

**GEODÆTISK INSTITUTS INTERNE RAPPORT NR. 14**  
**THE DANISH GEODETIC INSTITUTE**  
**INTERNAL REPORT NO. 14**

**Appendix 1.**

**USERS MANUAL FOR GPU**  
**General Processing Unit**  
**utilizing**  
**Double Processing Unit**

**General Processing Unit designed for double  
precision floating point arithmetic after Wilkinson.  
The double precision orders are processed by the DPU.**

**K. Engsager**

**1983**

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**References :**

- 1 GPU reference manual. K. Engsager. Geodætisk Institut 1982.
- 2 DPU reefernce manual. K. Engsager. Geodætisk Institut 1982.
- 3 GPU driverproc, GPU blockprocedure, GPU execute.  
K. Engsager. Geodætisk Institut 1982.
- 4 Algol 8. Users Guide. Part 2. RCSL 42 - i1278
- 5 Code procedures and the runtime organisaztion of ALGOL programs  
RCSL 31 - D119.
- 6 RC Slang Assembler programming Guide. RCSL 42 - i0785.

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## Preface

---

The GPU is a reprogrammed RC8000/35 central unit. The GPU executes a piece of code (slang), which is pointed by a message send to it. Some procedures have been designed to handle the device in a standard way in algol programs. E.g. `gpu_open`, `gpu_close` and `gpu_error`.

The `block_procedure gpu_error` will handle any error situation concerning the GPU in such a way that the user will know whether it is overflow, time\_out, bus-error, attempt to store in write protected area or attempt to execute instructions outside the pointed code area. If the installation has no GPU, the `block_procedure` will emulate the special `gpu-instruction` (in single precision only) using the escape facility (this might be a little slow).

## 2. Declarations

---

ALGOL

=====

The zone used to communicate with the GPU must be declared :

```
zone code(length, 1, gpu_error);
```

The name code is freely choosen among others.

The block\_procedure gpu\_error must be used because the GPU uses non standard status-bits in the answer.

NOT ALGOL

=====

An area may be laid out to the code and the communication may be handled by the message - answer monitor calls.

## 3. Open

-----

## ALGOL

=====

Use the procedure : gpu\_open(z, gpu\_name, code\_area, code);

Call: gpu\_open (return value, boolean) on the file code\_area is  
-----  
the codesegment(s) given by the name code searched  
and loaded to the zone\_buffer z. Then is the zone  
opened to the process gpu\_name or gpu\_dummy dependent  
of gpu\_emulat. The procedure will load code of any  
length roomed in the zonebuffer.  
result := code found else false.  
z (call and return value, zone) ready for open. At  
return ready for out\_rec.  
gpu\_name (call value, long array). name of gpu\_process.  
code\_area (call value, long array). name of code file.  
code (call value, short string). name of code to be  
searched for on last segment of code in long  
field 512.

Leading segments of code\_area are skipped until the segment having  
code\_area(1) = long z(127) and code\_area(2) = long z(128)  
are skipped (inclusive).

Integer field 2 of codes on codearea must contain the length of the  
code in bytes. A zero is replaced by 512.

gpu\_open will set the escape-mask used to emulate the gpu-instructions,  
when gpu does not exist.

A hard error during the actual execution of code will cause a  
break due to the fact that some code has been executed and the  
recovery will be very difficult if not impossible.

## NOT ALGOL

=====

The communication may be handled in a non standard way by the  
message - answer monitor calls. (see outrec).

#### 4. Close

ALGOL

=====

The communication by GPU is terminated by :

```
change_rec_6(code, 0); /* to ensure that no code is executed */  
close(code, true);
```

NOT ALGOL

=====

If no zone has been declared and opened there is no zone to close.

## 5. Out\_rec

---

ALGOL

=====

After open, set\_position, change\_rec\_6 or out\_rec\_6 a new block of code may be loaded to the gpu by a call of the procedure  
`out_rec_6(code, length);`

When the code is ready to be executed this may bee done in several ways :

```
if change_wanted then change_rec_6(code, no_of_code_bytes);
if pos_wanted then set_position(code, 0, 0) else
if close_wanted then close(code, true) else
    out_rec_6(code, length);
```

NOT ALGOL

=====

When a codearea has been created then it must be sent to the GPU by a message ad then GPU must be waited for by a wait\_answer.

The message buffer must contain the description of the codearea :

```
mess + 0 : next buffer
mess + 2 : previous buffer
mess + 4 : receiver or answer type
mess + 6 : sender
mess + 8 : operation = 5 < 12 <* write *>
mess + 10 : first code (see below)
mess + 12 : last code
```

GPU starts execution of code in word (first\_code + 4)  
+++++

The first 8 words are set by the monitor procedure.  
first\_code and last\_code are absolute addresses of the codearea.

6. Block-procedure                           gpu\_error

---

The block procedure will handle any error-situation from the GPU.

If the GPU is disconnected and gpu\_emulat permits it will the code be emulated by the block procedure, - but in single precision. The runtime will increase drastically.

If the GPU is disconnected during an execution of a code send from a process the process is breaked with an alarm and a message is sent to the operator.

The status bits from the GPU are :

bit no.	value	description
0	1 < 23	intervention
1	1 < 22	illegal instruction
2	1 < 21	time out
3	1 < 20	bus error
4	1 < 19	read/write address fault
5	1 < 18	floating point overflow *)
6	1 < 17	floating point underflow **)
7	1 < 16	integer overflow *)
8	1 < 15	integer underflow *)
9	1 < 14	code address fault ***)

10-23   see ref. 1.

\*) controlled by the statusbits of the senders process.  
e.g. integer underflow is only active when the senders process has set the status : integer exception active.

\*\*) in case of floating point underflow the result is set to floating point zero and no exception is set up.

++++++

\*\*\*) the next instruction to be executed is fetched from an address which lies outside the limits code\_low and code\_top - or the codearea-limits code\_low and code\_top are not found inside the write-limits llim and ulim of the senders process.

## 7. Gpu\_exec

---

The external procedure gpu\_exec emulates the gpu by executing the code.

Gpu\_exec is called from the blockprocedure if emulation is allowed, see gpu\_emulat.

The procedure may be called directly itself during debugging of the code to the GPU.

The execution of the special instructions is a little slow because the instructions are executed through the escape routine.

integer procedure gpu\_exec( code ).

---

gpu\_exec (return value, integer) the number of bytes processed.

code (call and return value, real array) the code to be emulated.

code 64 = gpu\_red will not be emulated.

The execution is locked in not\_blockfloating mode.

## 8. Gpu\_emulat

---

Own integer control parameter used by the blockprocedure  
gpu\_error.

The first two bits are controlbits and the last 22 are the  
number of emulations.

gpu\_open uses the first 2 bits of gpu\_emulat :  
= 3 : open is to gpu\_dummy. i.e. forced emulation.  
<> 3 : open is to gpu\_name. i.e. the GPU is used.

gpu\_error uses the first 2 bits of gpu\_emulat :  
= 0 : no emulation. i.e. break process.  
<> 0 : emulation. (unless the error situation is unrepairable).

## 9. Code

---

The GPU is a rebuilt RC 8000/35 which executes the normal RC 8000 instructions with some exceptions.

A sequence of instructions may be translated by slang into a code-file which may be input by an application program.

Example 1:

```
begin
    zone code(128, 1, std_error),
    -     gpu(128, 1, gpu_error);

    open(code, 4, <:codefile:>, 0);
    open(gpu , 0, <:gpul:>,      -2);

    in_rec_6(code, 512);
    out_rec_6(gpu, 512);
    to_from(gpu, code, 512);
    close(code, true);
    close(gpu, true);

end;
```

A sequence of instructions could also be placed in a codeprocedure, which moves the sequence to a given zone.

The code may be placed after the gpu\_exec segments and translated at the same time as that procedure (see the text gpu\_execx in the contract file gpu\_proc\_file).

The instructions must follow some simple rules :

---

- 1) The first instruction to be executed must be in byte 4 : code\_low + 4.
- 2) The instructions must be within the given limits  
code\_low = first of buffer and  
code\_high = code\_low + (length\_of\_buffer).
- 3) No privileged instructions may be executed except the special GPU-instructions.
- 4) The first word must contain the length in bytes (multiple of 512) when gpu\_open is used.

An easy way of composing a code is to use the gpu\_declaration\_block and follow the instructions given in page 19 ff.

## 10. Description of the gpu-instructions

```

the doubleword x (value real) is named : dwd(addr),
                                         where addr = absaddr(x)
modus_sign is set by gpu_init.
modus_blf1 is set by gpu_init.

gpu_add      (= 30)
+++++
ar := if addr = 0 then
-     (if modus_sign = 0 then (AR + (dwd(w) con 0.0))
-                  else (AR - (dwd(w) con 0.0)) ) else
-     (if modus_sign = 0 then (AR + (dwd(w) con dwd(w+addr)) )
-                  else (AR - (dwd(w) con dwd(w+addr)) ) );

```

gpu\_arm (= 59)  
+++++

AR := AR \* dwd(addr);

gpu\_cmv\_a (= 47)  
+++++

```

if w0 > 0 then
begin
  real x;
  x := dwd(addr);
  if x <> 0.0 then
    begin
      repeat
        AR      := + ( x * dwd(w_pre) ) + dwd(w);
        dwd(w) := AR;
        w_pre  := wpre + 4;
        w      := w      + 4;
        w0     := w0     - 1;
      until w0 = 0;
    end;
end;

```

gpu\_cmv\_s (= 31)

```

if w0 > 0 then
begin
  real x;
  x := dwd(addr);
  if x <> 0.0 then
  begin
    repeat
      AR      := - ( x * dwd(w_pre) ) + dwd(w);
      dwd(w) := AR;
      w_pre  := w_pre + 4;
      w      := w      + 4;
      w0     := w0     - 1;
    until w0 = 0;
  end;
end;

```

```

gpu_init    (= 63)
+++++++
begin
  integer md;
  md := word(addr);
  modulus_sign := (md shift (-22)) extract 1;
  comment 0: add, 1: sub ;
  modulus_bfl1 := (md shift (-21)) extract 1;
  comment 0: normalizing mode, 1: fix_point_mode with block_exp;
  see chapter 12;
  block_exp := md extract 12;
  if block_exp > 2047 then block_exp := ((-1) shift 12) + block_exp;
  if md shift (-23) extract 1 = 1 then AR := 0.0;
end;

gpu_inv     (= 60)
+++++++
begin
  comment the content of w_pre is destroyed;
  dwd(w) := 1 / dwd(addr);
end;

gpu_mla     (= 58)
+++++++
if w0 > 0 then
begin
  repeat
    AR := if modulus_sign = 0 then
      - (AR + dwd(w_pre) * dwd(w)) else
      - (AR - dwd(W_pre) * dwd(w));
    w_pre := w_pre + 4;
    w     := w     + 4;
    w0   := w0   - 1;
  until w0 = 0;
end;

gpu_red     (= 62)
+++++++
begin
  comment cholesky-reduction of a datamatic block of columns.

  The instruction must use w3 as regw and w1 as regx and the
  address must point to address c5 in the working area :
    gpured w3 xl c5.
  w1 need not be zero as it is subtracted from the address.

  In block_floating mode c6 and b8 is pointing the exponents
  which should be subtracted from the unreduced elements exponent
  to make the element a block_floating number with exponent equal
  zero (see chapter 13).

```

in normal mode c6 and b8 is dummy.

```

working area :
b0 = base of working area : CAT_I1U + nlu_base
c4 = b0 + 2 : AUTORED
c3 = b0 + 4 : LR = last_reduced column
b7 = b0 + 6 : (R_max + 1) Helmert block_red_limit.
n2 = b0 + 8 : SZU = Saved Zeroes in Unreduced coulmn
c1 = b0 + 10 : CAT_SZR + nlr_base, base of Saved Zeroes of
Reduced columns.
c0 = b0 + 12 : CAT_I1R + nlr_base, index of first nonzero
element in reduced column.
n11 = b0 + 14 : 4 * (R_MAX + 1) = 2 * word(b7)
b5 = b0 + 16 : tail_displacement.
c5 = b0 + 18 : accu_mode :
accumode(2) =
    if block_floating mode then 1 else 0.
    accu_mode(12:23) = if pos_accumulation then 0
    else ((-1) extract 12)
    <* neg_accumulation *>.
c6 = b0 + 20 : block_exp_r + nlr_base, base of catalog on
exponents of reduced columns to be subtracted
from the exponent of the element under reduc-
tion to make it be a block_floating number.
n13 = b0 + 22 : block_exp_u , to be subtracted from the
exponent of the element under reduction.
b8 = b0 + 24 : block_exp_u + nlu_base, base of catalog on
exponents of unreduced columns to be e.t.c.
b6 = b0 + 26 : exp_lim, max acceptable loss of binals.
b4 = b0 + 28 : status + nlu_base, base of catalog of status-
area of unreduced elements.
n0 = b0 + 30 : U, index of unreduced column.
b1 = b0 + 32 : cat_szu + nlu_base, base of catalog of saved
zeroes of unreduced columns
c2 = b0 + 34 : FR, first index of reduced column in datamatic
block.
b3 = b0 + 36 : LU, last index of unreduced column in datamatic
block.
b2 = b0 + 38 : FU, first index of unreduced column in data-
matic block.

```

Only the content of c5 is changed at return.

+++++  
+++++  
+++++  
+++++  
+++++  
+++++

The b-working\_locations refer to unreduced block.

The c-working\_locations refer to reduced block.

The n-working\_locations are pure working\_locations used by the microprogram.

;

```
<* check the registers used in call of bf *>
if reg_x <> w1 or reg_pre <> w2 or reg_w <> w3 then illop;
```

```

for col := fu step 2 until lu do
begin

  if auto_red then lr := col;
  szu := nlu.cat_szu.col * 2;
  if szu <= (R_max + 1) then
begin

  if bl_f1_mode then exp_u := nlu.cat_exp_u.col;

  for row := fr step 1 until lr do
begin

    szr := nlu.cat_szr.row * 2;
    if szr < szu then szr := szu;

    if szr < (R_max + 1) then
begin

      szr := szr + szr;
      reg_pre := nlr.cat_ilr + szr;
      reg_w   := nlu.cat_ilu + szr;

      ful_red := row*2 < (rmax+1)*2;
      w0 := (if ful_red then (2*row - szr)
              else ((R_max+1)*2 - szr) ) // 4;

      w1 := (1 shift 23) add
             ((accu_mode extract 1) shift 22) add
             (((accu_mode shift (-12)) extract 1) shift 21);
      gpu_init(w1);

      gpu_mla(w0,w);
      dia := reg_pre = reg_w;

      if bl_f1_mode then red_exp := nlu.cat_exp_u.row + exp_u
      else red_exp := 0;

      u_f := nlu.cat_ilu.row + row + row + tail_disp;
      x_t := nlu.u_f;
      if sign_extend(x_t extract 12) > -2048 or
         sign_extend(x_t extract 12) - red_exp > -2048 then
begin
      u_f := u_f - tail_disp;
      x_h := nlu.u_f;
      x_h := ((x_h shift (-12)) shift 12) add
             ((sign_extend(x_h extract 12) - red_exp) extract 12);
      if accu_mode extract 12 = 0 then gpu_add(x_h, x_t)
      else gpu_sub(x_h, x_t);
end
else
      x_h := 0.0;
    end
  end
end
end

```

```

if full_red then
begin
  if dia then
    begin
      wl := 0;
      gpu_init(wl);
      gpu_str(w, 0);

      if reg_pre <= 0 then
      begin
        s_f := status + 2 * col;
        nlu.s_f := (real <::> add
                     (if reg_pre = 0 then 4 else 3)) shift 24;
      end
      else
      begin
        exp_loss := (sign_extend(x_h extract 12)) -
                     (sign_extend(reg_w extract 12));
        s_f := status + col + col;
        nlu.s_f := ((real <::> add
                     (if exp_loss >= exp_lim then 3 else 1)
                     ) shift 24) add exp_loss;
        if exp_loss < exp_lim then
        begin
          gpu_sqrt(reg_w);
          wl := (if bl_f1_mode then 1 else 0) shift 21;
          gpu_init(wl);
          gpu_inv(reg_w, nlu.ilu+col+col);
        end
        else
        begin
          u_f := nlu.ilu + col + col;
          nlu.u_f := 0.0;
        end;
      end;
    end
  else
  begin
    gpu_arm(reg_pre);
    u_f := nlu.ilu + col + col;
    gpu_str(nlu.u_f, 0);
  end;
end
else
begin
  wl := 0;
  gpu_init(wl); /* normal mode */
  u_f := nlu.cat_ilu.col + col + col;
  gpu_str(nlu.u_f, tail_disp);
end;
end;
end;
end;

```

```

gpu_sqrt      (= 1)
+++++++
begin
  comment w_pre, w := sqrt( dwd(addr) );
  w_pre con w := sqrt( dwd(addr) );
end;

gpu_stop      (= 0)
+++++++
begin
  comment stop execution of user program in gpu and return
          status to host and user;
end;

gpu_str       (= 61)
+++++++
begin
  comment if the effective address is zero then store AR in dwd(w)
          else store AR as double floating point number in
          dwd(w) con dwd(w+addr).
  AR is unchanged;

  if addr = 0 then
    dwd(w) := float_rounded_to_36_bit(
      -         if -, bl_fl_mode then normalize(AR)
                  else bl_fl_norm(AR))
  else
    dwd(w) con dwd(w+addr) := float_rounded_to_71_bit(
      -         if -, bl_fl_mode then normalize(AR)
                  else bl_fl_norm(AR));
end;

gpu_sub       (= 31)
+++++++
AR := if addr = 0 then
  (if modus_sign = 0 then (AR - (dwd(w) con 0.0))
   -           else (AR + (dwd(w) con 0.0)) ) else
  -  (if modus_sign = 0 then (AR - (dwd(w) con dwd(w+addr)) )
      -           else (AR + (dwd(w) con dwd(w+addr)) ) );

ks            (= 51)
+++++++
no_op;

```

## 11. Description of the gpu\_declaration\_block

---

The gpu\_declaration\_block "gpu\_names" is copied into the code by the directive :

```
p.<:gpunames:>
```

and the code must be supplemented by an end-directive to the gpu\_declarations\_block :  
e.

before the last end-directive.

The names are :

```
u0      =    w0
u1      =    w1 ;    u5      =    x1
u2      =    w2 ;    u6      =    x2
u3      =    w3 ;    u7      =    x3

u8      =    indirect
u9      =    relative
u10     =    indirect and relative

u11     =    gpu_add
u12     =    gpu_arm
u13     =    gpu_cmv_a

u14     =    gpu_cmv_s
u15     =    gpu_init
u16     =    gpu_inv

u17     =    gpu_mla
u18     =    gpu_red
u19     =    gpu_sqrt

u20     =    gpu_stop
u21     =    gpu_str
u22     =    gpu_sub
```

## 12. Construction of the special gpu\_instructions.

---

An instruction consists of an operation and an address.

To ensure a correct address calculation a special construction must be used at least by relative addresses :

- a)  $(:uc + uw + ux + ua:) < 12 + \text{abs\_addr}$
- b)  $(:uc + uw + ux + ua:) < 12 + (:(:addr.:) < 12:) > 12$
- c)  $a0 = \text{addr.}$   
h.  
 $uc + uw + ux + ua, a0$   
w.

where uc is instruction ( $11 \leq c \leq 22$ )  
uw is word ( $0 \leq w \leq 3$ )  
ux is index ( $5 \leq x \leq 7$  or  $ux = 0$ )  
ua is addressmode ( $8 \leq a \leq 10$ )

Note that the address must not be negative and less than 4095 else the command part is destroyed.

When a negative relative address is used then it is recommended to use construction c with ua = if indirect then u10 else u9).

13. Block\_floating mode.

---

Block\_floating mode for the gpu is a kind of fixed\_point arithmetic.

The instruction gpu\_init delevers the modus bits and if needed the block\_floating exponent.

Block\_floating mode is set by : modus := 1 shift 21;  
 Block\_floating\_exponent is taken from : modus extract 12;

In block\_floating mode will the floating point stored by gpu\_inv and gpu\_str be stored with an exponent equal to or bigger than the block\_floating\_exponent, -

and all arithmetic on AR is followed with

no normalization shifts at all !!!  
 \*\*\*\*\*

IT IS RECOMMENDED TO USE ZERO AS BLOCK\_FLOATING\_EXPONENT  
 ++++++

Examples :

block\_floating\_exponent = 0;

number	is represented (octal)	with exponent
0.5	2000.0000.0000	0
1.0	2000.0000.0000	1
0.25	1000.0000.0000	0
2.0	2000.0000.0000	2
0.125	0400.0000.0000	0

It may then in block\_floating\_mode be necessary to scale all numbers with a known size by subtraction or addition to the exponent of the number.

In solution of normal equations this looks like :

A transposed = AT

obs-eq : A \* X = B

norm-eq : AT \* A \* X = AT \* B

scale\_factors : diagonal matrix S (n\*n)

transformation :

AT \* A \* X = AT \* B from right and left by S

S\*AT \* A\*S \* X = S\*AT \* B\*S

N \* X = S\*AT \* B\*S

scale\_factor to right-hand-side : R.

N \* X = S\*AT \* B\*S\*R \* (1/R)

N \* X = W \* (1/R)

triangulated NR \* X = WR \* (1/R)

back\_solution X = (NR\*\*(-1)) \* (WR \* (1/R))

## 14. Examples.

```

example 1 : code placed in bs_area gpu_code
===== product_sum result stored in a result area

gpu_code = slang

s. a20, b0, w.

    0          ; fill (evt. length)
    0          ; fill (evt. check_sum)

    j1.        b0.      ; goto code

a1: 1 < 23          ; modus = clear AR
a2: 0              ; addr a
a3: 0              ; addr b
a4: 0              ; repetitions
a5: 0              ; str addr
a6: 0              ; displacement

a20=a1.
h.
b0: 63<6 + 2<2 , a20 ; clear AR
w.
    d1. w2      a3.      ; load addr
    r1. w0      a4.      ; load repetitions
h.
58<6 + 2<4 + 0<2, 0 ; mla w2 ;
w.
    r1. w3      a5.      ; w3 := str addr
a19=a6.
h.
61<6 + 3<4 + 3<2, a19; str. w3 (a6.)
w.
    0          ; gpu_stop
i.
e.

; tail is missing. (see ref. 5).

end of example 1.
=====
```

example 2 : a code\_procedure moves the code to a zone  
=====

```

gpuinit = set 1
scope login gpuinit

(
if 9.yes
(gpuinit = slang names.yes list.yes xref.yes entry.no
gpuinit)
if 9.no
(gpuinit = slang entry.no
gpuinit)
lookup gpuinit
end
)

b.
d.

p.<:fpnames:>

l.

b. g1, e5
w.

s. a20, b20, j30, f20, g6

h.

g0 = 0           ; no of extrn.

e5:

g1:   g2,     g2      ; rel of last point, rel of last abs word

j13:   g0+13,   0      ; RS entry 13, last used
j30:   g0+30,   0      ;       - 30, save stack ref, saved w3
j17:   g0+17,   0      ;       - 17, index alarm
j6:    g0+ 6,   0      ;       - 6, end register expression

g2 = k-2-g1
w.
e0:   g0
0
s3,  date
s4,  time
e4:
```

```

w.

a1 = e5 + 6 ; modulus
a2 = a1 + 2 ; addr a
a3 = a2 + 2 ; addr b
a4 = a3 + 2 ; repetitions
a5 = a4 + 2 ; str addr
a6 = a5 + 2 ; displacement

; code for check and init of code in zone z,
el:
    r1. w2      (j13.)      ; get lastused
    ds. w3      (j30.)      ; save stackref and alarm addr

; check params
    d1  w1      x2+8          ;
    r1  w3      x1+h3          ; record base
    r1  w1      x1+h3+4        ; record length
    s1  w1      b1-1          ; if record length < code length
    j1. w3      (j17.)        ; then goto index alarm

; init code
f1:   al. w1      e5.-2      ; first of code
      ws w3      2           ; w3 := rel addr of zone buf
      rs. w3      f1.          ;
      al  w1      x1+6          ; first to store
      d1. w0      f2.          ; jump instr
      am.         (f1.)        ; zone displ
      ds  w0      x1          ;

; move code
f0:   al  w1      x1+4          ; increase pointer
      d1  w0      x1          ; load code
      am.         (f4.)        ; displ of zone
      ds  w0      x1          ; store in zone
      sh. w1      g3.          ;
      j1.         f0.          ;

; return to RS
      d1. w3      (      j30.) ; restore stackp and return
      j1.         (      j6.)  ; end reg expression

; work for code loader
f1:   0
f2:   0
f3:   j1.      f4          ; entry instr

c.     a6 - 1 - f3
      aw      0, r. (:a6 - f3:) > 1
z.

```

```

b0:
f4 = b0 - e5 - 4

a20 = a1.
h.
    63<6 + 2<2 , a20 ; clear AR
w.
    d1. w2      a3.      ; load addr
    r1. w0      a4.      ; load rep
    58<18+ 2<16 + 0      ; mla w2
    r1. w3      a5.      ; w3 := str addr
a19 = a6.
h.
    61<6 + 3<4 + 3<2, a19 ; str. w3 (a19.)
w.

```

```
w.
aw          0, r. 2
```

g3:

```
c. g3 - g1 -506
    m.code on segm too long
z.

c. 502 - g3 + g1, j1 -1, r. 252 -(:g3-g1:)>1 z.

<:gpuinit<0><0><0>:> ; alarm text

i.
e.      ; end of slang segm
```

```
g0:g1:  1           ; first tail, last tail
        0,0,0,0       ; name
        1<23 + e1 - e5 ; entry point for nllinitdpu
        2<18 + 8<12,0  ; spec bool proc, zone
        4<12 + e0 - e5 ; code proc start of extnr list
        1<12 + 0        ; 1 code segm, 0 bytes in core
```

d.

```
p.<:insertproc:>
e.
e.
e.
```

end of example 2.

=====

```

example 3 : call of gpu.
=====

real procedure scalar_product(a, b, length, res_tail_addr);
value                                length, res_tail_addr;
real array                           a, b;
integer                               length, res_tail_addr;
<* if res_tail_addr equals zero then is the result rounded
- to a normal real else is the result rounded to a double
- real result *>
begin
  zone code(128, 1, std_error),
    gpu(128, 1, gpu_error);
  real res;
  integer array field word;

  comment a procedure abs_addr(x), where x is type general
  -      is called. the result is an integer giving the
  -      absolute adress of th simple variable x or the
  -      absolute adress of the first element of an array
  -      variable x;

  open(gpu, 0, <:gpu1:>, -3);
  out_rec_6(gpu, 512);

  if example1 then
  begin
    open(code, 4, <:gpuode:>, 0);
    in_rec_6(code, 512);
    to_from(gpu, code);
    close(code, true);
  end
  else <* example 2 *>
    gpu_init(gpu);

  word      := 2;
  gpu.word(4) := abs_addr(a);
  gpu.word(5) := abs_addr(b);
  gpu.word(6) := n; <* length of vector *>
  gpu.word(7) := abs_addr(res);
  -           <* tail displacement = *>
  gpu.word(8) := if res_tail_addr = 0 then 0 else
  -                   res_tail_addr - gpu.word(7);

  close(gpu, true);

  scalar_product := res;

end scalar product;

end of example 3.
=====
```