

Raabo.



RC AS REGNECENTRALEN

SCANDINAVIAN INFORMATION PROCESSING SYSTEMS

RCSL No: 53-M18
Edition: December 1970
Author: Helge Elbrønd Jensen

Title: minimum

Keywords: RC 4000, Software, minimum, calculation of extrema, Algol procedure, ISO Tape

Abstract: The procedure, minimum, calculates extrema of a differentiable function in n variables. 17 pages.

1. Function and parameters

Let F denote a real, twice differentiable function in n variables, and suppose that the first order derivatives of F are given analytically (that is, as expressions depending upon the n variables). Suppose that in a given area the function is bounded below and has a minimum. From a reasonable good starting point the procedure finds this minimum by finding a point at which all the first order derivatives are zero (that is, smaller than a prescribed quantity).

Procedure head:

```
minimum(n, i, x, F, delta, eps, point);  
value n;  
integer i, n;  
real eps, F, delta;  
array x, point;
```

Call parameters:

n: the number of variables for the given function.

Call/Return parameters:

point: a real array point(1:n);
 at entry point contains the starting point for the
 procedure;
 at exit point contains the coordinates of the point
 at which the minimum is obtained;

eps: a real quantity affecting the precision to which the
 minimum is calculated. Consider the norm of the vector
 consisting of the first order derivatives. If this
 norm is smaller than eps, then the procedure will stop;
 at exit eps contains the norm of the vector described
 above.

Return parameters:

minimum: the value of the obtained minimum;

Other parameters:

F: a real procedure denoting the given function. In a program in which the procedure minimum is called, F must be declared in the following way:

```
real procedure F(x);  
array x;  
F:= the given expression;
```

delta: a real procedure delta(i, x) denoting for each i the partial derivative of F with respect to the variable x(i);
In a program in which the procedure minimum is called, delta must be declared in the following way:

```
real procedure delta(i, x);  
integer i;  
array x;  
delta:= case i of (... , ... , ...);
```

In the parenthesis there must be n expressions, where the i-th expression denotes the partial derivative of F with respect to the variable x(i);

2. The method

Let F denote a function in n variables, and let x denote the n-dimensional point with coordinates (x(1), x(2), ..., x(n)). F is said to have a minimum at a point x0, if there exist a small area including x0, in which the value of F at each point is greater than F(x0). Most of the various methods for finding a minimum for a function in n variables has one idea in common: They are all iterative processes based upon a roul, which for each point specifies a certain direction in which the next point of the process is to be found, and for each such direction specifies how to find the next point. Now, suppose that

the function is differentiable. By the gradient of F at the point x - denoted $\text{gradient}(x)$ - we understand the n -dimensional vector, which as the i -th coordinate has the partial derivative of F with respect to $x(i)$ at the point x .

The method used in the following program is essentially based upon to papers of A.A. Goldstein ((2), (3)). We suppose, that the function is twice differentiable and that the gradient is given analytically. It is well known, that the gradient will vanish at a minimum point.

Let the points of the iterative process be denoted $x_1, x_2, x_3, \dots, x_k, \dots$, where x_1 is given by the input array point.

For each k $f_i(x_k)$ denotes the n -dimensional vector which terminates the new direction.

We choose $f_i(x_1) = \text{gradient}(x_1)$.

For each k the number $h(k)$ is defined as:

$$h(k) = r \times \text{norm}(n, f_i(x_k)).$$

r is calculated at the beginning of the program in such a way that $h(1) < 1/5$.

norm is denoting the ordinary n -dimensional Euklidian norm.

Then the algorithm, at each point x_k , consists of the following two steps:

1. DIRECTION:

We compute an $n \times n$ matrix, which is an approximation to the matrix consisting of the second order derivatives of F .

For each j let $F(j)$ denote the vector, which has the j -th coordinate equal to 1 and the others equal to zero.

We then compute the matrix $Q(x_k)$ which has the j -th column equal to $(\text{gradient}(x_k + h(k) \times F(j)) - \text{gradient}(x_k))/h(k)$.

If the matrix $Q(x_k)$ is singular (it has no inverse) then we define the new direction $f_i(x_k)$ by

$$f_i(x_k) = \text{gradient}(x_k).$$

Suppose now, that $Q(x_k)$ has an inverse, which we denote $P(x_k)$.

If $(\text{gradient}(x_k), P(x_k) \times \text{gradient}(x_k)) > 0$

(where (\quad, \quad) denotes the ordinary innerproduct) then we define $f_i(x_k)$ by

$$f_i(x_k) = P(x_k) \times \text{gradient}(x_k).$$

If $(\text{gradient}(x_k), P(x_k) \times \text{gradient}(x_k)) \leq 0$
then we define $f_1(x_k)$ by
$$f_1(x_k) = \text{gradient}(x_k).$$

2. KONSTANT:

The next point in the process is now obtained on the form

$$x_k - g_k \times f_1(x_k)$$

where g_k is a constant calculated as follows:

Let $\text{product} = (\text{gradient}(x_k), f_1(x_k)).$

Let $f_1 = F(x_k).$

Let $f_2 = F(x_k - g_k \times f_1(x_k)).$

Then g_k is calculated such that

$f_2 < f_1$ and $(f_1 - f_2) < g_k \times \text{product}.$

It can be proved, by using the Taylor formula, that such a g_k always exists, and that x_k calculated in this way will converge to a minimum-point for $F.((2), (3)).$ From a numerical point of view however, g_k might fail to exist, and in this case the procedure will stop.

3. Accuracy, Time and Storage Requirements

Accuracy: As measure of accuracy we use the norm of the gradient. If the procedure succeeds, then at the end this norm is smaller than the call parameter $\text{eps}.$

Time: This depends on the wanted accuracy and first of all on the problem in question, so it is not possible to give general rules for this. (See 4. Test and Discussion).

Storage requirements: 10 segments of program

Typographical length: 248 lines.

4. Test and Discussion

The procedure have been tested on several functiones among which we describe the two most difficult problems:

1. Minimising the function in two variables

$$F = 100 \times (x(2) - x(1) \times 2) \times 2 + (1 - x(1)) \times 2$$

2. Finding a solution to the following three non-linear equations:

$$\sin(x(1) \times 2) + \exp(x(2)) \times x(3) - 4 = 0$$

$$x(1) + x(2) + x(3) - 3 = 0$$

$$x(1) + x(2) \times 2 + x(3) \times 3 - 14 = 0$$

This is done by minimising the square-sum of the three equations.

First we consider the problem 1:

The function F has minimum at the point (1, 1) with functionvalue 0.

Starting at the point $x(1) = -1.2$ and $x(2) = 1$ and using different values of the term eps, the following results were obtained:

	Value of eps			
	10^{-4}	10^{-6}	10^{-8}	10^{-10}
Minimum	0.999999592 0.999999183	0.999999592 0.999999183	1.000000000 1.000000000	1.000000000 1.000000000
Fc.-value	0.000000000	0.000000000	0.000000000	0.000000000
Gr.-norm	7.3×10^{-5}	4.2×10^{-7}	5.1×10^{-8}	1.7×10^{-10}
Ex.-time	0.76	0.73	0.75	0.75

(the execution time is in seconds).

It follows, that the procedure succeeds in all 4 situations, and that smaller values of eps does not affect the execution time. This last observation however can not be stated in general, (see below under problem 2).

Using $\text{eps} = 10^{-8}$ and using different starting points the following results were obtained:

Starting-point	-1.200000000 1.000000000	0.000000000 1.000000000	-0.500000000 -0.500000000	2.000000000 0.250000000
Execution-time	0.75	0.49	0.71	0.80

In all 4 situations the minimum was obtained at the point:

1.000000000
1.000000000

with the functionvalue 0.000000000 and gradient norm $5.1 \cdot 10^{-8}$.

Again the procedure succeeds in all 4 situations.

Next, consider the problem 2 in three variables. Starting at the point $x(1) = 0$, $x(2) = 0$, $x(3) = 2.5$ and using different values of the term ϵ , the following results were obtained:

	Value of ϵ			
	10^{-4}	10^{-6}	10^{-8}	10^{-10}
Minimum	0.097831561 0.512917627 2.389250732	0.097830233 0.512919004 2.389250762	0.097830224 0.512919014 2.389250762	0.097830224 0.512919014 2.389250762
Fc.-value	0.000000000	0.000000000	0.000000000	0.000000000
Gr.-norm	$5.8 \cdot 10^{-5}$	$3.9 \cdot 10^{-7}$	$3.2 \cdot 10^{-8}$	$3.2 \cdot 10^{-8}$
Ex.-time	1.96	2.55	3.81	3.97

It follows, that the procedure succeeds in the first three situations, but that it is not possible to make the gradient norm smaller than $3.2 \cdot 10^{-8}$, so in this sense the procedure does not succeed in the last situation. In this case smaller values of ϵ gives greater execution time, even if the obtained minimumpoints are practically the same in the last three cases.

Using $\epsilon = 10^{-8}$ and using different starting points the following results were obtained:

Starting-point	0.000000000 0.000000000 2.500000000	0.000000000 0.000000000 1.000000000	0.500000000 1.000000000 2.000000000	1.000000000 1.000000000 1.000000000
Execution-time	3.81	1.97	3.09	2.45

In all 4 situations the minimum was obtained at the point:

0.097830223
0.512919014
2.389250762

with the functionvalue 0.000000000 and gradient norm $3.2 \cdot 10^{-8}$
It follows, that the procedure succeeds in all 4 situations.

Example

Consider the function

$$F = 100 \times (x(2) - x(1) \times 2) \times 2 + (1 - x(1)) \times 2$$

Starting at the point $x(1) = -1.2$ and $x(2) = 1$ the following program might be used to find the minimum of F:

Testprogram

```
begin
  integer i, j;
  real a, eps;
  array x, point(1:2);

  real procedure F(x);
  array x;
  F:= 100 × (x(2) - x(1)×2) × 2 + (1 - x(1)) × 2;

  real procedure delta(i, x);
  integer i;
  array x;
  delta:= case i of (-400×x(1) × (x(2) - x(1)×2) - 2 × (1 - x(1)),
                    200 × (x(2) - x(1)×2));

  point(1):= -1.2; point(2):= 1; eps := 10-8;
  a:= minimum(2, i, x, F(x), delta(i, x), eps, point);
  write(out, <:Minimum obtained at the point <10>:>);
  for j:= 1 step 1 until 2 do
  write(out, <:<10>:>, <<-dddd.dddddddd>, point(j));
  write(out, <:<10><10> Minimumvalue =:>, <<-dddd.dddddddd>, a);
  write(out, <:<10><10> Gradient norm=:>, <<-d.d10-dd>, eps);
end;
```


This will give the following output:

Minimum obtained at the point

1.000000000

1.000000000

Minimumvalue = 0.000000000

Gradient norm = 1.0⁻⁸

end

In the program we use a boolean procedure inverse to find the inverse (if it exist) of an $n \times n$ matrix.

The procedure is based upon Simpel Gaussian illimination and is only introduced in order to make the program complete. One could use any other procedure of this sort, for ex. decompose-solve from RC mathematical procedure library.

Since a minimum of the function F is a maximum of the function $-F$, the procedure will of course be able to find maximum as well as minimum.

5. References

- (1) D. Fletcher and M.J.D. Powell:
A rapidly convergent descent method for minimisation.
Comput. Journal 6 p. 163-168 (1963)
- (2) A.A. Goldstein: On steepest descent.
Journal Siam Control Vol. 3 No 1 p 147-151 (1965)
- (3) A.A. Goldstein and J.F. Puce:
An effective algorithm for minimisation
Numerische Mathematik 10 p. 184-189 (1967)
- (4) E. Isaacson and H.B. Keller:
Analyses of numerical Methods
John Wiley and Sons, Inc. (1966)

6. Algol text

```
minimum = set 10  
minimum = algol  
external
```

```
real procedure minimum(n,i,x,F,delta,eps,point);  
value n;  
integer i,n;  
real eps,F,delta;  
array x,point;
```

```
begin  
integer j;  
real h,g,g1,gamma,r,f1,f2,f3,product,k,s;  
array psi,y,z,b(1:n),p,q(1:n,1:n);
```

```
real procedure norm(n,a);  
value n;  
integer n;  
array a;  
begin  
comment this is the ordinary norm in the n-dimensional Euklidian  
space;  
real h;  
h:=0;for i:=1 step 1 until n do h:=h+a(i)×2;  
norm:=sqrt(h);  
end;
```

```
real procedure innerproduct(n,a,b);  
value n;  
integer n;  
array a,b;  
begin  
comment this is the ordinary innerproduct in the n-dimensional
```

```
Euklidian space;
real h;
h:=0; for i:=1 step 1 until n do h:=h+a(i)Xb(i);
innerproduct:=h;
end;
```

```
procedure equal(n,a,b);
value n;
integer n;
array a,b;
begin
  comment the procedure identifies two arrays;
  for i:=1 step 1 until n do b(i):=a(i);
end;
```

```
boolean procedure inverse(n,a,b);
value n;
integer n;
array a,b;
  comment the procedure finds the inverse ( if it exists ) of the
  matrix a by Gaussian illimination.If the inverse exist,it is
  stored in b.If the inverse does not exist,inverse is false;
begin
  integer i,j,k,m,pivotnr;
  real pivot,s;
  array c(1:n,1:n),x(1:n),d(1:n);

  inverse:=true;
  for m:=1 step 1 until n do
  begin
    comment for each m one is solving the linear system,which on the
    wright side has the m-th column in the unit-matrix,and on the left
    side the given matix as coefficientmatrix and the m-th column in
    the wanted inverse as unknown;
```

```
for j:=1 step 1 until n do
for i:=1 step 1 until n do c(i,j):=a(i,j);
for i:=1 step 1 until n do d(i):=(if i=m then 1 else 0);
for k:=1 step 1 until n-1 do
begin
  comment among the last n-k+1 equations one is finding the equation,
  which has the numerical largest coefficient in x(k);
  pivot:=0; pivotnr:=0;
  for i:=k step 1 until n do if abs(c(i,k))>pivot then
  begin pivot:=c(i,k); pivotnr:=i; end;
  if pivot=0 then begin inverse:=false; goto END;end;
  comment if pivot=0 then the given matrix has determinant 0 and
  consequently no inverse;
  if pivotnr<>k then
  begin
    comment equation number k is replaced by equation number pivotnr
    and vica versa;
    s:=d(k); d(k):=d(pivotnr); d(pivotnr):=s;
    for j:=k step 1 until n do
    begin
      x(j):=c(k,j); c(k,j):=c(pivotnr,j); c(pivotnr,j):=x(j);
    end;
  end if pivotnr<>k;
  for i:=k+1 step 1 until n do
  begin
    comment x(k) is calculated from the k-th equation, and the
    expression inserted in the following n-k equations;
    d(i):=d(i)-d(k)*c(i,k)/c(k,k);
    for j:=k+1 step 1 until n do
    c(i,j):= c(i,j)-c(i,k)*c(k,j)/c(k,k);
  end;
end k;
if c(n,n)=0 then begin inverse:=false; goto END;end else
x(n):=d(n)/c(n,n);
for i:=n-1 step -1 until 1 do
begin
  comment for each i x(i) is calculated from the equation
  c(i,i)*x(i) + c(i,i+1)*x(i+1) + . . . +c(i,n)*x(j) = d(i),
```

```
where  $x(i+1), \dots, x(n)$  are known;  
s:=0; for j:=n step -1 until i+1 do s:=s+c(i,j)*x(j);  
x(i):=(d(i)-s)/c(i,i);  
end;  
for i:=1 step 1 until n do b(i,m):=x(i);  
end m;  
END: end;
```

```
procedure search(n,g,y,psi,f2);  
value n,g;  
integer n;  
real g,f2;  
array y,psi;  
begin  
comment the procedure finds the value of the function to be  
minimised, that is  $k \times F$ , at the point obtained from y by going the  
distance g in the direction -psi;  
for i:=1 step 1 until n do x(i):=y(i)-g*psi(i);  
f2:=k*F;  
end;
```

```
equal(n,point,x);equal(n,x,y); k:=1;  
for i:=1 step 1 until n do psi(i):=delta;  
comment psi is the gradient of F at the starting point;  
equal(n,psi,b); h:=product:=norm(n,psi);  
if h<1 then r:=1/5 else r:=1/(5*h);
```

KONSTANT:

```
; comment at each step of the iterative process the procedure will  
goto KONSTANT and run through the following. A point y and a  
direction psi is given, and the problem is to find a konstant g  
such that the point  $y-g \times \text{psi}$  can be used as the next point;  
h:=norm(n,psi); equal(n,y,x);  
if h/product<1/10 then
```



```
begin
  comment psi is too small relativ to the gradient which implies,
  that the greatest possible progress is too small. We therefore
  consider the function  $k \times F$ , where  $k$  is defined below;
   $k := (h / \text{product}) \times (1/n)$ ;
  for  $i := 1$  step 1 until  $n$  do  $\text{psi}(i) := (1/k) \times n \times \text{psi}(i)$ ;
  for  $i := 1$  step 1 until  $n$  do  $b(i) := k \times \text{delta}$ ;
   $h := \text{norm}(n, \text{psi})$ ;
  if  $h < 1$  then  $r := 1/5$  else  $r := 1/(5 \times h)$ ;
end;
 $h := r \times h$ ;  $f1 := k \times F$ ;
comment  $h$  is used below as the small quantity in the approximation
of the second order derivatives of  $F$ ,  $r$  is introduced in order to
insure, that this quantity is not too big at the beginning;
 $\text{product} := \text{innerproduct}(n, b, \text{psi})$ ;
 $g := 1$ ;  $g1 := 0$ ;
 $\text{search}(n, 1, y, \text{psi}, f2)$ ;
if  $f1 - f2 \geq 1/4 \times \text{product}$  then
begin
   $f1 := f2$ ;  $\text{equal}(n, x, y)$ ; goto DIRECTION;
end;
comment in this case we use  $g = 1/4$  and the next point is
therefore obtained as  $y - 1/4 \times \text{psi}$ ;
 $s := (\text{if } s < 1 \text{ then } 10^{-10} \text{ else } 1/s \times 10^{-10})$ ;
for  $g := g/2$  while  $f1 \leq f2$  do
begin
   $\text{search}(n, g, y, \text{psi}, f2)$ ;
  if  $g < s$  then begin  $\text{equal}(n, y, x)$ ; goto END; end;
end;
comment if  $g$  is smaller than  $s$  (see the definition of this term)
then the next point of the process will be practically equal to
the present, and we must therefore conclude, that the procedure is
unable to make further progress;
 $g := 2 \times g$ ;  $\text{equal}(n, x, z)$ ;
if  $(f1 - f2) < g \times \text{product}$  then goto SECOND else
begin
  comment in this case the functionvalue at  $y - g \times \text{psi}$  is smaller
  than  $f1$ , but the condition  $f1 - f2 < g \times \text{product}$  is not satisfied and
```

therefore g is too small;

g1:=g; g:=2×g;

FIRST:

g:=(g1+g)/2;

search(n,g,y,psi,f2);

if f1<f2 then goto FIRST else

begin

if (f1-f2)<g×product then

begin equal(n,x,y); f1:=f2; goto DIRECTION; end else

begin g:=2×g-g1; g1:=(g+g1)/2; goto FIRST; end;

end;

end;

SECOND:

; comment in this case the functionvalue at y-g×psi is smaller than f1 and the condition $f1-f2 < g \times \text{product}$ is satisfied. We therefore look for a smaller g for which this condition is satisfied and with a smaller functionvalue than before;

g:=(g1+g)/2; search(n,g,y,psi,f3);

if f2<=f3 then

begin equal(n,z,x);equal(n,x,y); f1:=f2; goto DIRECTION;

end else

begin

if (f1-f3)<g×product then

begin f2:=f3; equal(n,x,z); goto SECOND;end

else goto THIRD;

end;

THIRD:

; comment in this case the functionvalue is smaller than before, but the condition mentioned before is not satisfied,so g is too small;

g:=2×g-g1; g1:=(g+g1)/2; g:=(g+g1)/2;

search(n,g,y,psi,f3);

if (f1-f3)>=g×product then goto THIRD else

begin

if f3>=f1 then goto THIRD else

begin equal(n,x,y); f1:=f3; goto DIRECTION; end;

end;

DIRECTION:

```
; comment at each step of the iterativ process the procedure will
goto DIRECTION and run through the following. A point x is given
and the problem is to determine the direction in which the next
point is to be found;
for i:=1 step 1 until n do b(i):=k*delta;
product:=norm(n,b);
if product<k*eps or product<10-10 then goto END;
comment if product<k*eps then the wanted accuracy is obtained.
if product<10-10 then in most situations it will be meaningless
to look for further progress;
for j:=1 step 1 until n do
begin
comment an approximation to the matrix consisting of the second
order derivatives of kxF is calculated and the result stored in q;
for i:=1 step 1 until n do x(i):=(if i=j then y(i)+h else y(i));
for i:=1 step 1 until n do p(i,i):=k*delta;
for i:=1 step 1 until n do q(i,j):=(p(i,i)-b(i))/h;
end j;
if -,inverse(n,q,p) then goto STEEPEST else
begin
comment if the inverse of q exist, then the vector psi is
obtained by multiplying the inverse matrix with the gradient;
for i:=1 step 1 until n do
begin
psi(i):=0; for j:=1 step 1 until n do
psi(i):=psi(i)+p(i,j)*b(j);
end;
end;
if innerproduct(n,psi,b)<=0 then goto STEEPEST else
goto KONSTANT;
comment if innerproduct(n,psi,b)<=0 then we can not be sure to
find a point with smaller functionvalue in the direction psi,
and therefore psi can not be used. If the innerproduct is >0
then psi is the new direction;
```

STEEPEST:

```
equal(n,b,psi); goto KONSTANT;
```

```
comment the gradient is used as the new direction;
```

```
END:
```

```
; comment the present value of the relevant quantities are stored  
in the return parameters;
```

```
for i:=1 step 1 until n do b(i):=delta;
```

```
minimum:=F; eps:=norm(n,b);
```

```
for i:=1 step 1 until n do point(i):=x(i);
```

```
end;
```

```
end;
```