address) return Segment_Number;
function Extract_World (Addr : Address) return World_Number;
function Extract_Name (Addr : Address) return Module_Name;
function Extract_Offset (Addr : Address) return Sit_Offset;
function Construct_Address (Segment : Segment_Number;
                            world: World_Number;
                            Offset: Bit_Offset) return Address;
function Construct_Address (Name : Module_Name;
                            Offset: Bit_Offset) return Address;
type Page_Number is new Long_Integer range 0..2 ** 19 - 1;
function Extract_Page (Addr : Address) return Page_Number;
function Extract_Page (Offset : Bit_Offset) return Page_Number;
function Extract_Bit_Offset_Within_Page (Addr : Address) return Bit_Offset;
function Extract_Bit_Offset_Within_Page
              (Offset : Bit_Offset) return Bit_Offset;
function To_Long_Integer (Addr : Address) return Long_Integer;
function To_Address (Int : Long_Integer) return Address;
-- These are inverses of each other, and preserve bit representation.
subtype Page_Address is Address;
-- Same as an Address, except: bits 51..60 are 0, and bits 61..63 contain
-- the Segment_<ind; ie, it looks like:
--
       record
__
           Segment : Segment_Number;
___
           World : World_Number;
--
          Page
                 : Page_Number;
__
           5yte
                  : 0 .. 1023; -- all zero
--
           Kind
                 : Segment_Kind; -- in lieu of Bit #
--
       end record;
subtype Space_Address is Address:
-- Same as a page_Address, except: bits 32..50 are 0; ie, it looks like:
      record
__
           Segment : Segment_Number:
___
           World : World_Number;
___
           Page
                  : Page_Number: -- all zero
--
           Eyte
                  : 0 .. 1023; -- all zero
--
           Kind
                   : Segment_Kind: -- in lieu of Bit #
--
       end record:
-- Unlike the Gamma system, each segment is considered to have its own
```

-- space address. function Construct\_Space\_Address (Name : Module\_Name; Kind: Segment\_Kind) return Space\_Address; function Construct\_Page\_Address (Name : Module\_Name; Page : Page\_Number; Kind: Segment\_Kind) return Page\_Address; function To\_Data\_Address (Addr : Address) return Page\_Address; function To\_Address (Addr : Address; Kind : Segment\_Kind) return Page\_Address; function Extract\_Kind (Addr : Space\_Address) return Segment\_Kind; function Change\_Kind (Of\_Addr : Space\_Address; To\_Kind: Segment\_Kind) return Space\_Address: function Space (Of\_Page: Page\_Address) return Space\_Address; function Change\_Page\_Number (Of\_Page : Page\_Address; To\_Page\_Number : Page\_Number) return Page\_Address: type Address\_Array is array (Natural range <>>) of Address; -- Protected segment reference --type Segment\_Name is private; -- Same information content as a Module\_Name, but this form has the -- bits of the Module\_Name left justified, to be in the same position -- as in an Address. Also, this type is runtime private in the sense -- that the machine will zero its contents when declared and will not -- permit aliasing. Thus, one will not find any functions which can -- generate this type from integers (or other unprotected types). Null\_Segment\_Name : constant Segment\_Name; function To\_Long\_Integer (Seg : Segment\_Name) return Long\_Integer; function To\_Module\_Name (Seg : Segment\_Name) return Module\_Name; -- Basic Machine Operations --subtype Task\_Id is Module\_Name; -- Refers to a module, which can be a package or Ada task.

-- "task id" of the executing thread. Note that packages also have

function Get\_Machine\_Id return Long\_Integer;

function My\_Task\_Id return Task\_Id;

-- task ids.

```
-- Returns the id for this particular R1000. This value is known to not
-- exceed 10**6.
type Cpu_Priority_Level is new Long_Integer range 0..15;
Best_Cpu_Priority
                        : constant Cpu_Priority_Level := 0;
Next_Best_Cpu_Priority : constant Cpu_Priority_Level := 1;
Third_Best_Cpu_priority : constant Cpu_Priority_Level := 2;
Fourth_Best_Cpu_Priority : constant Cpu_Priority_Level := 3;
Worst_Cpu_Priority
                       : constant Cpu_Priority_Level := 15;
-- Note that these notions of priority reflect the internal scheduling
-- algorithms of the macine. Task priority (per LRM) is defined elsewhere.
-- Job priority (per MTS) is defined elsewhere.
function My_Priority return Cpu_Priority_Level;
-- If called while in rendezvous, will return that of the rendezvous, not
-- of the task calling the function.
function I_Am_wired return Boolean;
-- Returns true if the executing task's TCB is wired.
function Is_Callable (Tne_Task : Task_Id) return Boolean;
-- Returns the value of the 'Callable attribute (as per LRM).
package Memory is
   -- Memory related operations --
    function Number_Of_Sets return Natural;
    function Number_Of_Lines return Natural;
    function Number_Of_Frames return Natural;
    type Cache_Set_Number is new Natural range 0..Number_Of_Sets - 1;
    type Cache_Line_Number is new Natural range O..Number_Of_Lines - 1;
   type Cache_Frame_Number is new Natural range 0..Number_Of_Frames - 1;
   type Frame_Bits is array (Cache_Frame_Number) of Boolean;
    function Line (Frame : Cache_Frame_Number) return Cache_Line_Number;
    function Set (Frame : Cache_Frame_Number) return Cache_Set_Number;
    function Frame (Line : Cache_Line_Number;
                   Set : Cache_Set_Number) return Cache_Frame_Number;
   procedure Get_Page_Address (In_Frame : Cache_Frame_Number;
                               Empty: out Boolean;
                               Contains_Page : out Page_Address);
   -- If the specified cache frame contains no page, then return Empty;
   -- else not Empty and Contains_Page gives the virtual address of the
   -- page in the specified frame.
    function Exists (Page: Page_Address) return Boolean;
```

```
-- Returns true iff the specified page is currently in the cache.
   function Is_In_Transit (Page : Page_Address) return Boolean;
   -- Returns true iff the specified page exists in the cache and is
   -- in_transit.
   function Is_Dirty (Page : Page_Address) return Boolean;
   -- Returns true iff the specified page exists in the cache and is
   -- marked as dirty.
   function Is_Wired (Page : Page_Address) return Boolean;
   -- Returns true iff the specified page exists in the cache and is
   -- wired.
end Memory:
package Operations is
   -- Miscellaneous ucode'd operations --
   type Value_Generator is
       record
           Value : Long_Integer;
       end record:
   procedure Indivisibly_Increment (Generator: in out Value_Generator;
                                   Value : out Long_Integer);
   -- Performs the following:
         if Generator. Value < Long_Integer*Last then
             Generator. Value := Generator. Value + 1;
         Value := Generator.Value;
    -- This operation is indivisible (atomic), even with respect to page
    -- faults.
    type Boolean_Lock is
       record
           Value : Boolean;
       end record;
    procedure Test_And_Set (The_Lock: in out Boolean_Lock;
                           Value : out Boolean);
    -- Performs the following:
         value := The_Lock.Value;
         The_Lock.Value := True:
    -- This operation is indivisible (atomic), even with respect to page
    -- faults.
    procedure wait_For_Transfer_To_Complete (Page : Page_Address);
```

```
-- Used to wait for disk transfer completion: Caller should first -- issue the appropriate request to the Disk_Driver. Then call this -- routine with the page involved in that transfer. If the specified -- Dage does not exist or is not currently in transit, this call -- returns immediately; otherwise places the calling task in the disk -- wait queue, in which case the call returns when the appropriate -- invocation of Roust has been made by Disk_Driver.
```

## end Operations;

```
private
```

type Segment\_Name is access Boolean;
pragma Segmented\_Heap (Segment\_Name);
Null\_Segment\_Name : constant Segment\_Name := null;
end Machine;

```
PPPP
GGG
              AAA
      P
G
  G
        P
             A A
6
      Ρ
G GG
      PPPP
             AAAAA
G G
      P
                 Α
             A
      P
             Α
                 A
GGGG
      Р
                 A
```

User: GPA

Object: !DELTA\_KK.REV1\_0\_0.UNITS.OM.EXECUTION

Version: V(104) Request: 1274

Date: April 24, 1986

Queued: 11:25:50 AM Printed: 11:39:41 AM

```
!DELTA_KK.REV1_0_0.JWITS.OM.EXECUTION V(104)
with Job:
with Error:
```

page 1

```
with Object:
with Machine;
with Om_Definitions;
pragma Private_Eyes_Only;
with Execution_Layout;
with Load_Image_Layout;
package Execution is
    -- pragma Subsystem (Object_Management);
    -- pragma Module_Name (4, ?);
    -- Entities to be executed include main programs and subsystems, both are
    -- represented by load image objects (produced by the "loader"). The
    -- contents of a load image object is defined by the Load_Image backage
    -- (below). There are various operations which cause the execution
    -- (elaboration) of load images (representing main programs and
    -- subsystems). These are defined in the Operations package (below).
    package Om renames Om_Definitions;
    subtype Compatability_Vector is Om.Compatability_Vector;
    subtype Code_Segment_key is Natural range 0..2 ** 16 - 1;
    subtype Secondary_Number is Natural range 0..2 ** 16 - 1;
    subtype Imported_Unit_Number is Natural range 0..2 ** 16 - 1;
    subtype Imported_Subsystem_Number is Natural range 0..2 ** 16 - 1;
    subtype Imported_Program_Number is Natural range 0..2 ** 16 - 1;
    subtype Exported_Module_Fosition is Natural range 0..2 ** 16 - 1;
    subtype lmported_Module_Position is Natural range 0..2 ** 16 - 1;
    subtype Elaboration_Order_Array is Om.Simple_Object_Id_Array;
    subtype Exception_Information is Om_Exception_Information;
    subtype Root_Of_world is Om_Object_Handle;
    subtype Some_Object_In_World is Om.Object_Handle;
    -- Load images --
    package Load_Image is
        -- Recall the definition of "partial order defined by unit with
        -- clauses" in section 10.3 of the LRM. we use the term "unit with
        -- closure" to mean the LRM's partial order, treating the entire
```

- -- directory system as one big Ada library.
- -- Here we use the term "main program" to mean a library unit procedure -- which is (to be) loaded as a main program.
- -- Here we use the term "subsystem" to mean a world which is (to be) -- loaded as a subsystem (in the current view).
- -- One job of the loader is to generate code to cause library unit
- -- elaboration according to the semantics of section 10.5 of the LRM.
- -- The set of units (to be elaborated when the load image is

\_\_

--

-- executed) is calculated as follows:

- (1) For a main program, include every unit in the main program's closure, stopping at (and not including) units which are (i) in worlds loaded as subsystems or (ii) other main programs.
- (2) For a subsystem, for each "selected" library unit in the subsystem, include its unit closure, stopping at (and not including) units which are (i) in other worlds loaded as subsystems or (ii) main programs. The "selected" units could be specified in a variety of ways; examples include: (a) every unit in the world view, (b) every unit in some export list, (c) units specified by some wildcarded pathname; how they are specified is immaterial to this package.
  Note that, in both cases, this set of units may include units that
- -- Note that, in both cases, this set of units may include units that
  -- are in worlds other than the world containing the loaded main
  -- program or subsystem.
- -- The world reconstruction algorithms allow load images to reference -- only those code segments which live in the same world as the load -- image. The purpose of the primary/secondary mechanism is to allow -- the loading of a subsystem or main program across world boundaries, -- in spite of this restriction.
- -- Load images contain the following information:
- (1) Ignoring the primary/secondary issues, a load image references
   the set of code segments which contain the code for the set
   of units elaborated by executing the load image.
- Code segment names are NOT stored in code segments themselves, for the following reasons:
- First, code segment names are local to the machine on which the segment was created. Therefore, keeping code segment names in just the load image object simplifies the process of moving a loaded program/subsystem between machines.
- Second, code segment names are treated as protected objects by OM (for access control reasons); this would be considerably more difficult if code segment names could be stored in code segments (which would imply the ability to create them from integers).
- Third, expunge control for code segments follows these rules: A code segment cannot be expunged if there is a load image which references the code segment. A load image cannot be -expunged if it is Open. The execution mechanisms keep the ---load image open as long as there is an elaborated copy of the \_\_\_ load image. And the execution mechanisms prevent unelaboration while importers still exist. This acheives the invariant that a code segment cannot be expunded while being executed. \_ -Violating this invariant can cause executing programs to get - strange exceptions (like Non\_Existant\_Page\_Error) and risks \_\_ crashing the machine.
- -- (2) The load image stores a map from original Object\_Index to
  the Version\_Handle of the corresponding Diana tree. Debug
  tables use these object indexes to refer to units. Use of
  Version\_Handle allows tools to use the universe component

3,

- to record the identity of the view from which the unit
  was loaded. An option is available to expunge protect the
  referenced trees.
- -- (3) for subsystems, the load image stores export Compatability\_Keys for all of the units exported by the subsystem.
- -- (4) For subsystems, the load image identifies the set of exported runtime modules (corresponding in a perhaps package intergrated fashion to the exported units).
- (5) The load image identifies the set of subsystems which are imported by the given main program or subsystem. And for each imported subsystem, stores the import Compatability\_Keys for the imported units. At execution time, these import keys are matched against the exporter's export keys.
- -- (6) The importing load image specifies where (in its import space) it wants to find the module names of the imported runtime modules.
- -- (7) The load image identifies the set of main programs which it may call.
- -- (8) The elaboration order (chosen by the loader) is recorded in the load image. This may be useful information for the user. But the major purpose is so the subsystem integrator (discussed further below) can produce a single output load image from many input load images, and preserve elaboration order.
- -- In the event that a loaded main program or subsystem needs to elaborate units in more than one world, the primary/secondary -- mechanism is used. A primary load image is created in the world -- containing the main program or subsystem. Secondary load images -- are created in the other worlds. The primary references all of -- the secondaries. The secondaries name the primary. As long as -- the primary continues to exist, the secondaries will not (by -- default) be expunged. Each load image references the loaded -- code segments in the same world.
- -- Functions in this package return garbage (typically 0) when given -- garbaye parameters. Procedures which do not return status are -- noops when given garbage parameters.
- -- The functions require the handle to be open for either Read\_Object -- or write\_Object. The modifying procedures require the handle to -- be open for write\_Object.

-- Possible errors:

## Status : out Error.Condition);

```
-- This operation must be the first operatino to be applied to a
-- newly created version of a load image. Other operations are used
-- to "fill in the shape".
-- When Is_Primary is true, Number_Of_Secondaries indicates the
-- exact number of secondary load images which will be referenced by
-- this primary. Recall that secondary load images are only used when
-- the load closure spans worlds.
-- Note that the operations below which set string values will
-- generate garbage (in the load image) if the string value is set
-- more than once. The garbage cannot be reclaimed without expunging
-- the version.
-- Possible errors:
     Is_Execution_Error (Status)
           o Not_A_Load_Image
           o Load_Image_Is_Too_Big
           o.Exported_Module_Count_Out_Of_Bounds
procedure Set_Code_Segment
               (The_Load_Image : in out Object.Open_Handle;
                key : Code_Segment_Key;
                Code_Version : Om.Version_Handle;
                Status : out Error.Condition);
-- The Object.Code_Name function is used (on Code_Version) to set both
-- the Current_Segment and Original_Segment attributes. The
-- Code_Version attribute is assigned the value of the Code_Version
-- parameter.
-- Possible errors:
     Is_Execution_Error (Status)
__
          o Not_A_Load_Image
___
          o Code_Segment_Key_Out_Of_Bounds
          o Not_A_Code_Version
                The Code_Version parameter does not refer to an existing
___
--
                version of an object of Class_Code.
          o Must_Be_In_Load_Image_World
                Code_Version must be in the same world as the load
_--
__
                image.
procedure Set_Original_Segment
               (The_Load_Image: in out Object_Open_Handle;
                Key: Code_Segment_Key:
                Original_Segment : Machine.Module_Name);
-- Used by tool to move a load image between machines.
-- The Original_Segment and Current_Segment attributes might be
-- different (for a given key value) because the load image and the
-- corresponding code segments were moved from the producing
-- (original) machine to the current machine. Retaining
-- Original_Segments may be useful for debugging.
```

```
Is_Execution_Error (Status)
           o Not_A_Load_Image
           o Code_Segment_Key_Out_Of_Bounds
function Original_Segment (The_Load_Image: Object.Open_Handle;
                           Key : Code_Segment_Key)
              return Machine_Module_Name;
function Current_Segment (The_Load_Image : Object.Open_Handle;
                          Key : Code_Segment_Key)
              return Machine.Module_Name;
function Code_Version (The_Load_Image : Object_Open_Handle;
                       Key : Code_Segment_Key)
              return Om. Version_Id;
-- The above functions are used to return the values of the attributes
-- set by the Set_Code_Segment operation. Note that there is no
-- visible operation to retrieve the "protected form" of the code
-- segment name. The elaboration operations use internal facilities
-- which return the protected form and may generate access control
-- error conditions.
function Protect_Debug_Units (The_Load_Image : Object.Open_Handle)
              return boolean;
procedure Set_Protect_Debug_Units
               (The_Load_Image : in out Object.Open_Handle;
                Value : Boolean;
                Status: out Error.Condition);
-- Defaults to off. When turned on, the references to Diana trees
-- (supplied by the operation below) will be expunge protected.
-- Possible errors:
    Is_Execution_Error (Status)
         o Not_A_Load_Image
         o Cant_Protect_Debug_Units
               This operation cannot be applied once a call to
               Set_Debug_Unit has been made.
function Debug_Unit (The_Load_Image : Object_Open_Handle;
                     Original_Index : Om.Object_Index)
              return Om. Version_Handle;
procedure Set_Debug_Unit (The_Load_Image: in out Object_Open_Handle;
                          Original_Index: Om.Object_Index;
                          Debug_Unit : Om. Version_Handle;
                          Status: out Error.Condition);
-- Map from original Object_Index to the Version_Handle of the
-- corresponding Diana tree. Debug tables use these object indexes to
-- refer to units. Use of Version_Handle allows tools to use the
-- universe component to record the identity of the view from which
-- the unit was loaded.
-- Possible errors:
-- Is_Execution_Error (Status)
         o Not_A_Load_Image
         o Debug_Unit_Index_Out_Of_Bounds
         o Not_A_Debug_Unit
___
               The referenced Debug_Unit does not refer to an
```

```
existing object of Class_Ada. Only occurs when
               Protect_Debug_Units is turned on.
-- The following operations only apply to secondary load images --
function Protecting_Primary
              (The_Load_Image : Object.Open_Handle)
              return Om. Version_Handle;
procedure Set_Protecting_Primary
               (The_Load_Image : in out Object.Open_Handle;
                Protecting_Primary : Om. Version_Handle:
                Status : out Error.Condition);
-- Protecting_Primary refers to a primary load image which will
-- eventually reference this secondary. Recall that such a reference
-- protects the secondary from expunge in the secondary's world.
-- Possible errors:
     Is_Execution_Error (Status)
          o Not_A_Load_Image
          o Not_A_Secondary_Load_Image
          o Primary_Load_Image_Does_Not_Exist
-- The following operations only apply to primary load images --
function Secondary_Count (The_Load_Image : Object.Open_Handle)
              return Natural;
function Secondary_Reference
              (The_Load_Image : Object.Open_Handle;
               Key: Secondary_Number)
              return Om. Version_Handle;
procedure Set_Secondary_Reference
               (The_Load_Image: in out Object.Open_Handle;
                Key: Secondary_Number:
                Secondary: Om. Version_Handle;
                Status: out Error.Condition);
-- Used to record, in a primary, the Version_Id of one of its
-- secondaries.
-- Possible errors:
     Is_Execution_Error (Status)
--
--
          o Rot_A_Load_Image
_ _
          o Not_A_Primary_Load_Image
          o Secondary_Load_Image_Does_Not_Exist
function Starting_Pc
               (The_Load_Image : Object.Open_Handle)
              return *achine.Address:
procedure Set_Starting_Pc
```

```
(The_Load_Image : in out Object.Open_Handle;
                Starting_Pc : Machine.Address;
                Status : out Error.Condition);
-- Used to remember the starting pc.
-- Possible errors:
--
     Is_Execution_Error (Status)
          o Not_A_Load_Image
          o Not_A_Primary_Load_Image
          o No_Code_Segment_For_Pc
function Is_Subsystem
              (The_Load_Image : Object.Open_Handle) return Boolean;
procedure Set_Is_Subsystem
               (The_Load_Image: in out Object.Open_Handle;
                Value : Boolean;
                Status: out Error Condition);
-- Defaults to off. Determines whether you get subsystem or main
-- program behaviour during elaboration.
-- Possible errors:
     Is_Execution_Error (Status)
          o Not_A_Load_Image
          o Not_A_Primary_Load_Image
function Is_Persistent
              (The_Load_Image: Object.Open_Handle) return boolean;
procedure Set_Is_Persistent
               (The_Load_Image : in out Object.Open_Handle;
                Value : Boolean;
                Status : out Error.Condition);
-- Defaults to off. The elaboration of a persistent subsystem will
-- still exist after the command which caused its elaboration.
-- Possible errors:
     Is_Execution_Error (Status)
          o Not_A_Load_Image
          o Not_A_Primary_Load_Image
          o Not_A_Subsystem_Load_Image
function Exported_Unit_Count
              (The_Load_Image : Object.Open_Handle) return Natural;
function Exported_Module_Count
              (The_Load_Image : Object.Open_Handle) return Natural;
function Exported_Unit_Id
              (Tne_Load_Image : Object.Open_Handle;
               Unit : Um.Object_Index)
              return Om.Object_Id;
function Exported_Unit_Name
              (The_Load_Image : Object.Open_Handle;
               Jnit : Om.Object_Index)
              return String;
function Exported_Module_Pos
              (The_Loao_Image : Object.Open_Handle;
```

```
Unit : Om.Object_Index)
              return Exported_Module_Position;
function Exported_Decls
              (The_Load_Image : Object.Open_Handle:
               Unit : Um.Object_Index)
              return Compatability_Vector;
procedure Set_Exported_Unit_Information
               (The_Load_Image : in out Object.Open_Handle;
                Unit_Id : Om.Object_Id;
                Unit_Name : String:
                Module_Pos : Exported_Module_Position;
                Exported_Decls : Compatability_Vector;
                Status: out Error.Condition);
-- The exported unit information is used (for a loaded subsystem) to
-- describe the actual declarations that are exported (via the
-- Compatability_Vector for each exported unit). This information is
-- used for compatability checking during elaboration.
-- Under library unit package integration, more than one exported unit
-- may specify the same Exported_Module_Position. During elaboration.
-- the module will call Operations.Set_Exported_Module with the Module
-- parameter equal to the value of Module_Pos supplied here.
-- Possible errors:
     Is_Execution_Error (Status)
          o Not_A_Load_Image
          o Not_A_Primary_Load_Image
          o Not_A_Subsystem_Load_Image
          o Exported_Module_Pos_Out_Of_Bounds
          o Exported_Unit_In_Wrong_World
                The Unit_Id of an exported unit must specify the
                same world_Number as that of the load image being
                constructed.
          o Load_Image_Is_Too_Big
function Imported_Subsystem_Count
              (The_Load_Image : Object_Open_Handle) return Natural;
function Imported_Subsystem_Id
              (The_Load_Image : Object.Open_Handle;
               Subsystem: Imported_Subsystem_Number)
              return Root_Of_World;
function Imported_Subsystem_Name
              (The_Load_Image : Object.Open_Handle;
               Subsystem: Imported_Subsystem_Number)
              return String;
function Imported_Subsystem_First_Unit
              (The_Load_Image : Object.Open_Handle;
               Subsystem : Imported_Subsystem_Number)
              return Imported_Unit_Number:
function Imported_Subsystem_Last_Unit
              (The_Load_Image : Object.Open_Handle;
               Subsystem : Imported_Subsystem_Number)
              return Imported_Unit_Number;
```

```
procedure Set_Imported_Subsystem_Information
               (The_Load_Image : in out Object.Open_Handle;
                Subsystem : Imported_Subsystem_Number:
                Subsystem_Id : Some_Object_In_World:
                bubsystem_Name : String;
                First_Imported_Unit : Imported_Unit_Number;
                Last_Imported_Unit : Imported_Unit_Number;
                Last_Imported_Module : Imported_Unit_Number;
                Status: out Error.Condition);
-- Possible errors:
     Is_Execution_Error (Status)
          o Not_A_Load_Image
          o Not_A_Primary_Load_Image
          o Not_A_Subsystem_Load_Image
          o Imported_Subsystem_Number_Out_Of_Bounds
          o Imported_Subsystem_Does_Not_Exist
          o lmported_Unit_Number_Out_Of_bounds
          o Load_Image_Is_Too_Big
function Imported_Unit_Count
              (The_Load_Image : Object.Open_Handle) return Natural;
function Imported_Unit_Id
              (The_Load_Image : Object.Open_Handle;
               Unit : Imported_Unit_Number)
              return Om.Object_Handle;
function Imported_Unit_Name
              (The_Load_Image : Object.Open_Handle;
               Unit : Imported_Unit_Number)
              return String;
function Imported_Module_Pos
              (The_Load_Image : Object.Open_Handle;
               Unit : Imported_Unit_Number)
              return Imported_Module_Position;
function Imported_Decls
              (The_Load_Image : Object.Open_Handle;
               Unit : Imported_Unit_Number)
              return Compatability_Vector;
procedure Set_Imported_Unit_Information
               (The_Load_Image : in out Object.Open_Handle;
                Unit : Imported_Unit_Number:
                Unit_Id : Om.Object_Id:
                Unit_Name : String;
                Module_Pos : Imported_Module_Position;
                Imported_Decls : Compatability_Vector;
                Status : out Error.Condition);
-- Poth subsystems and main programs can import the facilities of
-- other subsystems. Since environment facilities are imported using
-- this mechanism, it is unlikely that any load image has 0
-- subsystem imports. For each imported subsystem, there is a set of
-- imported units (as indicated by the First_Imported_Unit and
-- Last_Imported_Unit values). For each imported unit, one specifies
-- the declarations (via the Compatability_Vector) which are actually
-- imported. This information is matched against the unit export
-- information (in the load image of the corresponding imported
```

```
-- subsystem) during compatability checking during elaboration.
-- During elaboration the names of the imported runtime modules appear
-- in the import space by Imported_Module_Position; this value may be
-- different from the Exported_Module_Position of the corresponding
-- unit.
-- Possible errors:
     Is_Execution_Error (Status)
--
          o Not_A_Load_Image
          o Not_A_Primary_Load_Image
          o Not_A_Subsystem_Load_Image
          o Imported_Unit_Number_Out_Of_bounds
          o Imported_Module_Pos_Out_Of_Bounds
          o Imported_Unit_Does_Not_Exist
          o Load_Image_Is_Too_Big
function Needs_Ucode_Assist (The_Load_Image : Object_Open_Handle)
              return Boolean;
procedure Set_Needs_Ucode_Assist
               (The_Load_Image : in out Object.Open_Handle;
                Value : Boolean;
                Status : out Error Condition);
-- Defaults to off. This is a special kludge which is used to cause a
-- subsystem to import the ucode assist package. It is required by the
-- "system call and command subsystems". A subsystem with this bit set
-- can only be elaporated by Privileged users.
-- Possible errors:
--
     Is_Execution_Error (Status)
          o Not_A_Load_Image
___
          o Not_A_Primary_Load_Image
          o Not_A_Subsystem_Load_Image
subtype Load_Image_Id is Om.Version_Handle;
function Imported_Program_Count
              (The_Load_Image : Object.Open_Handle) return Natural;
function Imported_Program_Id
              (The_Load_Image : Object.Open_Handle;
               Program : Imported_Program_Number)
              return Load_Image_Id;
function Imported_Program_Name
              (The Load Image: Object.Open_Handle:
               Program : Imported_Program_Number)
              return String;
procedure Set_Imported_Program_Information
                (The_Load_Image : in out Object.Open_Handle;
                Program : Imported_Program_Number;
                Program_Id : Load_Image_Id;
                Program_Name : String;
                Status: out Error Condition);
```

-- Both subsystems and main programs can call other main programs.

```
11
 -- The identity of these programs is recorded as shown above.
-- Possible errors:
      Is_Execution_Error (Status)
           o Not_A_Load_Image
           o Not_A_Primary_Load_Image
--
           o Imported_Program_Number_Out_Of_Bounds
           o Imported_Program_Does_Not_Exist
           o Load_Image_Is_Too_Big
function Elaboration_Order (The_Load_Image : Object.Open_Handle)
              return Elaboration_Order_Array;
procedure Set_Elaboration_Order
               (The_Load_Image : in out Object.Open_Handle;
                Order : Elaboration_Order_Array:
                Status: out Error Condition);
-- The elaporation order is stored in the load image so that the
-- subsystem integrator can produce an identical elaboration order
-- in the target integrated subsystem.
-- Possible errors:
     Is_Execution_Error (Status)
          o Not_A_Load_Image
          o Not_A_Primary_Load_Image
          o Elaboration_Order_Length_Mismatch
                The *length of the Order parameter must be the same
                as the Unit_Count parameter to Define_Shape.
function Enable_Code_Acl_Check
              (Ine_Load_Image : Object.Open_Handle) return Boolean;
procedure Set_Enable_Code_Acl_Check
               (The_Load_Image : in out Object.Open_Handle;
                Value : Boolean;
                Status: out Error.Condition);
-- Defaults to off. When enabled, a job must have Execute access to
-- the load image in order to elaborate a new copy.
-- Possible errors:
     Is_Execution_Error (Status)
          o Not_A_Load_Image
          O Not_A_Primary_Load_Image
function Enable_Module_Acl_Check
              (The_Load_Image : Object.Open_Handle) return boolean;
procedure Set_chable_Module_Acl_Check
               (The_Load_Image : in out Object.Open_Handle;
                value : Boolean;
                status : out Error.Condition);
-- Defaults to off. when enabled, a job must have Execute access to
-- the load image in order to import modules from an already elaborated
-- persistent copy of the subsystem.
```

-- Execution Operations --

```
-- Possible errors:
   --
        Is_Execution_Error (Status)
             o Not_A_Load_Image
             o Not_A_Primary_Load_Image
   function Original_Machine_Name
                 (The_Load_Image: Object.Open_Handle) return String;
   procedure Set_Original_Machine_Name
                  (The_Load_Image: in out Object.Open_Handle;
                   Name : String);
   -- The originating machine name is useful for error messages and
    -- debugging.
    -- Possible errors:
         Is_Execution_Error (Status)
              o Not_A_Load_Image
              o Not_A_Primary_Load_Image
              o Load_Image_Is_Too_Big
   function Loader_version
                  (The_Load_Image: Object.Open_Handle) return String;
   procedure Set_Loader_Version
                  (The_Load_Image: in out Object.Open_Handle:
                   Version : String);
    -- The name and/or version and/or timestamp of the "loader" is useful
    -- for error messages and debugging.
    -- Possible errors:
         Is Execution Error (Status)
    --
              o Not_A_Load_Image
              o Not_A_Primary_Load_Image
              o Load_Image_Is_Too_Big
   function Program_Version
                  (The_Load_Image : Object_Open_Handle) return String;
   procedure Set_Program_Version
                   (The_Load_Image: in out Object.Open_Handle;
                    Version: String);
    -- The program/subsystem might find this useful for its own
    -- diagnostic purposes. Can contain program name, and/or version,
    -- and/or timestamp.
    -- Possible errors:
          Is_Execution_Error (Status)
               o Not_A_Load_Image
     __
               o Not_A_Primary_Load_Image
               o Loau_Image_Is_Too_Big
     __
end Load_Image:
--pragma Integrate (Load_Image);
```

package uperations is

```
-- Operations to control elaboration state --
-- The execution stack follows the same conventions as the various
-- job stacks, with some exceptions: (a) certain kinds of stack
-- entries cannot be copied (namely, direct entries); (b) it is
-- not possible to modify stack entries below top of stack.
-- Each stack element contains an entity which comes in 2 classes,
-- known as "direct" and "indirect" execution states. A indirect is
-- simply a reference to a direct. This capability can be used to -- acheive "shared elaboration". A direct execution state is a map
-- from persistent subsystem to physically elaborated copy of the
-- persistent subsystem.
-- When we say "physically elaborate" we mean that the elaborating
-- job causes the execution of a new runtime copy of the subsystem.
-- Inis causes the top (direct) execution state of the job to get a
-- map entry which selects the new copy; this kind of map entry is
-- known as a local entry.
-- It is also possible to cause a direct execution state to contain
-- a map entry which selects a copy which is physically elaborated
-- by another job. This capability can be used to acheive "shared
-- elaboration". This kind of map entry is known as a non-local entry.
-- When a direct execution state contains a map entry for a subsystem
-- S, the subsystem S is said to be "elaborated" in the given
-- execution state (regardless of whether or not the entry is local).
-- Each direct execution state is "consistent". In particular, the
-- following invariants hold (within the context of the particular
-- execution state):
--
       (1) Consider subsystem A which is elaborated. Consider a
           subsystem & which is with'd by A.
               (a) X is also elaborated.
__
               (b) The map entry for X references (directly or
                   indirectly) the physically elaborated copy of X
                   which is imported by A.
               (c) The import Compatability_Vector's of A "match" the
--
                   corresponding export Compatability_Vector's of the
--
                   physically elaborated copy of subsystem X which is
                   imported by A.
--
       (2) By induction, one can infer that if 2 elaborated subsystems
--
           A and B with X, they import from the same physically
           elaporated copy of X.
-- Indirect execution state are by definition consistent, since they
```

-- simply reference a consistent direct execution state.

-- "Sharing" is acheived as follows: Job X elaborates some subsystems,

-- and invokes a Make\_Execution\_State\_Shareable operation on the -- particular execution state. Job X may optionally provide a string -- name to identify the execution state. Job Y (the "importer") can

```
-- push an indirect execution state by invoking the
  -- Share_Execution_State operation, supplying the string name, or the
  -- explicit job and stack location. At this point, Y can run programs,
  -- having them import from the shared subsystems (which Y imported
  -- from X). If Y wishes to physically elaborate its own subsystems, it
  -- can Push_Direct_State. Elaboration operations in job Y can then
  -- create local map entries for Y's private subsystems, and can
  -- create non-local map entries from the indirect execution state (at
  -- Top-1). This may become clearer after reading further.
  -- In order for persistently elaborated subsystems to last longer than
  -- the command which caused their elaboration, we introduce the
  -- concept of a "subsystem state" job in which persistently elaborated
  -- subsystems live. There is at most one of these jobs per session.
-- We will use the term "elaborated" to mean that an elaborated copy
  -- of the subsystem is selected by the execution state on top of the
  -- stack. The copy might be selected via an indirect execution state,
  -- or via a direct state with either a local or non-local map entry.
  -- Some consequences of this definition: (a) Pushing an empty direct
  -- state results in no subsystem being "elaborated". (b) A subsystem
  -- can be said to be "not elaborated", even though lower stack
  -- elements reference physically elaborated copies. (c) Thus, the
  -- term "elaborated" is considerably different than the term
  -- "physically elaborated".
  type Elaboration_Options is new Long_Integer;
  Create_Job_On_Demand : constant Elaboration_Options := 2 ** 0;
  Push_Direct_State_On_Demand : constant Elaboration_Options := 2 ** 1;
  Force_Private_Copy : constant Elaboration_Options := 2 ** 2;
Allow_Private_Copy : constant Elaboration_Options := 2 ** 3;
  Elaborate_Closure_On_Demand : constant Elaboration_Options := 2 ** 4;
  Ignore_Compatability_Errors : constant Elaboration_Options := 2 ** 5;
                               : constant Elaboration_Options :=
  Default_Elab_Options
                               Create_Job_On_Demand +
                               Push_Direct_State_On_Demand +
                               Allow_Private_Copy +
                               Elaborate_Closure_On_Demand;
   generic
       with procedure Visit (Subsystem : Root_Of_World;
                             Name: String:
                             Elaborated_Release : Om.Version_Handle;
                             Physically_Elaborated : Boolean;
                             Where_Elaborated : Om.Job_Id);
   procedure Elaborate_Persistent_Subsystem
        (Subsystem : Some_Object_In_World;
         Status : out Error.Condition;
         Lxception_Info : out Exception_Information;
         Options : Elaboration_Options := Default_Elab_Options;
         In_Job : Om.Job_Id := Job.Subsystem_State);
```

-- The purpose of this operation is to elaborate the Subsystem, if

-- its not already elaborated.

- -- If In\_Job is Nil, and Create\_Job\_On\_Demand is enabled, the -- operation will create a Subsystem\_State job for the session of the -- calling job.
- -- If Push\_Direct\_State\_On\_Demand is enabled, the operation will do -- an implicit call to Push\_Direct\_State, if necessary.
- -- If Force\_Private\_Copy is enabled, then this operation will
  -- physically elaborate a new copy of every subsystem (in the given
  -- Subsystem's with closure, as well as Subsystem itself), in the
  -- identified subsystem job, for every subsystem which is not
  -- already elaborated.
- -- If Allow\_Private\_Copy is not enabled, then this operation is limited -- to creating map entries to already physically elaborated copies of -- subsystems.
- -- If ilaborate\_Closure\_On\_Demand is enabled, then this operation -- will, if necessary, attempt to elaborate "lower level" subsystems -- (ie, subsystems in the "with" closure).
- -- This operation applies to the execution state stack of the -- identified job.
- -- In the course of elaborating the needed subsystems, consider the -- situation in which the with closure of persistent subsystem S is -- elaborated, but S itself is not yet elaborated (and is the current -- goal). Note that if stack top is not direct and -- Push\_Direct\_State\_On\_Demand is enabled, an implicit call to -- Push\_Direct\_State will be made. The following actions are taken to -- elaborate S:
- -- If Force\_Private\_Copy is disabled, then look for a shareable copy
  -- by scanning the execution stack stopping at the first element which
  -- contains a map entry for S. Note that when this scan encounters
  -- indirect execution states, it simply looks at the referenced
  -- execution state. If a map entry is found, and the referenced
  -- physically elaborated copy of S imports from the same physically
  -- elaborated copies that are currently selected by the top execution
  -- state, then simply create a non-local map entry (in stack top)
  -- referencing the located copy. Calls the supplied Visit procedure
  -- prior to setting the map entry; the call would be made with
  -- Physically\_Elaborated => False.
- -- If Force\_Private\_Copy is enabled, or the search for a shareable copy failed, then attempt to physically elaborate a new copy of subsystem S, importing from the copies elaborated in stack top, creating a local map entry in stack top. This step may fail with creating a local map entry in stack top. This step may fail with elaboration; the call supplied visit procedure prior to physical elaboration; the call would be made with Physically\_claborated => True. Note that enabling potn Force\_Private\_Copy and Include\_Closure can result in the physical elaboration of many subsystems.
- -- when physically elaborating a new copy, the load image is -- identified by the world view specified by the job's view stack.

```
-- when physically elaborating a new copy, the runtime modules
-- execute in the identified job.
-- The caller suspends until the physically elaborated copies have
-- finished elaboration of their library units. Clearly, the runtime
-- state of these units survives the completion of this operation.
-- In the event of an error, the offending subsystem is unelaborated,
-- but the other subsystems that were elaborated are left elaborated.
-- There are notes at the end of this package describing the context
-- in which the code (in the load image) begins execution.
-- Note that since Visit is called prior to elaboration, many of the
-- error conditions are known to apply to the subsystem last Visit'd.
-- The Name parameter to Visit comes from the Program_Version attribute
-- of the elaborated load image.
-- Possible errors:
     Is_Job_Error (Status)
__
          o Job_Does_Not_Exist
              Non-Nil In_Job parameter refers to non-existant job.
     Is_Execution_Error (Status)
          o Missing_Subsystem_Job
              Can only happen when Create_Job_On_Demand is disabled.
__
___
          o Not_A_Subsystem_Job
              The identified job must be the subsystem job for
- -
              some session.
___
          o Execution_Stack_Is_Empty
___
              Can only happen when Push_Direct_State_On_Demand
               is disabled.
___
          o Execution_State_Not_Direct
--
              Can only happen when Push_Direct_State_On_Demand
__
              is disabled.
          o Subsystem_Is_Already_Elaborated
--
              The identified subsystem is already elaporated. In some
__
               cases caller might consider this to be "success".
          o Private_Copy_Not_Allowed
               Couldn't find an import compatable atready physically
__
               elaborate copy of a subsystem, and Allow_Private_Copy is
               disauled.
          o Cant_Find_Load_Image_To_Elaborate
__
               Either the job's universe stack does not select a
___
               world view for the subsystem, or the world view does
__
               not identify an existing load image.
___
          o Cycles_In_Subsystem_Closure
___
               Subsystem with dependencies induce a cylic order.
___
           o Cycles_in_View_Indirection
               Cycle detected determining the target universe view.
___
           o Subsystem_Not_Persistent
--
               The identified Subsystem is not persistent.
--
           o Persistent_Withs_Non_Persistent
_ _
               A subsystem (in the closure) withs a non-persistent
--
               subsystem.
           o Runtime_Compatability_Mismatch
               Subsystem A needs to be elaborated, and it with's
```

subsystem B. A imports units and/or declarations which

```
are not exported by B. This is detected by the fact
             that an import Compatability_Vector in the load image
              for A does not "match" the corresponding export
              Compatability_Vector in the load image for 3.
          o broken_Indirect_Execution_State
--
              The search for an elaborated copy ran across an
              indirect execution state (in the job's stack) which
              has been broken (perhaps by termination of the
              referenced job).
          o Subsystem_Propagated_Exception
              The elaboration of a library unit propagated an
              exception. Note that it is a consequence of Ada
__
              semantics that such exceptions cannot be propagated
              once this operation returns. The Exception_Info
__
              parameter contains additional information.
__
     Is_Access_Control_Error (Status)
--
          o Couldn't load code segments or import modules or import
            ucode assist due to an access control violation. See
            Load_Image packages for details. Or don't have required
__
            access to the identified job.
procedure Make_Execution_State_Shareable
               (Status : out Error.Condition;
                Name : String := "";
                In_Job : Om.Job_Id := Job.Subsystem_State);
-- The top stack location (of the given job) is marked as shareable.
-- If the Name parameter is not null, then the name attribute of the
-- top stack location is set to the given Name. If the resulting top
-- stack location ends up with a non-null Name attribute, invocations
-- of Shared_Execution_State will able to identify this state using
-- its string name. Names longer than about 128 characters are
-- truncated on the left.
-- Note that this job will not be able to terminate (or be Job.Kill'd)
-- if other job's are currently sharing this job's subsystems. Some of
-- the motivation for this invariant is "politeness"; but more
-- importantly, the machine might crash if the rule were violated.
-- Possible errors:
___
     Is_Job_Error (Status)
___
          o Job_Does_Not_Exist
     Is_txecution_Error (Status)
--
          o Not_A_Subsystem_Job
          o Execution_Stack_Is_Empty
          o Execution_State_Not_Direct
          o Execution_State_Is_Shareable
___
          o Execution_State_Name_Not_Unique
--
              when non-null, the string name of the execution state
              must be unique among all of the execution states (in the
              entire system) which are marked as snareable.
procedure Make_Execution_State_Not_Shareable
               (Status: out Error.Condition;
                In_Job : Om.Job_Id := Job.Subsystem_State);
```

-- The top stack location (of the identified job) is made not shareable.

```
-- Will break any referencing indirect execution states. Note that
-- this operation does NOT unelaborate anything. And it does NOT
-- break non-local map entries (of direct execution states) that have
-- been constructed by elaboration via previously unbroken indirect
-- execution states.
-- Possible errors:
     Is_Job_Error (Status)
          o Job_Does_Not_Exist
     Is_Execution_Error (Status)
          o Not_A_Subsystem_Job
          o Execution_Stack_Is_Empty
          o Execution_State_Not_Direct
          o Execution_State_Not_Shareable
procedure Share_Execution_State
               (Source_State : String;
                Status: out Error.Condition;
                Target_Stack : Job.Stack_Id := Job.Job_Stack);
procedure Share_Execution_State
               (Source_State : Om.Job_Id;
                Status: out Error.Condition;
                Source_Job : Om.Job_Id := Job.Subsystem_State;
                Source_Location : Job.Stack_Location := Job.Top:
                Target_Stack : Job.Stack_Id := Job.Job_Stack);
-- The first form identifies the source execution state by string
-- name, whereas the second form identifies the source execution state
-- by explicitly giving the job and stack location. Clearly, the
-- presence of concurrency may make the second form subject to
-- misinterpretation of stack position.
-- In either case, an indirect elaboration state is pushed on the
-- identified target stack.
-- Invocation of Make_Execution_State_Unshareable (on the referenced
-- elaboration state) will cause the pushed indirect elaboration state
-- to become broken and cause Elaborate_Persistent_Subsystem to
-- generate errors. This may happen automatically when a job
-- terminates.
-- Recall that this is what allows the Elaborate_Persistent_Subsystem
-- operation to elaborate non-local map entries and thereby allows one
-- job to physically elaborate subsystems against subsystems
-- physically elaborated by another job. And it allows the
-- Elaborate_Program operation to directly import from the subsystem
-- physically elaborated by another job.
-- Possible errors:
     Is_Job_Error (Status):
          o Job_Does_Not_Exist
              Applies to either source or target job.
--
     Is_Execution_Error (Status):
          o Execution_State_Does_Not_Exist
___
              The given string (or given job & stack location) do
--
              not identify a execution state.
          o Execution_Stack_Is_Empty
```

```
Applies to source job.
          o Execution_State_Not_Direct
--
              Applies to source job.
          o Execution_State_Not_Shareable
              Applies to source job. The identified execution state
              has not been the subject of
              Make_Execution_State_Shareable operation.
--
          o Execution_State_Stack_Overflow
              Applies to target job.
procedure Un_Elaborate_Persistent_Subsystem
               (Subsystem : Some_Object_In_World;
                Status: out Error.Condition;
                Un_Elaborate_Clients : Boolean := True;
                In_Job : Om.Job_Id := Job.Subsystem_State);
-- Used to unelaborate a persistent subsystem. When
-- Un_Elaborate_Clients is enabled, will unelaborates all the with ing
-- subsystems as well. Unelaborate means: Deselect the subsystem and
-- physically unelaborate those subsystems which this job physically
-- elaborated. Can use this to kill a partial elaboration. Can be
-- invoked by the given subsystem itself, or one whose closure
-- includes the given subsystem.
-- Possible errors:
     Is_Job_Error (Status):
___
          o Job_Does_Not_Exist
              Applies to either source or target job.
     Is_Execution_Error (Status)
          o Execution_Stack_Is_Empty
__
          o Execution_State_Is_Not_Direct
--
          o Execution_State_Is_Shareable
              The identified execution state has been made shareable
              via the Share_Execution_State operation.
--
          o Subsystem_ls_Not_Elaborated
          o Subsystem_Still_Has_Clients
              The identified subsystem has clients, but
              Un_Elaborate_Clients is not enabled. Note that these
              clients are in the same execution state.
procedure Elaborate_Program
               (Program_Load_Image : Om.Version_Id;
                Status: out Error.Condition;
                Exception_Info : out Exception_Information;
                Options : Elaboration_Options := Default_Elab_Options;
                In_Job : Om.Job_Id := Job.Subsystem_State);
-- The purpose of this operation is to execute (ie, elaborate) the
-- identified program.
-- If the Program_Load_Image references a persistent subsystem that is
-- not already elaborated and Elaborate_Closure_On_Demand is enabled,
-- then this operation will invoke Elaborate_Persistent_Subsystem
-- (with the given Options and Job_Id) to get the subsystem elaborated.
```

-- Elaborates copies (private to this program) of all the

-- non-persistent subsystems in the program's subsystem closure. Runs

```
-- the program. When the program returns, unelaborates the
 -- non-persistent copies, but not the persistent copies.
 -- The caller is suspended during the program's execution.
-- If the task executing the program is aborted, the non-persistent
-- copies are automatically unelaborated.
 -- Possible errors:
 ___
      Is_Execution_Error (Status)
 --
           o Subsystem_Is_Not_Elaborated
               A with'd persistent subsystem is not elaborated, and
               Elaborate_Closure_On_Demand is not enabled.
           o Plus all those that can be produced by
             Elaborate_Persistent_Subsystem
5 Start
 procedure Create_Subsystem_State_Job
                (Status : out Error.Condition;
                 For_Session : Om.Session_Id := Job.Session);
 -- Can be used to explicitly create the subsystem state job for a
 -- session. (an be done implicitly by the elaboration operations.
 -- Possible errors:
 ___
      Is_Job_Error (Status):
           o Session_Does_Not_Exist
      all errors produced by Job. Initiate.
 procedure Push_Direct_State (Status : out Error.Condition;
                               Name : String := "";
                               In_Job : Om.Job_Id := Job.Subsystem_State);
 -- Pushes a fresh elaboration context in which nothing is elaborated.
 -- Possible errors:
      Is_Job_Error (Status):
           o Job_Does_Not_Exist
      Is_Execution_Error (Status):
           o Execution_State_Stack_Overflow
 procedure Pop_State (Status : out Error.Condition;
                       Stack: Job.Stack_Id:= Job.Job_Stack;
                       Location : Job.Stack_Location := Job.Top);
 -- Will automatically invoke Un_Elaborate_Persistent_Subsystem, if
 -- necessary. Can be applied to shareable execution state, provided
 -- there are no current importers. This operation is used by job
 -- termination (including Job.Kill), and may therefore be responsable
 -- for the inability to kill a job.
 -- Possible errors:
 --
      Is_Job_Error (Status):
 --
           o Job_Does_Not_Exist
 --
      Is_Execution_Error (Status)
           o Execution_Stack_Is_Empty
            o Subsystem_Still_Referenced
```

The identified execution state is direct and contains

```
a local entry for a physically elaborated copy which
              is still referenced by the direct execution state of
              some other job.
-- Information requests --
   function Stack_Depth (Stack : Job.Stack_Id := Job.Job_Stack)
              return Natural;
function State_Exists (Stack : Job.Stack_Id := Job.Job_Stack;
                       Location : Job.Stack_Location := Job.Top)
              return Boolean;
function Is_Persistently_Elaborated
              (Subsystem : Some_Object_In_world;
               In_Job : Om.Job_Id := Job.Current) return Boolean;
-- Returns true iff the identified job selects an persistent elaborated
-- copy of the Subsystem.
function Executing_Release
              (Subsystem : Some_Object_In_World;
               In_Job : Om.Job_Id := Job.Current)
              return String;
-- Simply returns Nil if the identified job has no persistent copy
-- selected and its view stack does not select an executable release
-- of the subsystem. Else, returns the Program_Version attribute of
-- the load image of the job's selected persistent copy or of the load
-- image that would be elaborated (from the view stack).
generic
    with procedure Visit (Job: Om.Job_Id;
                          Load_Image : Om. Version_tandle;
                          Is_Subsystem : Boolean;
                          Name : String);
procedure Locate_Un_Elaborate_Preventing_References
     (Subsystem : Some_Object_In_World;
      ibI_dcL.m0 : doL_nI
      Status : out Error.Condition);
-- will simply produce an error condition if the identified subsystem
-- is not elaborated. Else calls Visit for every subsystem/program
-- which must terminate before the identified subsystem can be
-- unelaporated. The Name string comes from the Program_Version
-- attribute of the visited load image. In some cases, the only valid
-- information will be the Job_Id; in this case, the Load_Image will
-- be mit.
-- Possible errors:
    Is_Execution_Error (Status)
          o Subsystem_Is_Not_Elaborated
    with procedure Visit (Subsystem : Root_Of_world;
```

Name : String); procedure Locate\_Job\_Termination\_Preventing\_Subsystems (In\_Job : Om.Job\_Id; Status : out Error.Condition); -- Will simply produce an error condition if the identified job does -- not exist. Else calls Visit for every persistent subsystem -- physically elaborated by the identified job and currently shared -- with other jobs and which therefore prevent the given job from -- terminating. The Name string comes from the Program\_Version -- attribute of the corresponding load image. -- Possible errors: Is\_Execution\_Error (Status) o Subsystem\_Is\_Not\_Elaborated -- Operations used by the executing program/subsystem -type Elaboration\_Handle is private; procedure Set\_Exported\_Module (Subsystem : Elaboration\_Handle; Module : Exported\_Module\_Position; Status : out Error Condition); -- This operation is intended for wise by an elaborating persistent -- subsystem. See code generation notes below. -- This operation is a noop if the subsystem is not persistent. It -- is used to record the name of an exported runtime module, for -- later importation by some subsystem/program. -- Possible errors: Is\_Execution\_Error (Status) o Bogus\_Elaboration\_Handle o Exported\_Module\_Pos\_Out\_Of\_Bounds type Load\_Image\_Handle is private; -- dandle on the load image object of the running program. procedure Elaborate\_Program (Program\_Load\_Image : Load\_Image\_Handle; Program : Imported\_Program\_Number: Number\_Of\_Arguments: Natural: Returns\_Result : Boolean; Status : out Error.Condition); -- Logically equivalent to the above form. Intended for use only by -- "linkage" code generated by the code generator/loader. The call -- site is passing Number\_Of\_Arguments and expects to receive a

-- Possible errors:

-- notes below.

-- all those that can be produced by the previous form of the

-- function result if Returns\_Result is true. See code generation

-- operation.

```
Is_Execution_Error (Status)
          o Bogus_Load_Image_Handle
          o Program_Number_Out_Of_Bounds
-- Load image execution environment --
-- With respect to code generation, the following information is
-- important (if you're not a code generator writer, you probably
-- don't care).
-- The context for executing a subsystem load image is as follows: A
-- skin package (the root module) is elaborated whose import space
-- has things pushed in the following order:
      o Elaboration_Handle (argument to pass to following op)
      o Proc ref for Operations.Set_Exported_Module
      o Load_Image_Handle (argument to pass to following op)
--
      o Proc ref for Uperations. Elaborate_Program (2nd form)
__
      o Code segment names are pushed, in their "protected form", in
        the following order: for each load image (starting with the
--
        primary, and proceeding to the secondaries, in
--
        Secondary_Tumber order) push the code segment names in
        Code_Segment_Key order.
--
      o The imported package variables are pushed in the following
___
        order: First push ucode assist (or Nil if not requested).
--
        Then, for each imported subsystem (in
        Imported_Subsystem_Number order) push the package variable
        for the root module (or nil if the subsystem was not
___
        elaborated by this job) followed by the package variables
        for each imported module (in Imported_Module_Position order).
-- The declaration list of the skin package (root module) runs the
-- code specified by the load image's start pc. This mechanism
-- assumes that the generated code is such that the skin package will
-- complete its elaboration as soon as all the library units (in the
-- subsystem) have completed their elaboration.
-- Note that the skin package is dependent upon some system package,
-- and not the caller's frame. The module name of the skin package
-- (root module) is recorded for later use by Un_Elaborate_Subsystem.
-- The context for executing a program load image is as follows: A
-- subprogram variable is created (using the load image's start pc).
-- The subprogram is called with arguments pushed in the following
-- order:
--
      o "Ada" arguments (Number_Of_Arguments)
      o Load_Image_Handle
      o Proc ref for Execution_Elaborate_Program (2nd form)
      o Code segment names (in same order as for subsystems)
      o Imported package variables (in same order as for subsystems)
-- It is expected to consume all of its arguments, returning \tilde{\epsilon} or 1
-- results (as specified by the Returns_Results parameter). A
-- result returned from a program executed via the first form of
-- Elaborate_Program is simply ignored. For command execution, it
-- is expected that the job's root thread will be used to execute
-- the program.
```

## -- Debugger operations --

function To\_Version\_Handle (Program\_Load\_Image : Load\_Image\_Handle) return Om. Version\_Handle;

- -- Can be used by the debugger to convert the handle (in the -- "imports" for the subsystem or program) to a Version\_Handle.
  - -- Observations on Debugging --
- -- How does one associate an executing program with a debugger, and
- -- how does its debug state (namely, interface subprogram and break
- -- mask) get set properly? By default, the debug state is mil. If you
- -- tell the code generator the right thing, it will import the
- -- appropriate debugger, and generate code to call it to establish
- -- debugging; it calls the debugger using the normal importing rules
- -- (ie, not with fixed module names).
- -- Suppose one has elaborated a load image, and now one wants to
- -- debug it, which implies finding the biana trees associated with
- -- the various code segments in the load image. This is an
- -- excellent question. How does it work?
- -- Amplification --
- -- One can get amplification in shared subsystems by doing the
- -- initial elaboration in a job with the appropriate rights.
- -- Currently, there are no mechanisms for automatic amplification.

## private

type Load\_Image\_Handle is new Load\_Image\_Layout.State\_Record\_Pointer; type Elaboration\_Handle is new Execution\_Layout.Subsystem\_State\_Pointer;

-- Both types must be "by value" types so that we can stick them

-- in import spaces without creating dangling references.

end Operations; --pragma Integrate (Operations);

end Execution;

```
GGG
       pppp
                 AAA
G
   G
       P
          P
G
       Р
            Ρ
               A
G
   GG
       PPPP
               AAAAA
G
       P
    G
               A
                    A
G
    G
       P
               A
                    A
GGGG
       Р
                    Α
```

User: GPA

Object: !DELTA\_KK.REV1\_0\_0.UNITS.OM.OM\_SERVICES\_1

Version: V(18) Request: 1275

Date: April 24, 1986 Queued: 11:27:01 Am

Printed: 11:41:31 Am

1

```
with Job;
with Disk:
with Error;
with Naming;
with Object:
with Machine:
with Calendar;
with Om_Definitions;
package Om_Services_1 is
    package Link is
         -- Pragma Subsystem (Object_Management);
         -- A "link pack" is an object that defines a map from simple Ada
         -- names to Ada library units. A "link" is one element of this map.
         -- Each world has exactly one link pack object. Eech world view
         -- potentially has its own version of the link pack.
         -- A link is Internal if its designated object is in the world of the
         -- link pack; otherwise it is External.
         -- As defined by the Object package, there is an internal link for
         -- every Ada library unit; its link name is the same as the simple
         -- name of the designated object. External links are another story.
         -- They are maintained by higher level tools and allow the link name
         -- to be different than the simple name of the designated object.
         -- ALL of these procedures are subject to a bad status resulting form
         -- the inability to acquire appropriate lock(s) or insufficient disk
         -- space on the unit.
        package Um renames Om_Definitions:
        subtype Context is Om.Naming_Context;
        procedure Create (In_World : Om.Version_Handle;
                          Link_Name : Om.Simple_Name_String:
                          Designated_Name : Om.Unique_Wildcard;
                          Id : Om.Action_Id;
                          Status: out Error.Condition;
                          In_Context : Context := Job.Default_Context;
                          Max_Wait : Duration := Job.Default_Wait);
        procedure Create (In_World : Om.Version_Handle:
                          Link_Name : Om.Simple_Name_String;
                          Designated_Object : Om.Object_Handle;
                          Id : Om.Action_Id;
```

-- Usual rules for Naming.Resolve, and interpretation of the -- Version\_Handle/Object\_Handle, as specified by the Object package.

Status: out Error.Condition;

dax\_wait : Duration := Job.Default\_wait);

<sup>--</sup> Creates an external link.

ブ

```
-- Acquires a Record_write_Object lock on the link pack, and one or
-- more Record_Write lock's within the link pack, with the given action
-- The In_world parameter may refer to any object/version in the
-- world of interest.
-- Possible errors:
--
      Is_Link_Error (Status):
 __
          Duplicate_Link_Name
              o The given Link_Name would duplicate that of an already
                existing internal or external link.
 --
          External_Link_To_Local_Object
 --
              o The Designated_Object is in the same world as the
 --
                link pack.
          Cant_Create_External_Link
              o Insufficient internal data structures exist to create
                this link.
 --
      Is_Bad_Version_Handle (Status)
procedure Delete (In_World : Om.Version_Handle;
                  Link_Name : Om.Simple_Name_String;
                  Id: Om.Action_Id;
                  Status: out Error.Condition;
                  wlax_Wait : Duration := Job.Default_Wait);
-- Deletes an existing link.
-- The In_world parameter may refer to any object/version in the
-- world of interest.
-- Acquire a mecond_Write_Object lock on the link pack, and one or
-- more Record_drite lock's within the link pack, with the given
 -- action.
 -- Possible errors:
      Is_Link_Error (Status):
           Link_Name_Does_Not_Exist
           Cant_Delete_Internal_Link
      Is_Bad_Version_Handle (Status)
procedure Resolve (Link_Name : Om.Unique_Wildcard;
                   Designated_Object : out Om.Object_Id:
                   Status: out Error.Condition);
-- The Link_Name is actually a Simple_Wildcard which must resolve to
 -- just a single link. If Status is successful, then the given
 -- Link_Name maps to Designated_Object.
 -- Used by Naming-Aesolve when it encounters a """ character.
 -- Possible errors:
    Is_Pad_Pathname (Status)
ueneric
    with procedure Visit (Pesult : Om.Object_Id);
procedure wild_Resolve (Link_Name : Um.Simple_Wildcard;
```

## Status: out Error.Condition);

- -- Similar to the Resolve procedure. Is happy with 0 or more matches, -- calling Visit for each match. Stops searching after first error.
- -- Used by Naming.wild\_Resolve when it encounters a "'" character.

procedure Get\_Full\_Pathname (Designated\_Object : Om.Object\_Id; Full\_Name : out Om.Pathname; Status : out Error.Condition);

- -- Typically used when Resolve returns an Object\_Id which turns out
- -- to be a dangling reference, and one wants to display a message
- -- indicating what we thought the link was supposed to reference.
- -- Possible errors:
- -- Is\_Link\_Error(Status):
- -- Cant\_Identify\_Link

-- Use of the following operations is restricted to Object'body --

- -- Called by Object. Open when undifferentiating a link pack. The
- -- given segment is open with mode; Write\_Object. The operation
- -- copies the map which accelerates resolution of link names.

procedure Create\_Internal\_Link

(Link\_Pack : Om.Version\_Handle; Link\_Name : Om.Simple\_Name\_String; Designated\_Object : Gm.Object\_Id; Id : Om.Action\_Id; Status : out Error.Condition; Max\_wait : Duration := Job.Default\_Wait);

- -- Creates an internal link.
- -- Acquires a Record\_Write\_Object lock on the link pack, with the -- given action (the same action as given to Object\_Create).
- -- Link\_Name must be the simple name of Designated\_Object.
- -- Acquires the Version\_Map Muxtex during the deletion of the link;
- -- No Log record is written; assumes Client will call
- -- Delete\_Internal\_Object if abandoning the action.
- -- Possible errors:
- -- Is\_Link\_Error(Status):
- -- Duplicate\_Link\_Name

procedure Delete\_Internal\_Link

(Link\_Pack : Om.Version\_Handle;

Link\_Name : Om.Simple\_Name\_String;

1d : Om.Action\_Id;

Status : out Error.Condition;

```
Max_Wait : Duration := Job.Default_Wait);
     -- Deletes the internal link, with the given Link_Name.
     -- Acquires a Record_Write_Object lock on the link pack, with the
     -- given action (typically a private action of expunge).
     -- Acquires the Version_Map Muxtex during the addition of the link;
     -- No Log record is written; assumes Client will call
     -- Create_Internal_Object if abandoning the action.
     -- Possible errors:
     -- Is_Link_Error(Status):
               Link_Name_Does_Not_Exist
end Link:
package User is
   package Um renames Om_Definitions;
   -- Users are denoted by a directory (with subclass User_world) under
   -- "!Users". Various security related parameters are stored in
   -- "!Users._Profiles_". Passwords are stored in "!Users._Passwords_".
    -- Accounting information in stored in the "!Users.Accounting_Log".
    -- Sessions are simply directories with subclass Session, typically in
   -- the users home directory. Groups'are denoted by a directory (with
   -- subclass Group) under "!Groups"; uses a link pac to construct the
   -- permanent representation of the set of users/groups in the group
   -- (denoted by the directory).
    procedure Create (User_Name : Om.Pathname_String;
                      Password: String;
                      Home_Directory : out Om.Object_Handle;
                      Users_ld : out Om.User_ld;
                      Status: out Error.Condition);
    -- Creates a new world for the user, returned in Home_Directory.
    -- Initially, User_Name must be a simple name. At some future point in
    -- time, might allow User_Name to also specify intervening directories
    -- between "!Users" and the new user's home directory.
    -- Possible errors:
         Is_Access_Control_Error (Status)
    procedure Get_Id (Name : Om.Unique_Wildcard;
                      Users_Object_Id : out Om.Object_Handle:
                      Users_Id : out Om.User_Id;
                      Status: out Error_Condition);
    procedure Set_Id (name : Om.Object_Handle;
                      Users_Id: out Om.User_Id;
                      Status: out Frror Condition);
```

-- kith the first procedure, User\_Name is resolved to an Object\_Handle

```
-- via Naming.Resolve. A resolve error implies the user does not
-- exist. Then, with either procedure, computes the corresponding
-- User_Id.
-- Possible errors:
--
     Includes all those from Naming.Resolve
     Is_User_Error (Status)
         o Not_A_User_World
               o The Object_Handle does not identify a user world.
function Simple_Name (User : Om.Object_Handle)
              return Om.Simple_Name_String;
-- Assuming User actually identifies an existing user, returns the
-- user's simple name, as defined by Naming.Get_Simple_Name.
function Home_Directory (User: Om.Object_Handle)
              return Om.Pathname_String;
-- Assuming User actually identifies an existing user, returns the
-- full pathname of the user, as defined by Naming.Get_Full_Name.
procedure Destroy (Jser : Om.Object_Handle;
                   Status : out Error.Condition;
                   Home_Directory_Must_Be_Empty : Boolean := False);
-- The identified user will no longer be able to login. The user's
-- home world is destroyed (in the sense defined by World Destroy,
-- supplying not Home_Directory_Must_Be_Empty for the
-- Destroy_Contained_Objects and Destroy_Contained_Worlds parameters).
-- Possible errors:
     All those returned by World. Destroy
function Last_Login_Time (User : Om.Object_Handle) return Calendar.Time;
function Last_Logout_Time
              (User : Om. Object_Handle) return Calendar.Time;
procedure Set_Password
               (User : Om. User_Id;
                Old_Password : String;
                New_Password : String;
                Status: out Error.Condition);
-- Possible errors:
     Is_Access_Control_Error (Status)
procedure Validate_Password
               (User : Om. User_Id:
                password: String;
                Status : out Error.Condition);
-- Possible errors:
   Is_User_[rror (Status)
         o Invalid_Password
```

```
page
            procedure Create (Name : Om.Pathname_String;
                              New_Id : Om.Group_Id;
                              Action : Om.Action_Id;
                              Status : out Error.Condition);
            procedure Get_Id (Name : Om.Unique_Wildcard;
                              Its_Id : out Om.Group_Id;
                              Status : out Error.Condition);
            procedure Delete (Group : Om.Group_Id;
                              Action : Om.Action_Id;
                              Status : out Error.Condition);
            package Member is
                procedure Add (Group : Om. Group_Id;
                               User : Om. User_Id;
                               Action : Om.Action_Id;
                               Status : out Error.Condition);
                procedure Remove (Group : Om.Group_Id:
                                  User : Om. User_Id;
                                  Action: Om.Action Id;
                                  Status : out Error.Condition);
            ena Member:
        ena Group;
    ena User:
   package View is
        package um renames um_Definitions;
        -- We assume reader is already familiar with the basic meaning and
        -- function universe and world views.
        -- A universe view is thought of as an array of elements, indexed by
        -- world number. Each element is called a "slot".
        -- There are 3 values stored in a universe view slot:
        --
                Bound_World_View : Om.Version_Id;
                Bound_Wv_Is_Protected : Boolean;
                Indirect_World_View : Om.Object_Id;
        -- and they have the following meaning:
        -- If Bound_world_View and Indirect_World_View are both Nil, then this
        -- slot is said to be Nil, and does not reference any world view.
        -- If bound_World_View is not Nil, this slot is said to be "bound" and
        -- to reference the identified world view version. Note that the world
        -- number (in the version id) is always the same as the index (of the sl
ot).
        -- If bound_world_View is Nil, this slot is said to be "indirect". The
        -- Indirect_worls_view references the corresponding element in the
        -- identified universe view. In this case, one "resolves" an object id
        -- by following these indirect links until a bound slot is
```

-- encountered. Cycles in the indirection are detected by having a

-- maximum depth (of 31, say).

```
page
-- If both Bound_World_View and Indirect_World_View are non Nile the
-- slot is said to be "refreshable", meaning that Bound_World_View can
-- be updated to reference to whatever Indirect_World_View resolves.
-- Bound_Wv_Is_Protected can only be true when Bound_World_View is not
-- Nil. If Bound_wv_Is_Protected is true, then the bound world view is
-- protected from being expunded, as long as the universe view still
-- exists.
procedure Create (Target_Universe_View : Om.Unique_Wildcard;
                  Action: Om.Action_Id;
                  New_Universe_View : out Om.Version_Handle;
                  New_World_View : out Om.Version_Handle;
                  Status: out Error.Condition;
                  Crippled: Boolean := False;
                  Indirect : Om. Version_Handle := Om. Nil_Version_Handle;
                  Max_Ada_Objects : Integer := - 1;
                  hax_Other_Objects :: Integer := - 1;
                  Version_Retention : Om.Version_Retention_Count := 1;
                  %ax_Wait : Duration := Job.Default_Wait);
-- This operation is typically used to initialize a new world, and
-- when creating a new path from scratch.
-- Max_Ada_Objects is interpreted only when the Indirect parameter is
-- nil; similarly for Max_Other_Objects. In this case, if
-- Max_Ada_Opjects is negative, then it defaults to the value returned
-- by World.Max_Ada_Objects (for the world in which the world view
-- version is created). Similarly for Max_Other_Objects.
-- Creates a new universe view, initialized as follows: it
      o is untrozen.
      o all slots are nil, except the slot for the world in which the
         universe view is created (note that a crippled universe view
___
         has just this one slot).
-- If the Indirect parameter is not nil, then the slot is set to be
-- and indirect reference through the specified universe view. Note
-- that no check is made to verify that the identified universe view
-- actually exists. If the Indirect parameter is nil, then this slot
-- references a new version of the world view, which is initialized as
-- follows: it
```

- -o is unfrozen
- \_\_\_\_ o is not protected
- o can support the specified number of Ada and other objects
- o each slot can reference the specified number of versions
- o all slots are nil except the slot for the link pack, which -references a new version of the link pack, which is empty. \_\_
- -o all slots have the Peserved bit set
- o all unit states are nil \_\_\_
- \_\_\_ o the dependency matrix is clear.
- -- Note that the new universe view is NOT entered into any world -- views.
- -- If Target\_Universe\_View is "", the simple name of the new universe -- view is computed using the same rules as for creating temporary

```
!DELTA_KK.REV1_0_0.UNITS.OM.OM_SERVICES_1.V(18)
```

e;

```
-- object names in Object.Create (with ""); the universe view will be
-- a child of the world in which it is created.
-- The new universe and world view are both held with write_Object
-- locks. But note that the Object operations and the universe view
-- slot operations (below) still work on the view (even though upgrade
-- from w0 to RWO is normally not supported).
-- Destroy of a Universe_View/World_View can be acheived via
-- Ooject Expunge.
-- Possible errors:
--
     Is_Bad_Pathname (Status)
     Is_View_Error (Status)
--
--
         o Cant_Recreate_View
               The object named by Target_Universe_View already exists.
--
         o Bad_Object_Index_Constraint
___
               The value of Max_Ada_Objects or Max_Other_Objects is not
               within the bounds prescribed by the parameters of the
--
               containing world.
--
     Is_Bad_Action (Status)
--
     Is_Lock_Error (Status)
     Is_Resource_Limit_Error (Status)
type Freeze_Option is new Long_Integer range 0...7;
Same_Frozenness_4s_Source : constant Freeze_Option := 1;
Freeze_Target
                  : constant Freeze_Uption := 2;
Unfreeze_Target
                         : constant Freeze_Option := 3;
Default_Res_Bits
                  : Om.Reservation_Bits (1..0);
procedure Copy (Source_Universe_View : Om.Version_Handle;
                Target_Universe_View : Om.Unique_Wildcard;
                Action: Om.Action_Id;
                Status : out Error.Condition;
                Version_Prune_Target : Boolean := False;
                This_World : Freeze_Option := Same_Frozenness_As_Source;
                Other_Worlds : Freeze_Option := Same_Frozenness_As_Sourc
                Ada_Res_Bits : Om.Reservation_Bits := Default_Res_Bits;
                Other_Res_Fits : Om.Reservation_Bits := Default_Res_Fits
                Max_Ada_Objects : Integer := - 1;
                Max_Other_Gbjects : Integer := - 1;
                Max_wait : Duration := Job.Default_Wait);
-- This operation can be used to implement higher level operations
-- such as fork a path, spawn a suppath, make a release, vanilla view
-- copy, etc.
-- The target universe view is created, and initialized as follows:
-- its frozen*ness is as specified by This_World Freeze_Option.
-- parameter. It has the same Cripplediness as the source. For empty
-- slots in the source universe view, the target universe view has
```

-- empty slots. For non-empty slots in the source, the corresponding

- -- target slot is initialized as follows:
- -- If the source slot is indirect, the target slot becomes identical
- -- to the source slot. But note that this is no guarantee that the
- -- indirection actually leads to someplace interesting!
- -- If the source slot is bound (the only remaining case), the
- -- following rules apply:
- -- The target slot will reference the same world view (as the source)
- -- only if the source world view is frozen and not protected by the
- -- source universe view.
- -- In all other cases, a new world view version is created, which is -- copy of the source world view, except as indicated by the following:
- -- If the target world view is in the same world as the target
- -- universe view, then the This\_World parameter applies, else the
- -- Other\_Worlds parameter applies. The target world view will be
- -- frozen iff
- (a) The Freeze\_Option is Same\_Frozenness\_As\_Source, and the source world is frozen, or
- (2) The Freeze\_Option is Freeze\_Target.
- -- The target world view will be version pruned if the target world
- -- view is to be frozen, or the Version\_Prune\_Target parameter is
- -- true. By "version pruned" we mean that the retention lists in the
- -- new world view will be of length: 1 (may expunge deleted versions).
- -- The target world view will be object pruned if the target world
- -- view is to be frozen. By "object pruned" we mean that the maximum
- -- number of Ada units and other objects will be reduced to cover just
- -- the objects that are actually selected by the world view (makes the
- -- view smaller).
- -- The target world view will be protected by the target universe view
- -- if the source world view is protected by the source universe view.
- -- If Max\_Ada\_Objects is negative, then it defaults to the value
- -- returned by World.Max\_Ada\_Objects (for the world in which the
- -- target world view version is created). Similarly for
- -- Max\_Other\_Objects.
- -- If the target world view is unfrozen, then the following apply: (1)
- -- The target view can support the number of Ada and other objects
- -- specified by the Max\_Ada\_Objects and Max\_Other\_Objects parameters.
- -- (2) If both Ada\_Res\_Bits and Other\_Res\_Bits are null arrays, then
- -- the reserved pits in the target world view are identical to those
- -- of the source world view. Otherwise, the reserved bits (in the
- -- target world view) are set as specified by the value of the
- -- Ada\_Res\_Bits and Other\_Res\_Bits parameters; objects not specified
- -- have their reserved bits True; objects which appear in both the
- -- Aga\_Res\_Bits and Other\_Res\_Bits get an unspecified value for their
- -- reserved bit; the operation ignores slices of Ada\_Res\_bits or
- -- Other\_Res\_Fits which are out of the range of Object\_Indexes defined
- -- for the view.
- -- The new universe and world views are held with write\_úbject locks.

-- object lock).

```
-- But note that the Object operations and the universe view slot
-- operations (below) still work on the views (even though ubgrade
-- from w0 to RWO is normally not supported).
-- Note that link packs are differentiated on demand, just like
-- versions of other objects; therefore no special rules are required.
-- Possible errors:
--
     Is_Bad_Version_Handle (Status)
--
         o The Source_Universe_View does not exist
     Is_Bad_Pathname (Status)
     Is_View_Error (Status)
--
         o Cant_Recreate_View
               The object named by Target_Universe_View already exists.
         o Bad_Object_Index_Constraint
               The value of Max_Ada_Objects or Max_Other_Objects is not
___
               within the bounds prescribed by the parameters of the
--
               containing world.
     Is_Bad_Action (Status)
--
     Is_Lock_Error (Status)
     Is_Resource_Limit_Error (Status)
procedure Get_Universe_View_Slot
               (Universe_View : Om.Version_Handle;
                Action: Om.Action_Id;
                Bound_World_View : out Om.Version_Handle;
                bound_wv_Is_Protected : out Boolean;
                Indirect_World_View: out Om.Version_Handle;
                Status: out Error.Condition;
                Max_Wait : Duration := Job.Default_Wait);
procedure Set_Universe_View_Slot
               (Universe_View : Um.Version_Handle;
                For_world: Om.Object_Handle:
                Action : Um.Action_Id;
                Cound_world_View : Om.Version_Handle;
                Bound_wv_Is_Protected : Boolean;
                Indirect_world_View : Om.Version_Handle;
                Status: out Error.Condition;
                Max_Wait : Duration := Job.Default_Wait);
-- These operations are typically used to change "imports".
-- The get operation reads all 3 values from the identified slot. The
-- meaning of these values is defined at the beginning of this
-- package. The set operation writes all 3 values into the identified
-- slot. Client can combine the 2 operations to create operations
-- which change just a single value.
-- For_world identifies some object in the world whose corresponding
-- slot (in Universe_View) is to be read/written.
-- The get operation acquires a Record_Read lock on the identified
-- slot. The set operation acquires a Record_write lock on the slot
-- (or a Frite_Object lock if the action already holds a read/write
```

```
-- Possible errors:
     Is_bad_version_Handle (Status)
     Is_View_Error (Status)
___
         o World_View_Not_Bound
_ _
               Attempting to set Bound_wv_Is_Protected, but
               Bound_World_View is Nil.
     Is_bad_Action (Status)
___
     Is_Lock_Error (Status)
--
__
     Is_Resource_Limit_Error (Status)
procedure Expunge (Universe_View: Om.Version_Handle;
                   Number_Expunged : out Natural;
                   Disk_Blocks_Freed : out Natural;
                   Status: out Error.Condition);
-- Calls Object. Expunge on every object selected by the world view
-- (selected by Universe_View) in the same world as Universe_View.
-- On return, Number_Expunged and Disk_Blocks_Freed are returned equal
-- to the sum of the values returned by the calls to Object. Expunge.
-- Also, cleans up some internal OM garbage. Specifically: (1) Due to
-- outstanding locks at expunge time, it is possible for versions to
-- exist and not be selected by any view; these versions are expunsed.
-- (2) Also due to outstanding locks at expunge time, it is possible
-- for expungeable objects with no versions to exist; these objects
-- are expunged. (3) If an object goes from having 0 versions to 1 or
-- more versions, an empty Nil version is left around (invisible to
-- users); these versions are expunded. In future implementations,
-- these sources of garbage may be eliminated, and this paragraph
-- stricken.
-- Possible errors:
     Is_bad_Version_Handle (Status)
     plus all those returned by Object. Expunge.
procedure Destroy (View: Om. Version_Handle;
                   Number_Expunged : out Natural:
                   Disk_Blocks_Freed : out Natural;
                   Status: out Error.Condition);
-- Class World View: For a world view which is no Longer referenced,
-- this operation will destroy the identified version of the "_View_
-- object and expunse any version (and possibly object) which is
-- no longer referenced as a result.
-- (Except that the "_View_" and "_Links_"
-- objects are never expunged).
-- Class_Universe_View: This operation will destroy the identified
-- universe view and destroy any world view versions (as defined
-- apove) which are no longer referenced as a result.
-- Possible errors:
     Version_Is_Still_keferenced
         o Attempt to destroy a world view version which is still
            referenced by some universe view.
___
```

e:

-- unit name space.

```
package World is
       -- pragma Subsystem (Object_Management)
       package Om renames Om_Definitions;
       subtype Version_Retention_Count is Om.Version_Retention_Count;
       subtype Context is Om.Naming_Context;
       -- Basic operations --
       ------
       function Default_Unit return Disk.Unit_Number;
       -- Returns a legal Unit_Number according to a policy that is currently
       -- unspecified. Probably a policy similar to "best/usable volume" in
       -- Gamma.
       function Default_Cluster_Size return Disk_Block_Count;
       -- Currently 1.
       function Default_Warning_Threshold return Disk.Region_Count;
       -- Currently infinite.
       function Default_Allocation_Increment return Disk.Block_Count;
       -- Currently 1 track.
       procedure Create (Name : Om. Unique_Wildcard;
                         Max_Ada_Objects : Positive;
                         max_Other_Objects : Positive;
                         New_World : out Om.Object_Handle;
                         Status: out Error.Condition:
                         In_Context := Job.Default_Context;
                         Unit : Disk.Unit_Number := Default_Unit;
                         Cluster_Size : Disk.Block_Count := Default_Cluster_Siz
                         warning_Threshold : Disk.Block_Count := Default_Warnin
g_Threshold;
                         Allocation_Increment : Disk_block_Count := Default_All
ocation_Increment);
        -- Name and In_Context are used to specify the name of the new world,
        -- in a fashion similar to Object.Create.
        -- The following objects are created (as children of the world
        -- object): The object with simple name "_Views_" has
       -- Class_World_View: versions of this object are used to construct
        -- various "views" of the world. The object with simple name
        -- "_Links_" has Class_Link_Pack; versions of this object are selected
        -- by versions of the world view in order to represent the Ada library
```

```
-- The values of Max_Ada_Objects and Max_Other_Objects control the
-- size of a table: (of object information), which is stored in the
-- world object. If one specifies maximal values you chew up something
-- like a 1/3 of a Mbyte in the world object. These values constrain
-- the corresponding values when creating new views (via the
-- View.Create and View.Copy).
-- Some additional objects are created (as children of the world
-- object): The object with simple name "_Initial_" has
-- Class_Universe_View and is created by calling View. Create with the
-- appropriate string name and default parameters. The object with
-- simple name "_Default_" has Class_Universe_View and is created by
-- calling View.Create with the appropriate string name, Indirect to
-- the "_Initial_" view, Crippled, and remaining parameters taking
-- defaults.
-- The creation takes permanent effect immediately - there is no
-- action involved.
-- Note that the created world version is not put in any world view,
-- and there is no way to put it into one.
-- The System_View is updated to contain an indirect reference to the
-- "_Default_" view created for this world.
-- Possible errors:
     Is_Bad_Pathname (Status):
          o Might be because the last segment of the name contains
___
            wildcards.
___
     Illegal_Class
          o The "'C" attribute specified something other than a director
     Is_kesource_Limit_Error (Status)
          o Like ran out of world numbers.
procedure Destroy (Rame: Om. Unique_Wildcard;
                   Status: out Error.Condition;
                   In_Context : Context := Job.Default_Context;
                   Destroy_Contained_Objects : Boolean := False;
                   Destroy_Contained_Worlds : Boolean := False);
procedure Destroy (The_world : Om.Object_Handle;
                   Status: out Error.Condition;
                   Destroy_Contained_Objects : Boolean := False;
                   vestroy_Contained_Worlds : Boolean := False);
-- Name, and In_Context are used according to the rules of
-- Naming. Resolve to compute the Object_Id of the world to be destroyed.
-- All of the regions allocated to this world are made available to
-- the pool of free regions. This operation takes permanent effect
-- immediately. This operation DDES NOT follow action semantics. So,
-- once the world is destroyed there is no way to get the objects back.
```

-- WARMING: Supplying True to either of the destroy options will cause

-- This operation attempts to produce its errors prior to actually

-- the objects which live in the world to be destroyed.

```
-- destroying anything. However, the destroy operation is not atomic
       -- with respect to concurrent operations (like Object.Create);
       -- consequently, it is possible to get errors after contained objects
       -- have been destroyed.
       -- The entry (for this world) in the System_View is removed.
       -- Possible errors:
        __
             Is_Bad_Object_Id (Status)
        --
             Is_Bad_Object_Handle (Status)
        --
             Is_Bad_Pathname (Status)
             Is_mad_Action (Status)
       __
             Is_Lock_Error (Status)
        --
             Version_Is_Still_Referenced
        __
                 o Destroy_Contained_Objects is true, but the (sub) world
        --
                   contains at least one version which cannot be expunded.
       --
             World_Still_Contains_Objects
        --
                 o Destroy_Contained_Objects is false, but the world contains
        --
                   objects (with unexpunged versions) other than the link pac
       --
                   and world view.
       __
             World_Still_Contains_Subworlds
       __
                 o Destroy_Contained_Objects is false, but the world contains
                   objects which are themselves worlds.
       function System_View return Om.Version_Handle;
       -- Returns the id of the "system view".
       procedure Expunse (World: Om.Unique_Wildcard:
                           Number_Expunged : out Natural;
                           Disk_Blocks_Freed : out Natural;
                           Status: out Error Condition);
       procedure Expunge (world : Om.Object_Handle;
                           Number_Expunged : out Natural:
                           Disk_Blocks_Freed : out Natural;
                           Status: out Error.Condition);
       -- Calls View.Expunge (not Object.Expunge!) on every universe view in
        -- the identified world. Number_Expunged and Disk_Blocks_Freed are
       -- returned equal to the sum of the values returned by the calls to view
.Expunge.
        -- Possible errors:
        __
              Illegal_Class
        ___
                  o Can only be applied to an object with Class_Directory and
        --
                    subclass world.
              All those returned by View. Expunge.
        procedure Move (Name : Om.Unique_Wildcaro:
                        lo_Unit : Disk.Unit_Number:
                        Status : out Error.Condition;
                        ln_Context : Context := Job.Default_Context);
        procedure fove (world : Om_Object_Handle;
                        To_Unit : Disk.Unit_Number:
```

Status: out Error.Condition; In\_Context := Context := Job.Default\_Context); -- Name and In\_Context are used according to the rules of -- Naming.Resolve to compute the Object\_Id of the world to be moved. -- Moves the specified world to the specified disk. Requires that -- there be no locks of any kind on any object in the world. Ok if the -- world is already on that disk. Atomic with respect to concurrent -- activity. Atomic with respect to crashes. Probably not available in -- the 1st atr implementation. -- Possible errors: \_ -Illegal\_Class o Can only be applied to an object with Class\_Directory and -subclass world. Is\_Bad\_Object\_Id (Status) \_\_\_ Is\_Bad\_Pathname (Status) --Is\_Bad\_Action (Status) --Is\_Lock\_Error (Status) Is\_Resource\_Limit\_Error (Status) -- world parameters ---- These functions return garbage (typically 0) when given garbage -- input parameters. Procedures which do not return status are noops -- when the input parameters are garbage. The set operations take -- permanent effect immediately. These operations take a world -- parameter, which can identify any object in the world of interest. function Cluster\_Size (World: Om. Gbject\_Handle) return Disk.Block\_Count; -- Returns the cluster size for this world. Currently, we only -- support a cluster size of 1. It is not possible to change the -- cluster size once a world has been created. function Warning\_Threshold (World : Om.Object\_Handle) return Disk.Block\_Count; procedure Set\_Warning\_Threshold (World : Om.Object\_Handle; Threshold : Disk\_Region\_Count); -- Offending sessions get a warning message when the number of -- allocated regions (in this world) goes above this value. Default: -- Region\_Count\* Last. function dax\_Header\_Allocation\_Increment return Disk\_3lock\_Count; function Object\_Heaver\_Allocation\_Increment (World : Om.Object\_Hanale) return Disk.Block\_Count; procedure Set\_Object\_header\_Allocation\_Increment (world: Om.Object\_Handle: Increment : Disk.Block\_Count);

-- To minimize crash recovery time, it is advantageous for object

-- headers to be allocated contiguously on the disk. When a new object -- neader must be allocated, and the world does not contain any free

function Data\_dtock\_Capacity

inq

um

```
-- object neaders, the system allocates N plocks for use as object
-- headers, where N is the "allocation increment" parameter controlled
-- by these 2 operations. Larger values of the allocation incrment
-- tend to increase contiguity, but also increases fragmentation
-- (since free object header blocks cannot be used for anything else).
-- The value must not exceed Max_Header_Allocation_Increment.
function Max_Ada_Objects (World: Om.Object_Handle) return Positive;
procedure Set_Max_Ada_Objects (World : Om.Object_Handle;
                               Max_Ada_Objects : Positive;
                               Status : out Error.Condition);
function Max_Other_Objects (World: Om_Object_Handle) return Positive;
procedure Set_Max_Other_Objects (World : Om.Object_Handle;
                                 Max_Other_Objects : Positive;
                                 Status: out Error.Condition);
-- As indicated under the Create operation, the value of these
-- parameters constrains the corresponding values in View. Create and
-- View.Copy operations.
-- Possible errors:
       Is_View_Error (Status)
--
           o Bad_Object_Index_Constraint
                 Can be returned for any of the following: The given
--
                 value of Max_Ada_Objects is smaller than the correspond
                 value for some version of the world view for this world
                 Similar situation for Max_Other_Objects. The sum of
                 Max_Ada_Objects and Max_Other_Objects exceeds the maxim
                 number of object indices.
-- World information --
-- These operations return garbage when the supplied world_Number
-- refers to a world which does not exist. These operations take a
-- World parameter, which can identify any object in the world of
-- interest.
function Unit (World: Om.Object_Handle) return Disk.Unit_Number;
function Is_A_Job_world (world : Om.Object_Handle) return Boolean;
function Consumed_Regions (World: Om.Object_Handle)
              return Disk.Region_Count:
-- Returns the number of regions which are currently allocated to
-- Forid.
subtype block_Count is Disk_Elock_Count;
```

function Flock\_Capacity (World : Om.Object\_Hangle) return Block\_Count; function Header\_Capacity (World : Om.Object\_Handle) return Block\_Count;

page 17 (World : Om.Object\_Handle) return Block\_Count; function Consumed\_Headers (World : Om.Object\_Handle) return Block\_Count; function Consumed\_Data\_Blocks (World: Om.Object\_Handle) return Block\_Count; function Structural\_Map\_Size (World : Om.Object\_Handle) return Block\_Count; function Version\_Map\_Size (World : Om.Object\_Handle) return Block\_Count; -- Block\_Capacity is simply Consumed\_Regions times the number of -- blocks in a region. Header\_Capacity is the number of blocks -- currently reserved for use as headers, and Data\_Block\_Capacity is -- the number of blocks remaining for other uses. Slock\_Capacity less -- Header\_Capacity and Data\_Block\_Capacity identifies the built in -- overhead of a world. Consumed\_Headers is the number of header -- blocks currently in use. Similarly, Consumed\_Data\_Blocks is the -- number of blocks currently in use for other things. Each of the -- following expressions represent different kinds of internal -- fragmentation: (1) header\_Capacity - Consumed\_Headers (2) Data\_Block\_Capacity - Consumed\_Data\_Blocks (3) (Header\_Capacity + Data\_Block\_Capacity) = (Consumed\_Headers + Consumed\_Data\_Blocks) -- Structural\_Map\_Size and Version\_Map\_Size represent the only other -- forms of overhead in the world which is not accounted for by -- looking at the amount of space consumed by individual versions of -- objects. -- Window of Vulnerability Operations ---- These operations are noops when given invalid World\_Number's. function Is\_Vulnerable (World : Machine.World\_Number) return Boolean; procedure Make\_Vulnerable (World: Machine.World\_Number); procedure Make\_Invulnerable (World : Machine\_World\_Number);

- -- A world enters its window of vulnerability when its contents are -- modified. Vulnerable worlds have to be examined, object by object,
  - -- by crash recovery. Thus, causing a world to leave its window of -- vulnerability may reduce the cost of crash recovery. There is a
  - -- daemon which attempts to use idle cycles to remove worlds from
  - -- their window of vulnerability.

end World:

package Access\_Control is

- -- pragma Subsystem (Object\_Management);
- -- The following is a summary of proposed access control rules.
- -- General rules --

1 %

```
-- Conventional definition for user. Max 1023 users. Jser "George"
-- exists iff there is an object with name "!Users...Seorge", where
-- "!Users" is a predefined world, "George" is a direct child world of
-- "!Users", and "..." can be an intervening chain of 0 or more
-- non-world directories. By default, only the Privileged group has
-- the ability to create/delete users. OM will not allow one to create
-- any other objects in the "!Users" world.
-- Conventional definition for group. Max 4095 groups. Group
-- "Software" exists iff there is an object with name
-- "!Groups.Software", where "!Groups" is a predefined world, and
-- "Software" is a directory with Group_Subclass. George is a member
-- of Software iff there is a directory with the name
-- "!Groups.Software.George", with Member_Subclass. By default, only
-- the Privileged group has the ability to create/delete/modify
-- groups. OM will not allow one to create any other objects in the
-- "!Groups" world.
-- Every user is implicitly a group. Implicitly defined group Public
-- includes all users on the machine. Implicitly defined group
-- Network_Bublic includes all users with account on remote machines.
-- Predefined users Operator and Rational.
-- Predefined group Privileged includes users Operator and Rational.
-- Can use basic mechanisms to include arbitrary users in the
-- Privileged group. The Privileged group is special in the sense that
-- the privileges conferred by being a member of this group have to be
-- explicitly enabled. When Privileged capabilities are enabled, there
-- are no access control restrictions in force.
-- As indicated by the Job package, each session/job has an associated
-- "owner" user id. In addition, there is an amplification list of up
-- to 3 user ids. The rights of the session/job are the union of the
-- rights of the owner and the users named in the amplification list.
-- Recall that jobs are initialized to have this state equal to that
-- of the creating session.
-- ACLs are defined in OM_Definitions, and should be fairly obvious.
-- ACLs are applied to individual versions. There is no object-level
-- ACL.
-- The following privileges are defined:
--
     Create
     Delete
     Pead
     Write
     Execute
__
__
     Owner
-- The specific meaning of the privileges is defined below:
     Create
                  Applies only to worlds and controls the amility to
--
                  create views and direct world descendents.
--
    Delete
```

Applies only to worlds and views. For worlds,

controls the ability to destroy this world. For views, controls the ability to destroy the view \_ \_ itself. ----Applicable to all objects except directories (and Read therefore worlds), and link packs. For world views, controls the ability to open objects, via the view. --For other objects, controls the ability to open the \_\_\_ specific version for reading. Applicable to all objects except directories (and Write \_\_ therefore worlds). For world views, controls the -ability to modify the view; this includes -creation/deletion of objects/versions via the view. --For other objects, controls the ability to open the specific version for modification. --Applicable only to load images and Ada units/images. Execute \_\_ Further discussion below. \_\_ Applicable only to worlds. Controls the ability to Owner change ACLs on versions stored in the world, as well -as other miscellaneous things.

-- Note that worlds have an associated "owner" which is used for -- accounting. The world's owner need not have Owner access granted in -- the ACL of the world object (although this would certainly be true -- by default).

- -- world views contain a default ACL for new objects/versions created -- in this view. Recall that objects/versions (except for views and -- worlds) cannot be created without doing it in a view. The world -- object contains defaults for the ACL of views as well as for the -- default ACL contained in the view.
- -- The ACL of a new version of an object which is currently "deleted"
- -- from the perspective of the view is set from the view's default
- -- ACL; this also covers the case where the object is being created.
- -- The ACL of a new version (of an object which already "exists" from
- -- the perspective of the view) is identical to that of the
- -- predecessor version.

-- Initial ACLs:

-- World Operations --

-- Create:

-- Requires create access to the parent world.

-- Destroy:

- -- Requires destroy world access to the world being destroyed.
- -- Move: Requires owner access to the world being moved.
- -- Read world attributes:
- -- Does not require any access.

```
2 🗇
-- Write world attributes:
          Requires owner access to the world.
-- Change ACL:
          Requires owner access to the world.
-- Naming Operations --
-- No access control is applied.
-- Object Operations --
-- Create:
      (1) To create a new object requires write access to the view.
          The new version will be writeable by the creator only if the
___
          default ACL in the view so specifies.
__
      (2) To create a new version of an object requires write access to
          the view. The new version will be writeable by the creator
          only if the ACL of the predecessor version so specified.
  Open: Opening an object for read via a view requires read
          access to the view. Opening an object for update via a view
___
          requires write access to the view, regardless of whether or
__
          not differentiation is actually required. Note that when
          differentiation is performed, the new version will be
          writeaule by the creator only if the ACL of the predecessor
--
          version so specified. Recall that one can open a version
--
          without going through a view by simply specifying the
--
          correct version number. Also recall that many interesting
--
          classes of objects (such as Ada units) have a version policy
___
          which prohibits opening the version for update without going
___
          through a view. The ACL of the version itself must grant the
          desired access (read/write).
-- Delete:
          Requires write access to the view.
-- Indelete:
          Requires write access to the view.
-- Expunge:
          Requires write access to the view.
-- Copy: Requires write access to the target view. If the version is
--
          being copied by reference, no other acess is required.
___
          Otherwise, read access to the source view and/or version is
          also required.
-- Read version attributes:
          Requires read access to the version.
-- *rite version attributes:
          Requires write access to the version. This rule does not
```

apply to change of ACL.

of the added user.

```
-- Change ACL:
          Requires owner privilege to the enclosing world. Does not
          require write access to the version itself.
-- View Operations --
-- Create:
___
          Requires create access to the containing world.
-- Copy: Requires read access to the source views, and create access
          to the containing world.
-- Read/Write Univ view:
          Requires read/write access to the universe view object.
-- Expunge:
         Requires write access to the view.
-- Links --
-- The ACL for a link pack has no effect when creating/deleting
-- internal links or resolving link names. The ACL for a link
-- pack only has effect when creating/deleting external links.
-- Note that ACLs for link packs are inherited according to the
-- same rules as for other objects.
-- Jobs --
-- Various operations, such as setting state and job-kill require
-- requestor to have a group map which includes the owner of the
-- job/session in question.
-- Execution --
_____
-- Elaporate program/subsystem:
      (1) Importing code segments named by the load image requires
___
          the executing job to have Execute access to the load image.
      (2) Importing module names from already elaborated subsystem
_ _
--
          requires (a) load image for imported subsystem enables
          module access check, and (b) the executing job has Execute
          access to the load image.
-- Amplification --
      (1) Job/session can amolify its rights by supplying the password
```

(2) The Privileged group does not normally give its members any

```
page
          additional capabilities. A Job/Session can enable privileges
--
          if user is a member of the Privileged group.
___
      (3) Amplify_nights bit in primary load image will add rights of
          user, identified by the load image, to the elaborating job;
--
          rights go away when elaboration finished; prevents more than
          1 main program or subsystem from being elaborated by the
--
          job; setting the amplification rights requires the executing
--
          job to have the identified rights.
--
      (4) Can acheive amplification by elaborating subsystem in its
--
          own job, and running called procedure's on tasks of the
          elaborated subsystem.
```

```
-- Scenarios --
------
```

-- Strict partition: \_\_\_ No problem other than one can "see" (in the Naming sense) the existence of controlled objects. -- Give read-only access to subsystem importers: No problem.

-- Give Fred edit access to only a set of Ada units in a world: --We assume that the Create/Destroy/Owner privileges (for the world) are not already given to Fred. Spawn a new view. Give \_\_ Fred Read/Write access to the view. Set the view's default --ACL to give Fred Pead/Write access to units he creates. For -all existing versions (in Fred's view) set the Read/Write -privileges in the version's ACL) appropriately. Only modify -versions whose ACLs actually need changing. Versions which -get their ACLs changed are automatically differentiated. Fred \_\_\_ can now edit selected units in the world. Fred might be able to install/code (some issues with ACLs on central databases \_\_\_ kept by tools).

-- Give Fred write access to a file in an unmanaged world, but prevent Fred from deleting the file. Can do it provided Fred is not allowed to create new versions.

```
-------
-- Miscellaneous Notes --
------
```

-- Note that read/write privileges in the universe view are NOT -- checked during Object operations, such as Create and Open. This is -- true even when there is indirection in the views.

-- Note that link packs have no ACL.

```
end Access_Control;
end Om_Services_1;
```

```
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       Р
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               A
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   GG
       PPPP
               AAAAA
       P
   G
               A
                    Á
   G
       P
               A
                    A
 GGGG
       P
                    A
```

User: GPA
Object: !DE

!DELTA\_KK.REV1\_0\_0.UNITS.KK.OM\_DEFINITIONS

Version: V(117) Request: 1276

Date: April 24, 1986 Queued: 11:27:10 AM Printed: 11:43:09 AM

```
!DELTA_kk.REV1_0_0.UNITS.Kk.OM_DEFINITIONS!V(117)
```

```
page 1
with Sys;
with Filler:
with Bounded;
with Machine;
package Om_Definitions is
    -- pragma Subsystem (Kernel);
    -- General guidelines for use of OM specs:
    -- (1) Unless explictly stated otherwise, no operation propagates
    -- exceptions.
    package System renames Sys;
    Nil : constant := 0;
    -- All visibly scalar types have 0 as their Nil value.
    -- Machine ids --
    type Machine_Id is new _ong_Integer range 0..2 ** 24 - 1;
    -- Guaranteed unique for all Rational machines. Stored in the NOVRAM.
    -- Unique numbers --
    type Boot_Number is new Long_Integer range 0..2 ** 16 - 1;
    -- The system boot number is incremented each time the virtual memory
    -- system tries to boot itself. This value is stored in the label of the
    -- root disk unit. It is initialized when that disk unit is
    -- formatted/replaces. The value is allowed to "wrap". Under normal use,
    -- this means that the system boot number is unique and monotonically
    -- increasing.
    type Unique_Sequence_Number is new Long_Integer range 0..2 ** 32 - 1;
    -- Starts at 1 each time the system boots. Incremented each time the
    -- function Get is invoked.
    type Unique_Id is new Long_Integer range 0..2 ** 48 - 1;
    -- record
                  : Poot_Number;
          Sequence: Unique_Sequence_Number;
    -- end record;
    function boot return Boot_Number;
    function To_Unique_Id
                  (Boot : poot_Number;
                   Sequence: Unique_Sequence_Number) return Unique_Id;
    function foot (Id : Unique_Id) return Boot_Number;
```

function Sequence (In : Unique\_Id) return Unique\_Sequence\_Number;

function Get return Unique\_Id; -- Within this run of the system, guarantees that the returned value -- will be ">" than all those returned by previous calls (even in the -- face of concurrent calls to Get). With high probability, -- this property is also true in the case of crashes and recovery -- (from backup). Costs roughly 3-4 usec to call this function. -- Concurrent calls are ok. Calls by abortable threads are ok. -- At a consumption rate of 1 id every msec, will run out of Id's in -- 49 days, at which point the system crashes. type Net\_Unique\_Id is private; -- record Machine : Machine\_Id; Boot : Boot\_Number: Sequence: Unique\_Sequence\_Number: -- end record; Nil\_Net\_Unique\_Id : constant Net\_Unique\_Id; function To\_Unique\_Io (Machine : Machine\_Id; Soot : Boot\_Number; Sequence: Unique\_Sequence\_Number) return Net\_Unique\_Id; function Extract\_Machine (Id : Net\_Unique\_Id) return Machine\_Id; function Boot (Id : Net\_Unique\_Id) return Boot\_Number; function Sequence (Id : Net\_Unique\_Id) return Unique\_Sequence\_Number; function Get return Net\_Unique\_Id; -- Get the same properties as the previous Get function. -- additional property that with high probability, the returned value is -- unique among the values generated by all Rational machines. -- Object\_ld, Version\_Id and Object\_Handle --\_\_\_\_\_ type Class\_Number is new Long\_Integer range 0..31; -- Uniquely identifies the "type" of data contained in the object. For -- example, an object containing a Diana tree probably has a different -- type than an object containing a text file. All versions of an object -- have the same class. : constant Class\_Number := 1; Class\_Directory : constant Class\_Number := 2; Class\_world\_view : constant Class\_Number := 3; Class\_Universe\_View -Class\_File : constant Class\_Number := 4; : constant Class\_Number := 5; Class\_Pipe Class\_Link\_Pack : constant Class\_Number := 6;

: constant Class\_Number := 7;

: constant Class\_Number := 8;

: constant Class\_Number := 9;

: constant Class\_Number := 11;

Class\_bad\_Image\_Object : constant Class\_Number := 10;

Class\_Ada

Class\_Code

Class\_Ada\_Attributes

```
: constant Class_Number := 12;
__
                       : constant Class_Number := 13;
Class_Tape
                        : constant Class_Number := 14;
Class_Terminal
                        : constant Class_Number := 15;
-- The parent-child relationship makes most sense when the parent is
-- a directory or Ada unit. However, any object, regardless of class,
-- may take the role of parent in the directory system.
-- All classes of object may have multiple versions, except:
      World, Directory, and Device_Classes.
type Subclass_Number is new Long_Integer range 0..2 ** 16 - 1;
-- Rational reserves the first 1024 and provides a service for assigning
-- subclasses to independent software vendors. The Nil subclass indicates
-- that there is no special interpretation for the object. Examples:
-- Subclasses for Class_Directory: World, Path ...
-- Subclasses for Class_File: Text, Switches, Editor_Macros, ...
-- Subclasses for Ada: @eneric_Procedure_Spec, Package_Body, ...
type Object_Index is new Long_Integer range 0..2 ** 14 - 1;
-- Uniquely identifies a set of objects, within a particular world. All
-- versions of the object have the same pathname, class, file subclass
-- (if appropriate), and Object_Index.
type Simple_Object_Id is new Long_Integér range 0..2 ** 32 - 1;
-- Uniquely identifies a set of objects, within this machine. Beware that
-- World_Number's and Object_Index's are subject to reuse. Thus, there
-- are dangling reference issues. In some parts of the system (like
-- Diana) there is (hopefully) sufficient obsolescence protection that the
-- system will never interpret a dangling Object_Index.
__
     record
         Class : Class_Number;
--
                 Filler.Filler_1;
__
         world: Machine.world_Number;
--
                 Filler.Filler_2;
--
         Index : Object_Index;
     end record
function Class (Id : Simple_Object_Id) return Class_Number;
function World (Id : Simple_Object_Id) return Machine.World_Number;
function Index (Io : Simple_Object_Id) return Object_Index;
function To_Simple_Object_Id (Class: Class_Number;
                              Forld: Machine.world_Number:
                              Index : Object_Index)
              return Simple_Object_Id;
-- Field extraction and aggregate operations.
type Simple_Object_Id_Array is array (Natural range (>) of
                                     Simple_Object_Ia:
type Version_\umber is new Long_Integer range 0..2 ** 24 - 1:
```

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```
-- Uniquely identifies a particular object, in the set of objects which
-- have the same name. Each member of the set must have a unique version
-- number assigned. Version numbers are assigned sequentially, in time,
-- and are never reused. Note, however, that objects can be updated
-- without changing the version number.
type Version_Id is private;
subtype Object_Id is Version_Id;
-- record
                : filler.Filler_8;
       Machine : Machine_Id;
       Instance : Unique_Id;
--
       Object
               : Simple_Object_Id;
       Version : Version_Number;
-- end record; -- 17 bytes
Nil_Version_Id : constant Version_Id;
Nil_Object_Id : constant Object_Id;
-- The Machine field is used to make the object/version is unique across
-- machines.
-- The Instance field is used to match against the Instance field in the
-- Segment_Descriptor of the referenced object and catches dangling
-- references caused by reuse of World_Number's and Object_Index's.
-- Note that a Version_Id does not indicate the generation of the version.
-- To do that, one needs to drag along the Last_Update_Id field from the
-- Segment_Descriptor for the version in question.
-- The following conventions are used for the identification of job
-- worlds: (machine => (the usual),
            Instance => (that of the job world),
            Object => (Class => Class_Directory,
                        world => <job #>>
                        Index => (contant for world roots>)
            Version => Nil)
-- The following conventions are used for the identification of job
-- segment: (Machine => <the usual>,
             Instance => <that of the segment>,
--
             Object => (Class => Class_File,
                         World => <job #>.
--
                          Index => <1 + segment kind>)
            Version => (segment #>)
function Instance (Id : Version_Id) return Unique_Id;
function Extract_Machine (Id : Version_Id) return Machine_Id;
function Object (Id : Version_Id) return Simple_Object_Id;
function Version (Id: Version_Id) return Version_Number;
function Class (Id : Version_Id) return Class_Number;
function World (Id : Version_Id) return Machine.World_Number;
function Index (Id : Version_Id) return Object_Index;
function To_version_ld (Instance : Unique_Id;
                        Class: Class_Number;
```

world : Machine.World\_Number;

ς

```
Index : Object_Index;
                        Version : Version_Number) return Version_Id;
function To_Version_Id (Instance : Unique_Id;
                        The_Machine : Machine_Id:
                        Class: Class_Number;
                        World: Machine.world_Number;
                        Index : Object_Index;
                        Version : Version_Number) return Version_Id:
function To_Version_Id (Instance : Unique_Id;
                        Object : Simple_Object_Id:
                        Version: Version_Number) return Version_Id;
function To_Version_Id (Instance : Unique_Id;
                        Machine : Machine_Id;
                        Object : Simple_Object_Id;
                        Version : Version_Number) return Version_Id;
function To_Version_Id (Object : Object_Id;
                        Version : Version_Number) return Version_Id;
function To_Object_Id (version : Version_Id) return Object_Id;
type Version_Id_Array is array (Natural range (>) of Version_Id;
type Version_Handle is private;
subtype Object_Handle is Version_Handle;
-- record
       Object : Version_Id:
       Universe : Version_Id;
-- end record;
Nil_Version_Handle : constant Version_Handle;
Nil_Object_Handle : constant Object_Handle;
-- It is presumed that Version. Machine and Universe. Machine are identical.
-- As the source of an operation: If Object_Handle.Object.Version is not
-- Nil, then the specified version is used, and Object_Handle.Universe
-- is ignored. Otherwise, the Version_Number comes from the selected
-- universe/world view; the "selected view" is specified by the .Universe
-- component or job state when the .Universe component is Nil.
-- As the target of an operation which does not create a new version, the
-- rules are identical to those above.
-- As the target of an operation which creates a new version,
-- Object_Handle.Object.Version is ignored, and the new version is created
-- in the selected universe/world view; the definition of "selected" is
-- identical to that above.
-- The subtype Object_mandle is used in contexts which simply ignore
-- Object_Handle.Object.Version and Object_Handle.Universe, or in contexts
-- which return Nil for those components. The subtype Version_Handle is
-- used in the remaining contexts.
function Object (handle : Object_Handle) return Version_Id;
function Universe (Handle: Object_Handle) return Version_Id;
```

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```
function To_Object_Handle (The_Object : Version_Id;
                          In_Universe : Version_Id) return Object_Handle;
type Version_Handle_Array is
   array (Natural range (>) of Version_Handle;
-- Directory names --
subtype Bounded_String is Bounded.Variable_String;
Max_Simple_Name_Length : constant := 63;
subtype Simple_Name is Bounded_String (Max_Simple_Name_Length);
subtype Simple_Name_String is String;
subtype Padded_Simple_Name_String is String (1..63);
                   : constant Character := * *;
Pad_Character
-- Names in the directory system have a maximum length, and are
-- occasionally padded with blanks. System objects often have a """ at
-- the beginning. For example, the code segment for package "Foo" might
-- be called "'Code", and be a child of "Foo".
-- Every object has a name, even temporary heads, system objects, etc.
-- All versions of the same object have the same value for Name.
                   : constant Simple_Name_String := "";
Nutt_Name
-- A simple name is legal iff
           (1) Name'length <= max_simple_Name_Length.
       --
        and (3) Name ">" Null_Name (1 .. Name Length)
-- Should just stick the BNF here.
Max_Pathname_Length : constant := 511;
subtype Pathname is dounded_String (Max_Pathname_Length);
subtype Pathname_String is String;
-- A full pathname is a sequence of simple names separated by dots.
-- Should stick the BNF here.
subtype Simple_Wildcard is String;
-- A simple name plus wildcard characters.
-- Should stick the BNF here. Special characters restricted to just
-- "#" and "a".
subtype Full_Wildcard is String;
-- A full wildcard name spec.
-- Should stick the ENF here.
subtype Unique_hildcard is Full_wildcard;
-- By convention, must resolve to a single object/version.
subtype Naming_Context is Object_Handle;
```

-- Can be used as the starting context for various operations which

```
!DELTA_KK.REV1_0_0.UNITS.KK.OM_DEFINITIONS!V(117)
                                                                      page
    -- work on pathnames.
   -- Version Control --
   Max_Retained_Versions : constant := 15;
   subtype Version_Retention_Count is Natural
         range O.. Max_Retained_Versions;
    -- A World_View can keep track of up to Max_Retained_Versions of any
   -- particular object.
   type Version_Control_Policy is new Long_Integer range 0..2 ** 12 - 1;
   -- Applied to individual objects. See Object.Create.
   Non_Nil_Version_Control_Policy
                                         : constant Version_Control_Policy :=
   -- Kludge bit to satisfy the Nil=O convention.
   Allow_Explicit_Create_Version
                                         : constant Version_Control_Policy :=
   -- When false, Object.Create will refuse to create new versions.
   Private_Version_On_First_Update
                                         : constant Version_Control_Policy :=
                                         4: .
   Private_Version_Un_Every_Update
                                         : constant Version_Control_Policy :=
                                         8:
   Implicit_Versions_Replace_Current
                                         : constant Version_Control_Policy :=
                                         16:
    -- These 3 bits control when and how Object.Open (for update) creates
    -- implicit versions. See that spec for details.
   Freeze_Versions : constant Version_Control_Policy :=
   -- Controls whether or not world. Freeze causes selected versions to be
    -- frozen-
   Prevent_Update_Without_View
                                         : constant Version_Control_Policy :=
                                         64:
    -- When set, prevents Object. Open for update with non-nil version #.
    Shared_Version_Policy
                                         : constant Version_Control_Policy :=
                                         Non_Nil_Version_Control_Policy;
   Standard_Differential_Version_Policy : constant Version_Control_Policy :=
                                         Non_Nil_Version_Control_Policy +
                                         Allow_Explicit_Create_Version +
                                         Private_Version On_First_Update +
                                         Implicit_versions_Replace_Current +
                                         Freeze_Versions +
                                         Prevent_Jpdate_without_View;
    type Reservation_Bits is array (Object_Index range <>) of boolean;
```

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function To\_Job\_Id (Instance : Unique\_Id;

```
-- Info about Ada units --
type Compilation_State is
    (UU)
    Unparsed,
                      -- Not fully parsed
    Parsed,
                      -- Syntactically correct, not fully semanticized
    Installed,
                      -- Fully semantized
    Phase_1_Coded,
                      -- Unit attributed by the code generator
    Coded,
                      -- Code has been generated
    U6. U7);
type Ada_Unit_State is
   record
        State
                             : Compilation_State;
        Potentially_Obsolete: Boolean;
        -- An imported unit may have changed its meaning in a way that
        -- invalidates the current interpretation of the unit's semantics.
        Modified
                             : Boolean;
        -- Declarations have been edited which may have changed the meaning
        -- of one or more declarations exported by this unit.
    end record:
type Dependency_Matrix is
    array (Object_Index range <>>.
           Object_Index range (>) of Boolean;
-- DM(J/K) is true iff unit J contains a semantic pointer to unit K.
-- Additional information about the nature of the reference from J
-- to K is stored in unit J.
type Dependency_Vector is array (Object_Index range (>) of Boolean;
-- Can contain either a row or column of the Dependency_Matrix.
-- Jobs and Logins --
subtype Job_Number is Machine.Job_Worlds:
function Number (Tsk: Machine.Task_Id) return Job_Number;
type Job_Id is new Long_Integer range 0..2 ** 58 - 1;
-- record
_ _
        Instance : Unique_Id:
        Job
                 : Job_Number;
-- end record;
-- Jon_Number uniquely identifies a job. Use of Job_Id protects one from
-- the fact that Job_Number's are reused.
function Instance (Id : Job_Id) return Unique_Id;
function Number (Id : Job_Id) return Job_Number:
```

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```
Job : Job_Number) return Job_Id;
type Job_Kind is new Long_Integer range 0..2 ** 4 - 1;
Terminated_Job : constant Job_Kind := 1;
Attached_Job : constant Job_Kind := 2;
Detached_Job : constant Job_Kind := 3;
Server_Job : constant Job_Kind := 4;
Ce_Job
             : constant Job_Kind := 5;
0e_Job
              : constant Job_Kind := 6;
type Job_Priority is new Long_Integer range 1..7;
-- Priority 7 is "better" than priority 1. These values are used by the
-- medium term scheduler. Not the same as Machine.Cpu_Priority_Level,
-- which is used by the short term scheduler.
type Job_State is new Long_Integer range 0..2 ** 4 - 1;
Idle_Job : constant Job_State := 1: -- no unblocked tasks
Running_Job : constant Job_State := 2: -- has unblocked tasks
Waiting_Job : constant Job_State := 3: -- withheld by mts
Disable_Job : constant Job_State := 4: -- disabled by external agent
www.dueued_Job : constant Job_State := 5; -- detached and withheld by mts
type Session_Number is new Long_Integer range 0..2 ** 8 - 1;
type Session_Id is new Long_Integer range 0..2 ** 56 - 1;
-- record
        Instance : Unique_Id;
--
        Session : Session_Number;
-- end record;
-- The Session_Number uniquely identifies each "session". Use of
-- Session_Id protects one from the fact that Session_Number's are reused.
function Instance (Id : Session_Id) return Unique_Id;
function Number (Id: Session_Id) return Session_Number;
function To_Session_Id (Instance : Unique_Id;
                       >ession : Session_Number) return Session_Id;
-- Execution --
type Compatability_Index is new Long_Integer range 0..2 ** 16 - 1;
type Compatability_Vector is
    array (Compatauility_Index range <>) of Boolean;
type Compatability_Key (Static_Length : Compatability_Index := 256) is
    record
                     : Compatability_Vector (1..Static__ength);
        Dynamic_Length : Compatability_Index;
    end record:
type Exception_Information is
    record
        Raise_Fc
                        : Machine.Address;
        Exception_Number : Long_Integer;
        Other_Info
                        : Long_Integer;
```

```
page
    end record;
-- Users, Groups and ACL's --
type User_Number is new Long_Integer range 0..2 ** 10 - 1;
type User_Id is new Long_Integer range 0..2 ** 58 - 1;
--
   record
        Instance : Unique_Id;
--
                 : User_Number:
-- end record;
-- The Instance and User_Number are computed from the directory object
-- which corresponds to the user.
function Instance (Id : User_Id) return Unique_Id;
function User (Id : User_Id) return User_Number;
function To_User_Id (Instance : Unique_Id;
                      User : User_Number) return User_Id;
type Group_Number is new Long_Integer range 0..2,** 12 - 1;
Public : constant Group_Number := 1;
type Group_Id is new Long_Integer range 0..2 ** 60 - 1;
--
     record
         Instance : Unique_id;
___
         Group
                 : Group_Number:
     end record;
-- The Instance and Group_Number are computed from the directory object
-- which corresponds to the group.
function Instance (Id : Group_Id) return Unique_Id;
function Group (ld : Group_Id) return Group_Number;
function To_Group_Id (Instance : Unique_Id;
                        Group : Group_Number) return Group_Id;
type Group_Map is array (Group_Number) of Boolean;
type Acl_Entry_Kinds is new Long_Integer range 0..3;
Prohibit_Access : constant Acl_Entry_Kinds := 0;
                 : constant Acl_Entry_Kinds := 1;
Grant_Access : constant Acl_Entry_Kinds := 1;
Perform_Audit : constant Acl_Entry_Kinds := 2;
-- code 3 reserved for future use
type Privilege_Bits is new Long_Integer range 0..2 *# 5 - 1;
Create_Privilege : constant Privilege_Bits := 2;
Delete_Privilege : constant Privilege_Bits := 1;
Kead_Privilege : constant Privilege_Bits := 2;
Write_Privilege : constant Privilege_Bits := 4;
```

Execute\_Privilege : constant Privilege\_Bits := 1; Owner\_Privilege : constant Privilege\_Bits := 4;

-- codes 8 % 16 reserved for future use

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```
type Acl_Entry is
   record
       Kind : Acl_Entry_Kinds;
       -- Prohibit_Access causes the failure of an operation which
       -- requires the rights listed below and is requested by the
       -- identified Group, even if there is an entry later in the list
       -- which grants the desired access. A user in the Privileged group
       -- with privileges enabled is never denied access, even if such an
       -- entry exists.
       -- Grant_Access allows an operation which requires the rights
       -- listed below and is requested by the identified Group. There
       -- is an implicit entry giving the Privileged group all rights
       -- (this entry is not physically represented in the ACL). If no
       -- entry grants the desired access, the operation fails.
       -- Perform_Audit means that the system will audit any opreration
       -- which requires the rights listed below and which is attempted
       -- and/or performed by the identified Group; the audit entry goes
       -- to the system log. The fact that an operation is audited does
       -- not effect whether the request meets the requirements specified
       -- by the remainder of the ACL. Auditing is performed even if the
       -- requester has Privileged access enabled.
       -- It is recommended (for performance reasons) that one place the
       -- entries which are most likely to match at the beginning of the
       -- list (subject to ordering by Kind, of course).
       Group : Group_lumber:
       -- The group governed by this entry. Nil terminates the list.
       -- Completely filling the array also terminates the list.
       Rights : Privilege_Sits;
       -- See visible part of the Access_Control spec, and the various
       -- Om operations for details.
   end record;
type Access_Control_List is array (Natural range 1..6) of Acl_Entry;
pragma Assert (Access_Control_List*Size = 120);
Hull_Acl_Entry : constant Acl_Entry :=
               (Kind => Nil, Group => Nil, Rights => Nil);
Null_Act
               : constant Access_Control_List :=
               (others => Null_Acl_Entry);
-- An element containing a wil Group is considered to terminate the
-- list.
-- Note: moving an object from machine to machine or even to/from tape
-- makes for interesting issues with respect to ACLs in the target copy.
```

<sup>--</sup> Actions --

```
type Action_Number is new Long_Integer range 0..2 ** 12 - 1;
-- Used as an index into the action manager's internal tables.
-- OM subsystem reserves the right to change the *size of this type.
type Action_Id is new Long_Integer;
-- Handle for an open action.
    record
         Instance : Unique_Id;
                 : filler.filler_4;
--
         Index
                 : Action_Number;
    end record
-- "
function Boot (Id : Action_Id) return Boot_Number;
function Instance (Id : Action_Id) return Unique_Id;
function Index (ld : Action_Id) return Action_Number;
function To_Action_Id (Instance : Unique_Id;
                       Index : Action_Number) return Action_Ia;
type Lock_Mode is new Long_Integer range 0..2 ** 4 - 1;
                         : constant Lock_Mode := 1;
Unsynchronized
Read_Object
                         : constant Lock_Mode := 2;
Write_Object
                         : constant Lock_Mode := 3;
Record_Read_Object
                         : constant Lock_Mode := 4;
Record_write_Object
                         : constant Lock_Mode := 5;
Record_Read
                          : constant Lock_Mode := 6;
Record_write
                          : constant Lock_Mode := 7;
Supersedeable_Read_Object : constant Lock_Mode := 8;
-- Documented in Action spec.
subtype Record_Key is Long_Integer range 0..2 ** 32 - 1;
-- When acquiring action locks, this value is used to identify a "record",
-- within some object.
-- Miscellaneous --
type Relational_Op is (Lt, Le, Eq, Ne, Ge, Gt);
type Format_Number is new Long_Integer range 0..2 ** 8 - 1;
-- Used to tag various representations to facilitate detecting disk
-- incompatable changes.
type Segment_Position is new Long_Integer range 0..2 ** 32 - 1;
-- Values of type Machine. Bit_Offset are used in Machine. Addresses to
-- denote the physical position within a segment. Values of type
-- Segment_Position are used to denote the logical position within a
-- segment. The two values differ by the amount of ucode 3 om overhead
-- present at the beginning of the segment.
```

end Om\_Definitions;

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```
First_Valid_Segment_Position : constant Segment_Position := 1;
   -- The Segment_Position type follows the usual rule for interpretation of
   -- Nil. The first bit of "user data" in a segmented heap occurs at
   -- Segment_Position*(1).
   type File_Organization_Kinds is new Integer range 0..15;
   -- Identifies the internal organization of a file, as defined by the file
   -- system. Hypothetical example: a file might be organized as raw bits (as
   -- in the Gamma system), or it might be organized as a sequence of records
   -- (which retain length information and support record level action
   -- semantics. In either case, one could imagine that a subclass of text
   -- would make sense; in the first organization, one might have raw
   -- characters in the file, with special characters used to denote line
   -- boundaries, etc; in the second organization, one might have lines
   -- stored in individual records.
   -- Private --
   ---------
private
   type Net_Unique_Id is
       record
           Machine : Machine_Id;
           Boot : Boot_Number:
           Sequence : Unique_Sequence_Number;
       end record:
   Nil_Net_Unique_Id : constant Net_Unique_Id := (0, 0, 0);
   type Version_Id is
       record
           Fill
                : Filler.Filler_8;
           Machine : Machine_la;
            Instance : Unique_Id;
                   : Simple_Object_Id:
           Object
           Version : Version_Number;
       end record:
   Nil_Version_Id : constant Version_Id := (0, 0, 0, 0, 0);
   Nil_Object_Id : constant Object_Id := Nil_Version_Id;
   type Version_Handle is
       record
           Version : Version_Id;
           Universe : Version_Id;
       end record;
   Nil_Version_mandle : constant Version_Handle :=
                      (Vil_Version_Id, Nil_Version_Id);
   Nil_Object_Handle : constant Object_Handle :=
                       (Nil_Version_Id, Nil_Version_Id);
```

```
bbbb
 GGG
                AAA
G
   G
       Ρ
            þ
               A A
G
       Р
                    A
       pppp
   GG
               AAAAA
G
   G
       P
                   Α
G
   G
       P
               A
                    Α
       P
 GGGG
                    Α
```

User: GPA Object: !DE

!DELTA\_KK.REV1\_0\_0.UNITS.KK.KERNEL\_SERVICES\_1

Version: V(4)
Request: 1277

Date: April 24, 1986 Queuea: 11:27:18 AM Printed: 11:44:04 AM

```
with Sysi
with Machine;
package Kernel_Services_1 is
    -- pragma Subsystem (Kernel, Private_part => Closed);
    -- pragma Module_Name (4, ?);
    package Byte_String_Conversions is
        package System renames Sys:
        function To_String (B : System.Byte_String) return String;
        function To_Byte_String (S: String) return System.Byte_String;
    end byte_String_Conversions;
    package Wait_Service is
        type Wait_State is private;
        Nil_wait_State : constant Wait_State;
        function Register (Wait_State_Name : String) return Wait_State;
    -- Defines a new wait state. If the given string is too long, or
    -- already defined, or the database is full, returns Nil_Wait_State.
    -- The following are predefined:
          Port_Wait, Port_Input_Wait, and Port_Output_Wait
        function Get_Wait_State (Wait_State_Name : String) return Wait_State;
        function Get_Wait_State_Name ((State): Wait_State) return (String)
    -- Converts from string name to Vait_State, and vice-versa.
        procedure Roust (The_Client : Machine.Task_Id;
                         State: Wait_State);
        procedure Un_Roust (State: Wait_State);
        Forever: constant Duration := Duration*Last;
        procedure Wait_For_Roust (State : Wait_State;
                                  Max_Wait : Duration := Forever);
    -- Assume the existance of some service S. The clients of S, known as Ci,
    -- ask S to do work, and wish to be able to wait for the work to get done
    -- prior to proceeding. We use Satask to denote the "worker" task of the
    -- service S. we use S. skin to denote the procedures in the body of S
    -- which run on Ci's task, which is denoted Ci.task.
    -- Ci.task calls S, asking it to do some operation. S.skin should call
    -- Un_Roust; this serves the purpose of "throwing away" an outstanding
    -- roust that Ci.task did not see (these can happen because (a) the
    -- "roustedness" of a task is not initialized when the task is created,
    -- and (b) the task might have exited the skin of some service via delay
```

```
2
-- expiration). S.skin then calls S.task, which attempts the operation.
-- S.task then returns to S.skin a boolean indicating whether or not
-- Ci-task should suspend. If Ci-task is supposed to suspend, then S.skin
-- calls Wait_For_Roust, on behalf of Ci.task; Ci.task becomes blocked
-- inside wait_For_Roust. At some later point, the operation completes,
-- and Si-task invokes Roust; this will in turn cause Ci's call to
-- Wait_For_Roust to return; S.skin may then want to call S.task to get
-- the results of the operation. (Its also possible for the Roust to
-- happen prior to client invoking Wait_For_Roust.) It is also possible
-- that the Wait_For_Roust will return due to delay expiration, in which
-- case S.skin will need to call S.task in order to get the results of the
-- operation.
-- The wait service has, in effect, a boolean semaphore per client.
-- Therefore, when S. task plans to do a Roust on CI.task, S. skin should
-- either guarantee that Ci.task has not left S.skin, or have some other
-- guarantee that Ci.task is not attempting to wait on more than one event.
-- Roust is a noop when the specified client no longer exists.
-- Roust may be invoked by non-blocking tasks.
-- The wait state parameter to the Roust and Un_Roust procedures is used
-- to ensure that services only roust clients that "are their skin". That
-- is, try to roust someone that is not in the specified state, and you
-- will get constraint error.
-- Notes for non-blocking (ie., priority 0) tasks: If the tcb of the
-- rousted task is not in memory, the call to Roust will place a ghost call
-- in the queue space of the roust server (a task internal to this
-- package). When the roust server accepts the call, it will execute the
-- same roust, but at priority 1, thereby taking the necessary page fault.
-- The roust server has a finite amount of queue space reserved. In other
-- words, up to N rousts of non-resident tcbs can be "outstanding". Rousts
-- remain outstanding because (a) priority 0 tasks hog the cpu keeping the
-- roust server from running, or (b) the roust server is tied up in page
-- fault servicing one of the rousts. The current value for N is about
-- 128. Should the roust server's queue be full, the rousting priority 0
-- task will attempt to create a new queue page; if the page cannot be
-- created without calling the space manager, the system will crash.
-- Currently, the only non-blocking service is the Port_Manager. It is
-- believed that the propability of crashing in this scenario is very low.
private
    type Wait_State is new Integer range 0..2 ** 7 - 1;
    Nil_Wait_State : constant Wait_State := 0;
end Wait_Service;
package Trigger is
    -- Encapusulates the wait service protocal. Suitable for fairly
    -- trivial wait service applications.
    type Task_Count is new Natural range 0..2 ** 16 - 1;
```

type Descriptor (Count : Task\_Count) is private;

```
3,
!DELTA_KK.REV1_0_0.UNITS.KK.KERNEL_SERVICES_1"V(4)
                                                                      page
        procedure Initialize (Desc: in out Descriptor;
                              wait_State : Wait_Service.Wait_State);
        -- Recommend COUNT to be prime. You must INITIALIZE your descriptors.
        -- All the waiting tasks are assumed to be in the same wait state.
        procedure Pull (Desc: in out Descriptor);
        -- Rousts all the waiting tasks.
        generic
            with procedure Change (Tid: Machine.Task_Id: Is_Waiting: Boolean);
        procedure Wait_In_Skin
             (Wait_State : Wait_Service.Wait_State;
              Max_Wait : Duration;
              Non_Blocking_Service : Boolean := False);
        procedure Change (vesc: in out Descriptor:
                          Tid : Machine.Task_Id;
                          Is_Waiting : Boolean);
        -- Changes record of who is waiting for the trigger to be pulled.
        -- Feeu this to WAIT_IN_SKIN generic. Note that after calling change,
        -- you typically need to check if the "roust condition" holds, and if
        -- so pull the trigger.
        function Somebody_Is_Waiting (Desc : Descriptor) return Boolean;
    private
        type Waiter is
            record
                Tid : Machine.Task_Id;
            eno recordi
        type Waiter_Array is array (Task_Count range <>) of waiter;
        type Descriptor (Count : Task_Count) is
            record
                Wait_State : Wait_Service.Wait_State:
                           : Waiter_Array (1..Count);
                waiters
            end record:
    end Trigger;
```

end Kernel\_Services\_1;

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